

# **P2P Applications**

Reti di Elaboratori

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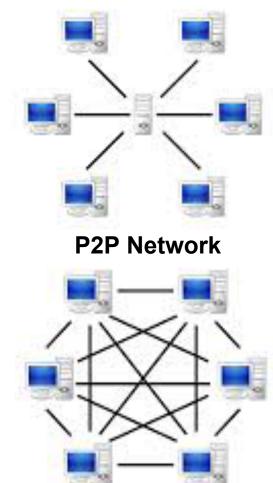
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#### **Server-based Network**

### Peer-to-peer networks

- A type of network in which each workstation has equivalent capabilities and responsibilities
- Differs from client/server architectures, in which some computers are dedicated to serving the others



# P2P Paradigm

Late 80's

- napster
- Became popular in 1999-2001 thanks to
- Napster was shut down by court order in 2001 due to copyright violation
- New P2P clients were developed: Gnutella, Kazaa, BitTorrent
- As of today, 43-70% of Internet traffic is generated by P2P applications (Feb 2009)

#### P2P file sharing

- Computer systems connected via Internet form a network of peers among which digital documents and files are distributed and shared
- P2P programs search for other connected computers on a P2P network and locate the desired content
- Each peer can both download and upload files from/to other peers
- How to keep track of which peers hold a given content?
  - Centralized vs decentralized solutions

#### **Understanding P2P protocols**





#### Spetify: Overview

- Spotify is a <u>peer-assisted</u> on-demand music streaming service
- Large catalog of music (over 15 million tracks)
- Available in more than 32 countries (USA, Europe, Asia)
- Very popular (over 10 million users)
- Only ~ 250ms playback latency on average!



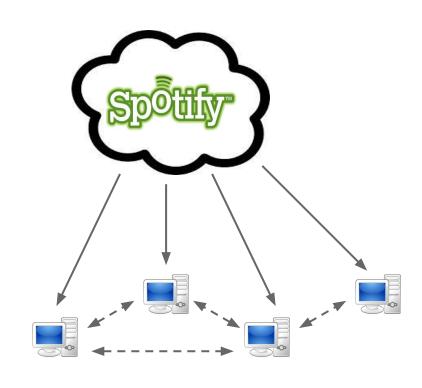
#### Spetify: Overview (cont'd)

- Spotify uses a proprietary protocol, but:
  - some of its internals have been described by researchers working at Spotify (<a href="http://www.csc.kth.se/~gkreitz/spotify-p2p10/">http://www.csc.kth.se/~gkreitz/spotify-p2p10/</a>)
  - a third-party OSS alternative client has been released (<a href="http://despotify.sourceforge.net/">http://despotify.sourceforge.net/</a>)
  - update: the alternative client is not compatible anymore with the new protocol



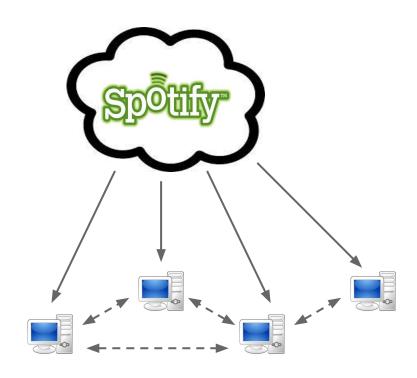
#### Spetify: Architecture (cont'd)

- Spotify uses a *hybrid* content distribution method, combining:
  - a client-server access model
  - o a P2P network of clients
- Main advantage: only ~ 8.8% of music data comes from the spotify servers! The rest is shared among the peers (although mobile devices do not participate to the P2P network)
- Possible drawbacks:
  - playback latency (*i.e.*, the time the user has to wait before the track starts playing)
  - complex design



#### Spetify: Architecture (cont'd)

- As we are about to see, Spotify uses a number of good design choices in order to:
  - keep the playback latency low
  - simplify its design (with respect to pure P2P systems)



### Spetify: Load Balancing

- To balance the load among its servers (at least 2, one in London, one in Stockholm), a peer randomly selects which server to connect to.
- Each server is responsible for a separate and independent P2P network of clients.
  - advantage: does not require to manage inconsistencies between the servers' view of the P2P network
  - advantage: the architecture scales up nicely (at least in principle). If more users join Spotify and the servers get clogged, just add a new server (and a new P2P network)
- To keep the discussion simple, we assume there is only one server.

