

Reti di Elaboratori

Corso di Laurea in Informatica
Università degli Studi di Roma “La Sapienza”
Canale A-L
Prof.ssa Chiara Petrioli

Parte di queste slide sono state prese dal materiale associato al libro
Computer Networking: A Top Down Approach , 5th edition.

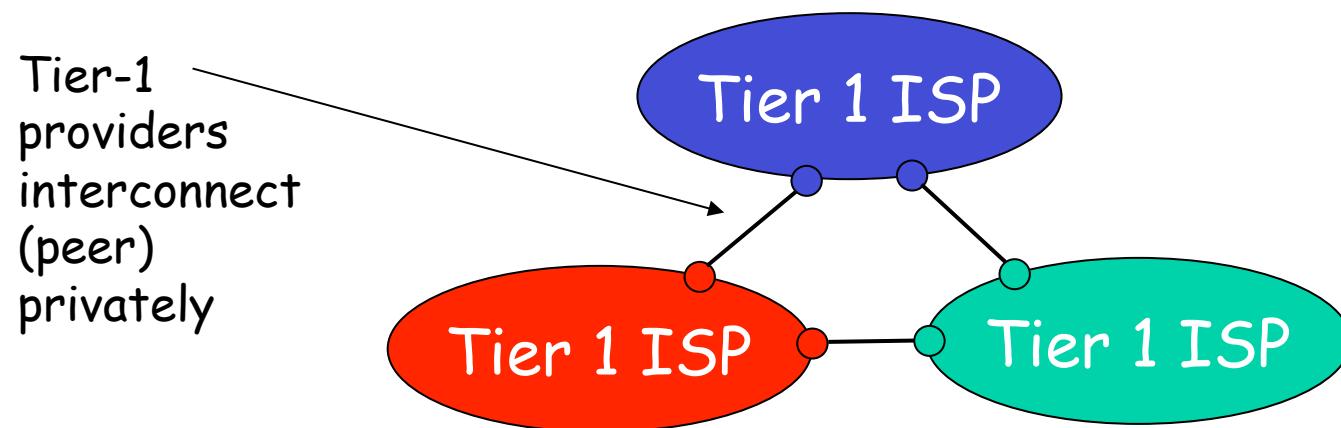
All material copyright 1996-2009

J.F Kurose and K.W. Ross, All Rights Reserved