#### Sp**stify: P2P Network**

- Spotify uses an unstructured P2P overlay topology.
  - the network is built and maintained by means of trackers (similar to BitTorrent)
  - no super nodes with special maintenance functions (as opposite to Skype)
  - o **no** Distributed Hash Table to find peers/content (as opposite to eDonkey)
  - no routing: discovery messages get forwarded to other peers for one hop at most

#### Advantages:

- keeps the protocol simple
- keeps the bandwidth overhead on clients down
- reduces latency
- This is possible because Spotify can leverage on a *centralized* and *fast* backend (as opposite to the completely distributed P2P networks)

### Spetify: Caching

- Spotify clients store the already played tracks in a cache. By default, the cache uses at most 10% of disk space (capped to 10GB, but never less than 50MB).
- Around 56% of clients have a maximum cache size of 5GB.
  - advantage: reduces the chances that a client has to re-download already played tracks.
  - advantage: increases the chances that a client can get a track from the P2P network (lower load on the Spotify servers).
  - drawback: impacts on the users' disk
    - an LRU cache-eviction policy is used that removes the Least Recently
       Used (i.e., played) track.
    - caches are large (as compared to the typical track size), so this is not a big deal.

#### **Spetify: Sharing Tracks**

- A client cannot upload a track to its peers unless it has **the whole** track
  - advantage: this choice greatly simplifies the protocol and keeps the overhead low, as clients do not have to communicate (to their peers or to the server) what parts of a track they have.
  - o **drawback**: reduces the number of peers a client can download a track from (*i.e.*, slower downloads).
    - tracks are small though (few MB each), so this has a limited effect

#### **Spetify: Locating Peers**

- There are two ways a client can locate the peers:
  - ask the server
  - ask the other peers

### Sp**stify: Locating Peers (tracker)**

- The server maintains a **tracker**, similarly to BitTorrent.
  - as opposite to other systems, however, the server does not keep track of all the peers who can serve each track
  - rather, it keeps a list of the ~ 20 most recent clients that played each track
  - clients do not report to the server the content of their caches!

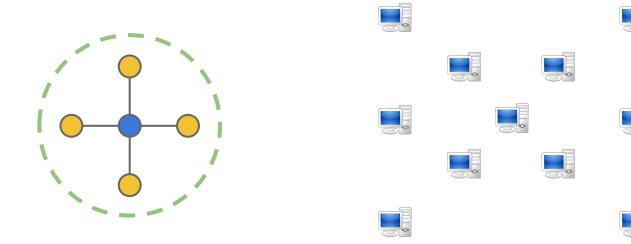
#### Advantages:

- less resources on the server side
- simplifies the implementation of the tracker
- **Drawback**: only a **fraction** of the peers can be located through the tracker
  - this is not a big issue, since clients can ask the other peers (next slide)

### **Sps**tify: Locating Peers (P2P)

- Each client has a set of neighbors (other clients) in the P2P network.
  - these are the peers the client has previously uploaded a track to, or has previously downloaded a track from
- When a new track has to be downloaded, a client can search its neighborhood for peers that have stored in their cache
- The peers can, in turn, forward the search request to their own peers in the network
  - the process stops at distance 2 in the overlay network
- each query has a unique ID, to allow ignoring duplicate queries

### Sp⊜tify: Locating Peers (P2P) (cont'd)



#### Spetify: Neighbor Selection

- A client uploads to at most 4 peers at any given time
  - helps Spotify behave nicely with concurrent application streams (e.g., browsing)
- Connections to peers do not get closed after a download/upload
  - advantage: reduces time to discover new peers when a new track has to be played
  - drawback: keeping the state required to maintain a large number of TCP connections to peers is expensive (in particular for home routers acting as stateful firewall and Network Address Translation (NAT) devices)
- To keep the overhead low, clients impose both a soft and a hard limit to the number of concurrent connections to peers (set to 50 and 60 respectively)
  - when the soft limit is reached, a client stops establishing new connections to other peers (though it still accepts new connections from other peers)
  - when the hard limit is reached, no new connections are either established or accepted

#### Spetify: Neighbor Selection (cont'd)

- When the soft limit is reached, the client starts pruning its connections, leaving some space for new ones.
- To do so, the client computes an utility of each connected peer by considering, among the other factors:
  - the number of bytes sent (received) from the peer in the last 60 (respectively
     10) minutes
  - the number of other peers the peer has helped discovering in the last 10 minutes
- Peers are sorted by their utility, and the peers with the least total scores are disconnected.

#### Sp⊜tify: Playing a Track

- Around 61% of tracks are played in a predictable order (*i.e.,* the previous track has finished, or the user has skipped to the next track)
  - playback latency can be reduced by predicting what is going to be played next.
- The remaining 39% are played in random order (e.g., the user suddenly changes album, or playlist)
  - predicting what the user is going to play next is too hard. Playback latency may be higher

#### **Spetify: Random Access**

- When tracks are played in an unpredictable (random) order, fetching them
  just using the P2P network would negatively impact the playback delay.
- Why?
  - searching for peers who can serve the track takes time (mostly because of multiple messages need to be exchanged with each peer)
  - some peers may have poor upload bandwidth capacity (or may be busy uploading the track to some other client)
  - a new connection to a peer requires some time before start working at full rate (check out the lectures about TCP congestion control)
  - P2P connections are unreliable (e.g., may fail at any time)