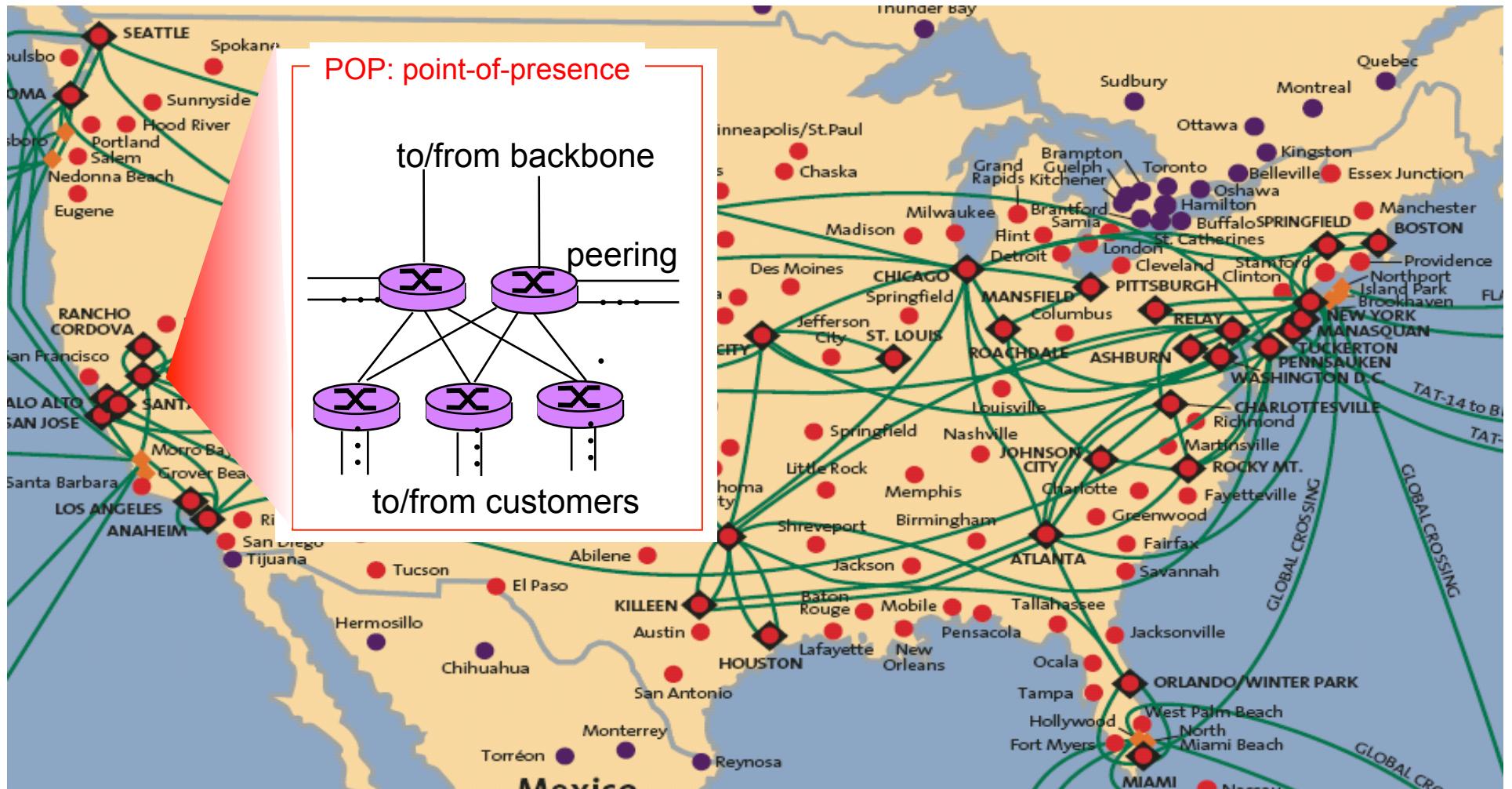
Thanks also to Antonio Capone, Politecnico di Milano, Giuseppe Bianchi and
Francesco LoPresti, Un. di Roma Tor Vergata

Internet structure: network of networks

- roughly hierarchical
- at center: “tier-1” ISPs (e.g., Verizon, Sprint, AT&T, Cable and Wireless), national/international coverage
 - treat each other as equals

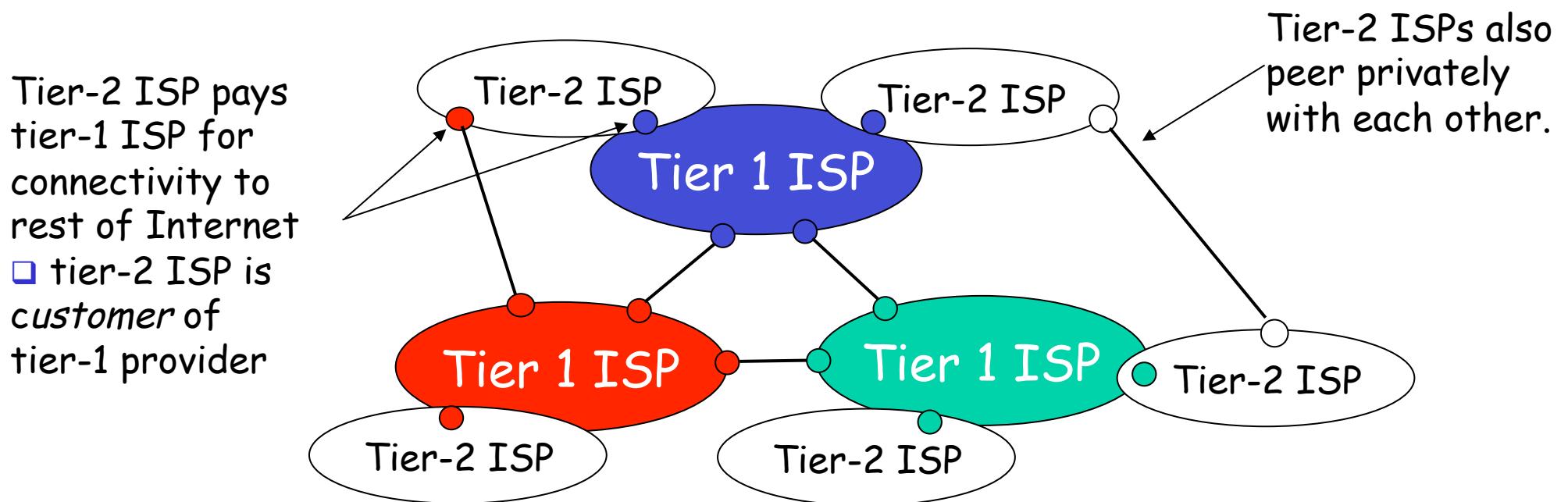


Tier-1 ISP: e.g., Sprint



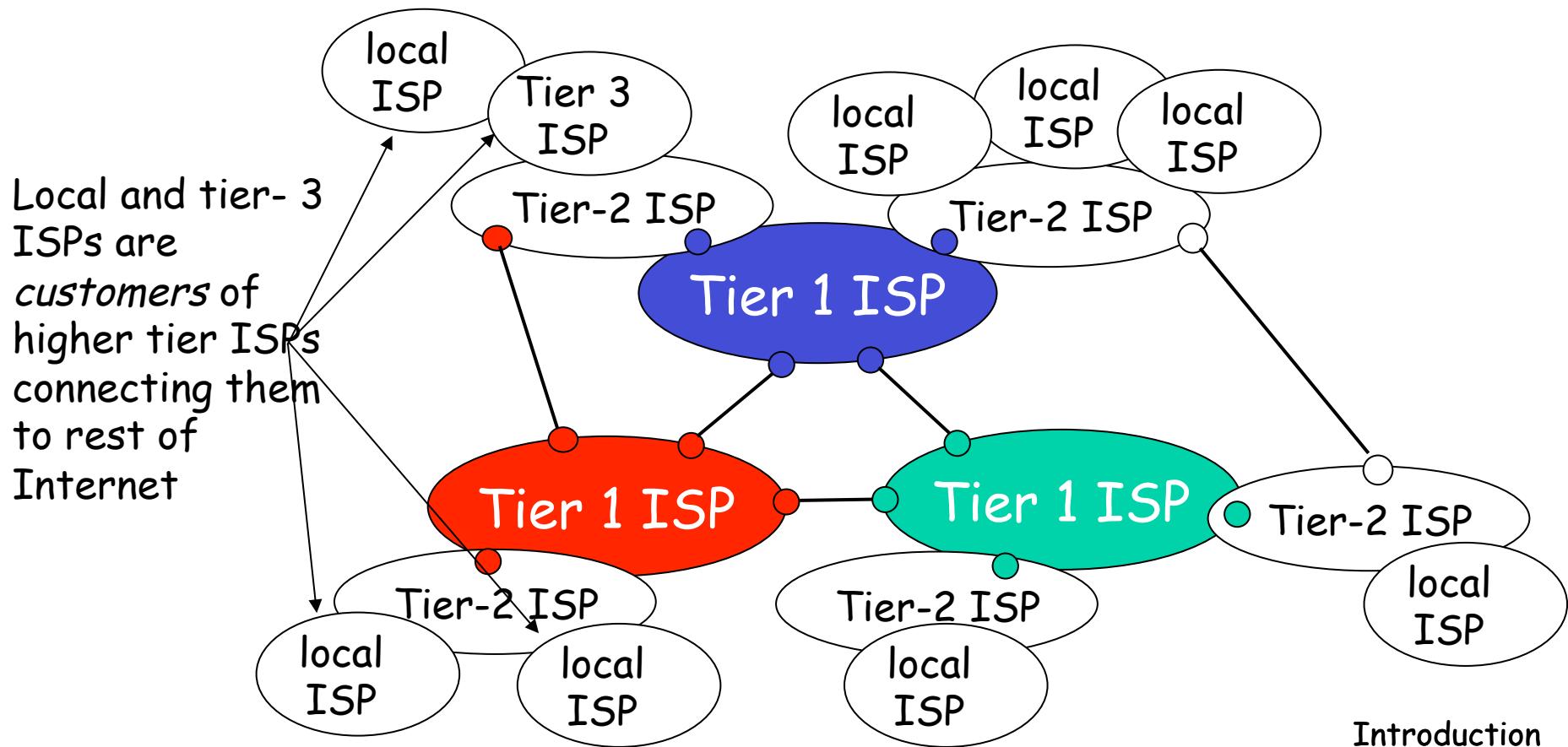
Internet structure: network of networks

- “Tier-2” ISPs: smaller (often regional) ISPs
 - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



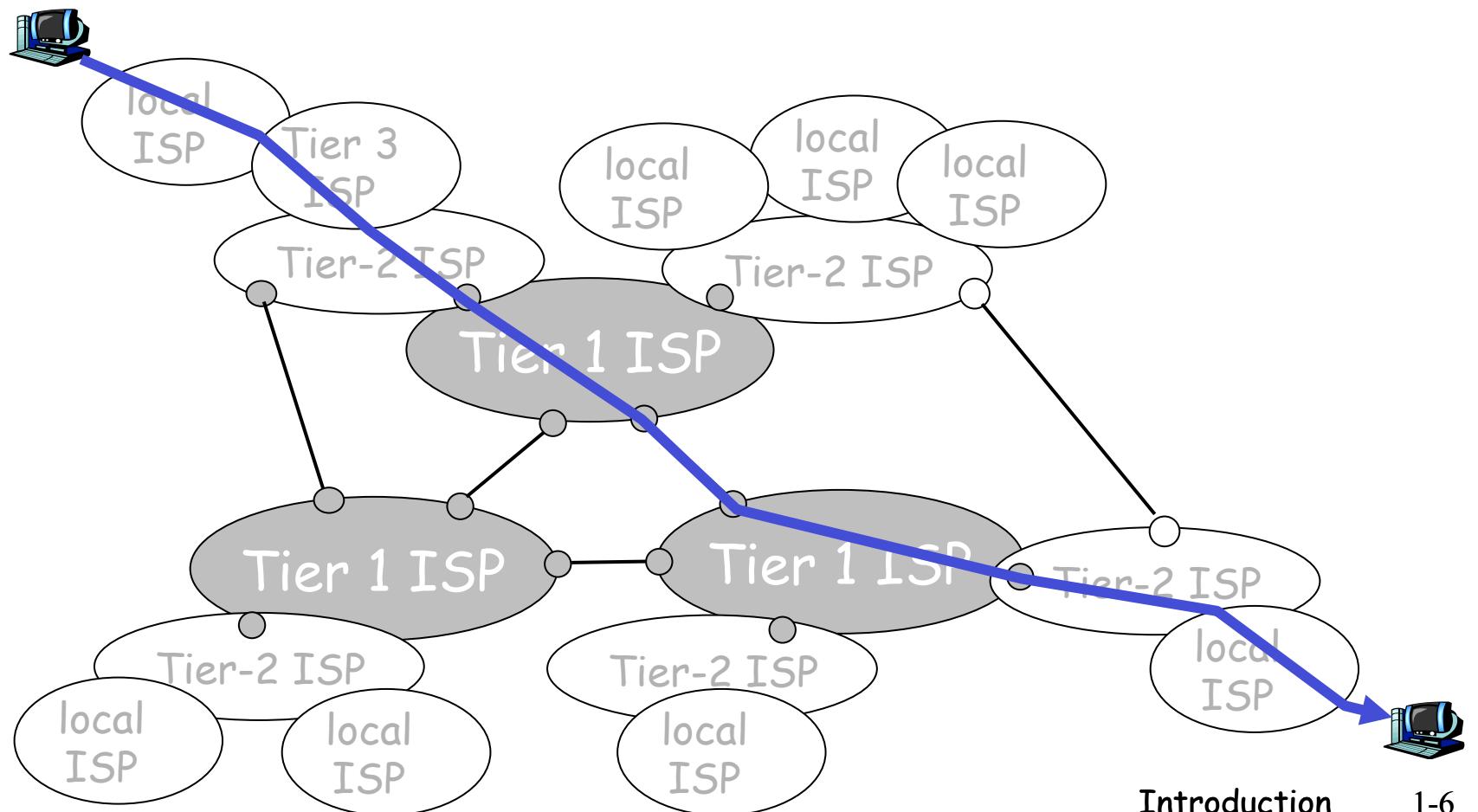
Internet structure: network of networks

- “Tier-3” ISPs and local ISPs
 - last hop (“access”) network (closest to end systems)