#### Sp⊜tify: Random Access (cont'd)

- How to solve the problem?
- Possible solution: use the fast Spotify Content Delivery Network (CDN)
  - drawback: more weight on the Spotify CDN (higher monetary cost for Spotify.. and possibly to its users too)
- Better solution: use the Spotify CDN asking for the first 15 seconds of the track only.
  - advantage: this buys a lot of time the client can use to search the peer-topeer network for peers who can serve the track.
  - advantage: the Spotify CDN is used just to recover from a critical situation (in this case, when the user has started playing a random track)

#### Spetify: Sequential Access

- When users listen to tracks in a predictable order (*i.e.*, a playlist, or an album), the client has plenty of time to **prefetch** the next track before the current one finishes.
- **Problem**: you don't really know whether the user is actually going to listen to the next track or not. If the user plays a random track instead of the predicted one, you end up having wasted bandwidth resources.
- **Solution**: start prefetching the next track only when the previous track is about to finish, as Spotify has experimentally observed that:
  - when the current track has only 30 seconds left, the user is going to listen to the following one in 92% of the cases.
  - when 10 seconds are left, the percentage rises to 94%
- The final strategy is:
  - 30 seconds left: start searching for peers who can serve the next track
  - 10 seconds left: if no peers are found (critical scenario!), use the Spotify CDN

### Spetify: Regular Streaming

- Tracks are split in 16KB chunks.
- A track can be simultaneously downloaded from the CDN and the P2P network.
- If both CDN and P2P are used, the client never downloads from the Spotify
   CDN more than 15 seconds ahead of the current playback point.
- To select the peers to request the chunks from, the client sorts them by their expected download times and greedily requests the most urgent chunk from the top peer.
  - expected download times are computed using the average download speed received from the peers
  - if a peer happens to be too slow, another peer is used

### Sp⊜tify: Regular Streaming (cont'd)

- The client continuously monitors the **playout buffer** (*i.e.*, the portion of the song that has been downloaded so far but not already played)
- If the buffer becomes too low (< 3 seconds) the client enters an emergency mode, where:
  - it stops uploading to the other peers
    - this is especially useful in asymmetric connections (e.g., aDSL), whose download capacity is negatively affected by concurrent uploads (check out the lectures on TCP)
  - it uses the Spotify CDN
    - this helps in the case the client fails to find a reliable and fast set of peers to download the chunks from

#### **Understanding P2P protocols**





#### **BitTorrent**



- P2P file distribution system
- Designed and implemented (Python) by Bram Cohen in 2001
- Dozen of free clients
- January 2012: 150 million active users
- Used to distribute large amounts of data over the Internet: not only media content, but also Linux distributions, scientific data sets, ...



- **BitTorrent overview**
- The set of all peers participation in the distribution of a file is called a torrent
- A separate torrent for each file
- Peers simultaneously upload and download pieces of file within the torrent
- The set of all active peers in a torrent is called the swarm

#### Two types of peers

For each *torrent* the set of active peers is divided into:



Seeds: clients that have a complete copy of the file and that continue to serve other peers

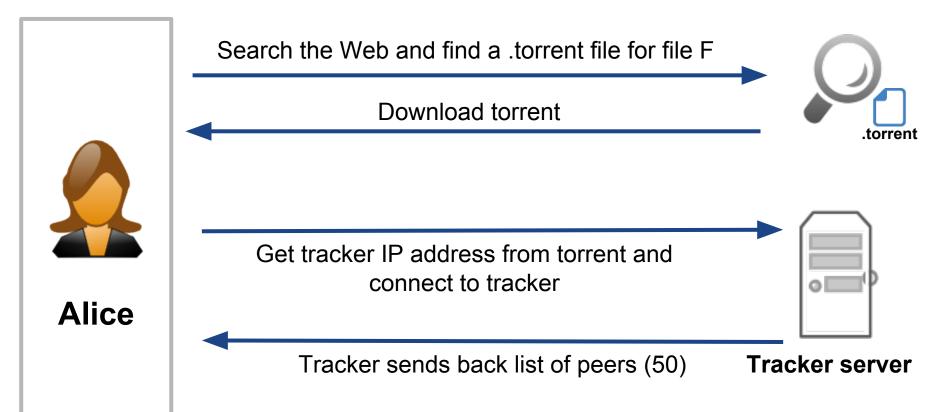


Leechers: clients that are still downloading the file (Alice)

#### How to download a file?

- Users need to discover which peers hold a copy of the file (at least a seeder!)
- Search for a .torrent file on the Web
- Torrent file include the address of a centralized server (the *tracker*) that helps peers finding each other
- Connect to the tracker and receive the list of peers having a copy of the file