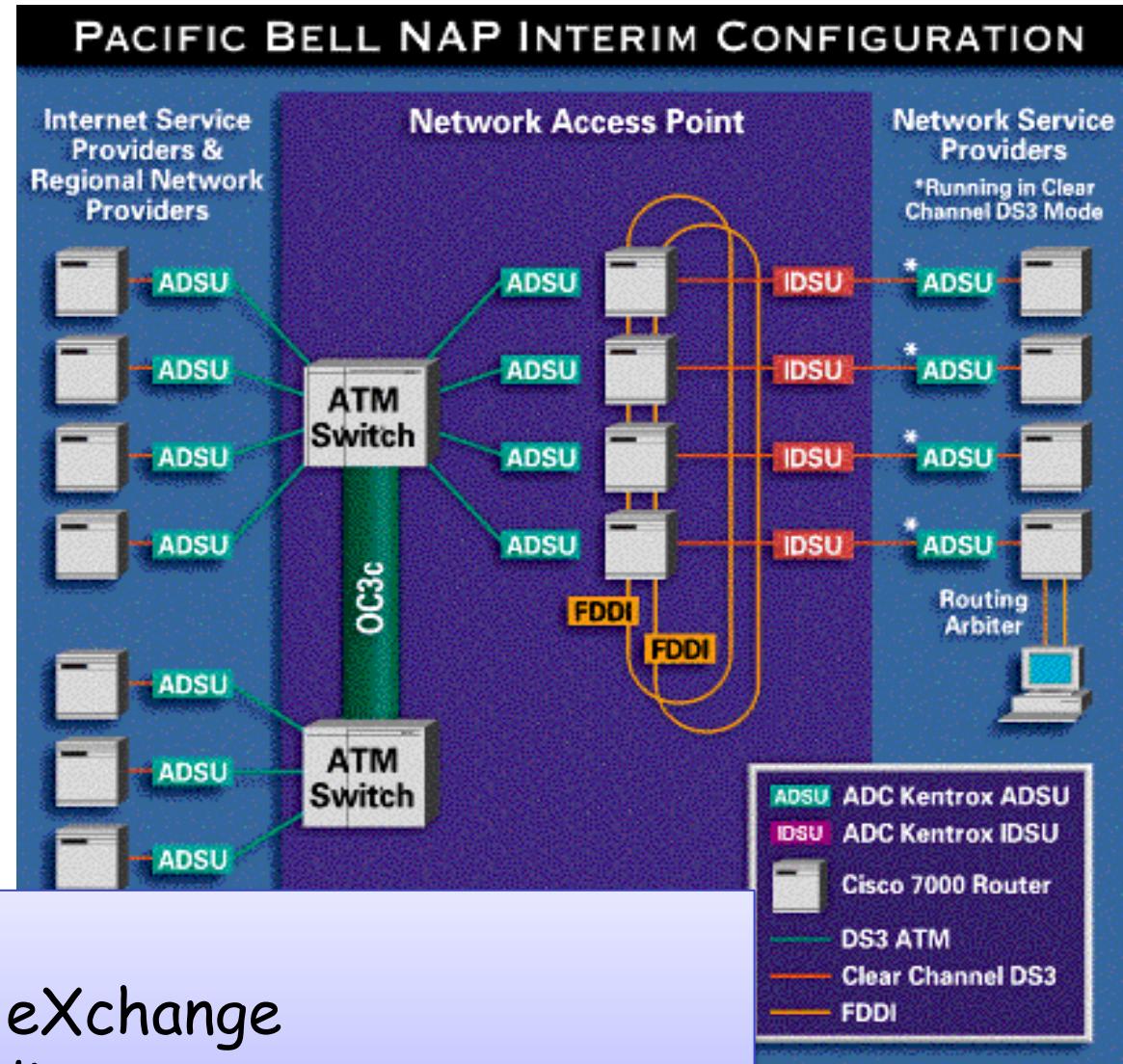
Internet structure: network of networks

- ❑ a packet passes through many networks!