### Discovering peers for a file F



#### The Tracker



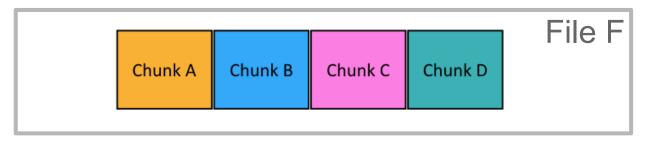
- Not involved in the actual distribution of files!
- Keeps information about peers currently active
- Peers report their state to the tracker every 30 minutes, and when joining or leaving the network
- New clients receive from the tracker the IP address of 50 randomly chosen active peers

#### **Contacting peers**

- Once received the list of IP addresses from the tracker, Alice tries to establish a TCP connection with each of them
- Peer set: peers to which Alice is connected
- It changes over time!
- If nodes in the peers set become less than 20, Alice contacts the tracker again to obtain a new list

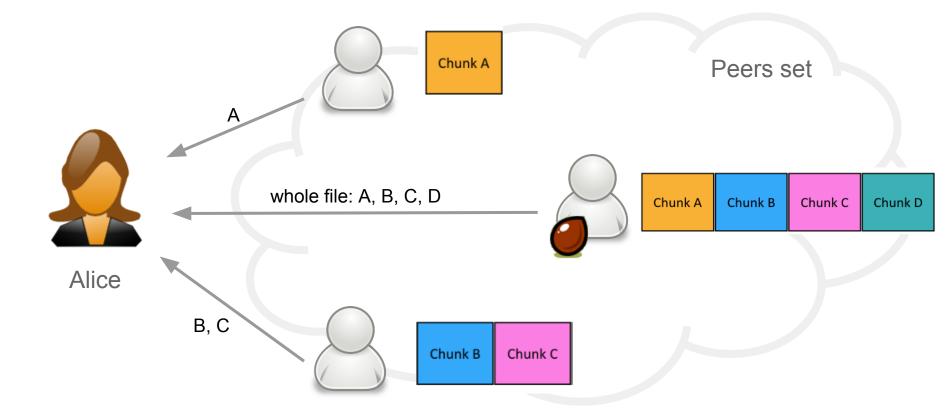
#### File chunks

In BitTorrent files are divided into pieces (chunks) of size between
 64 KB and 1 MB (typically 256 Kb)



- When Alice enters the torrent for file F, she has no chunks
- Each peer in the peers set have a subset of chunks from F
- Alice periodically asks each node in the peers set for the list of chunks they have

#### Peers send their chunks list to Alice



#### **Downloading chunks**



- Alice downloads chunks from multiple peers and keeps track of the download rate from each of them
- In which order file chunks are downloaded?
- (Local) rarest first: based on the chunks list received by her peers set, Alice determines which chunk (among those she does not have) is the rarest one in her peers set

#### **Uploading chunks**

- As soon as Alice downloads her first chunk, she can start uploading to other peers
- Alice has a limited number of upload slots to allocate to other peers
- How to choose which peers to serve?
- Tit for tat: exchanging upload bandwidth for download bandwidth

## **Trading chunks**



- Alice continuously measures her download rate from the other peers
- She uploads chunks to the 4 peers from which she is downloading at the <u>highest rate</u>
- Every 10 seconds she recalculates the four top peers
- In addition, every 30 seconds she picks a peer at random and uploads chunks to her





- The five peers to which Alice uploads are said to be unchoked
- All the other peers in the swarm are choked,
   i.e., they do not receive any chunk from Alice
- Unchoking a random peer every 30 seconds
  - allows to discover better partners
  - ensures that newcomers get a chance to join the swarm

#### **BitTorrent Pros and Cons**

#### Pros:

- Proficiently uses partially downloaded files
- Discourage free-loading by rewarding fast uploaders
- Works well for hot content

#### Cons:

- High latency and overhead for small files
- Less useful for unpopular content
- Does not support streaming
- Leech problem
- Not a pure P2P protocol: single point of failure

#### Want to know more?

Incentives Build Robustness in BitTorrent

Bram Cohen bram@bitconjurer.org

May 22, 2003

#### Abstract

The BitTorrent file distribution system uses tit-fortat as a method of seeking pareto efficiency. It achieves a higher level of robustness and resource utilization than any currently known cooperative technique. We explain what BitTorrent does, and how economic methods are used to achieve that goal. each peer's download rate be proportional to their upload rate. In practice it's very difficult to keep peer download rates from sometimes dropping to zero by chance, much less make upload and download rates be correlated. We will explain how BitTorrent solves all of these problems well.

#### 1.1 BitTorrent Interface

Incentives Build Robustness in BitTorrent, Bram Cohen Workshop on Economics of Peer-to-Peer Systems, June 2003