A NAP: just another router...?

Pacific Bell
S.
Francisco
NAP



In Italia:
-MIX Milan Internet eXchange
-NaMeX Nautilus Mediterranean Exchange Point

Chapter 1: roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

Protocol “Layers”

Networks are complex!

□ many “pieces”:

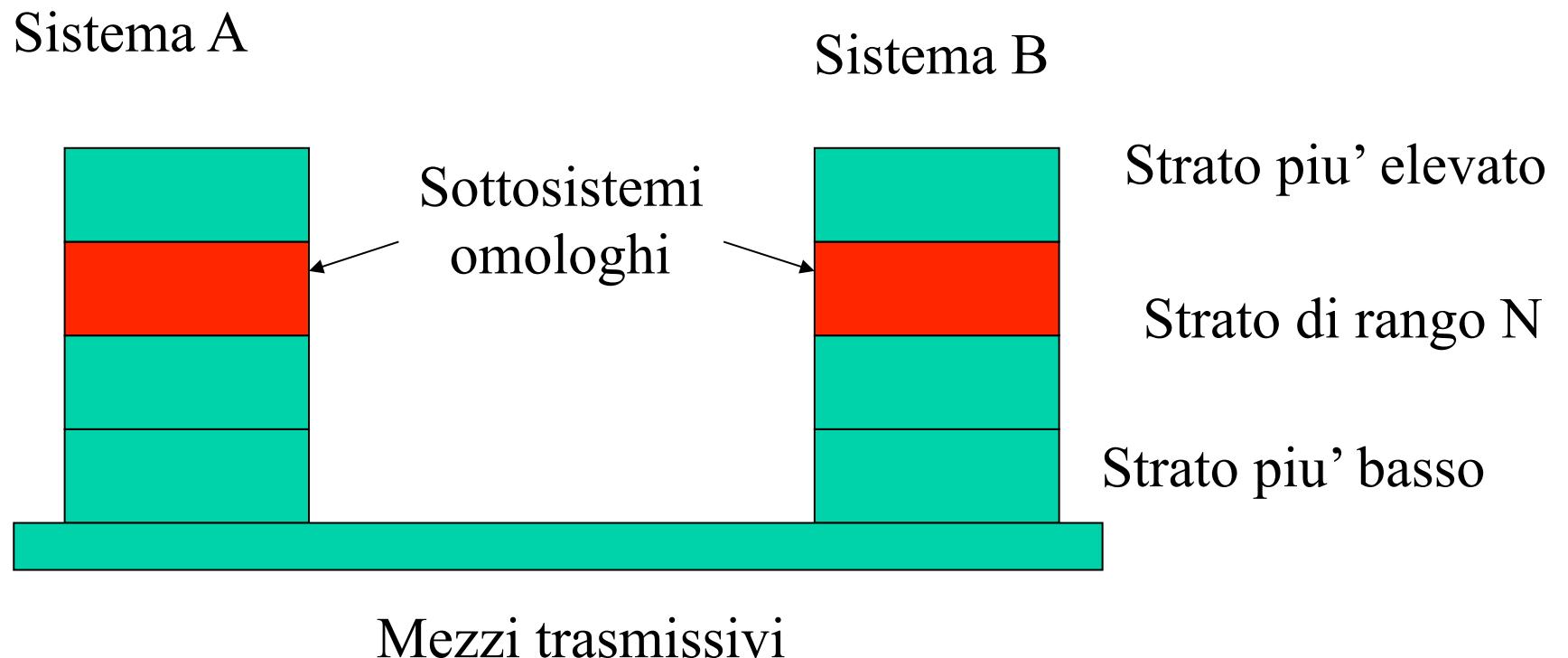
- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question:

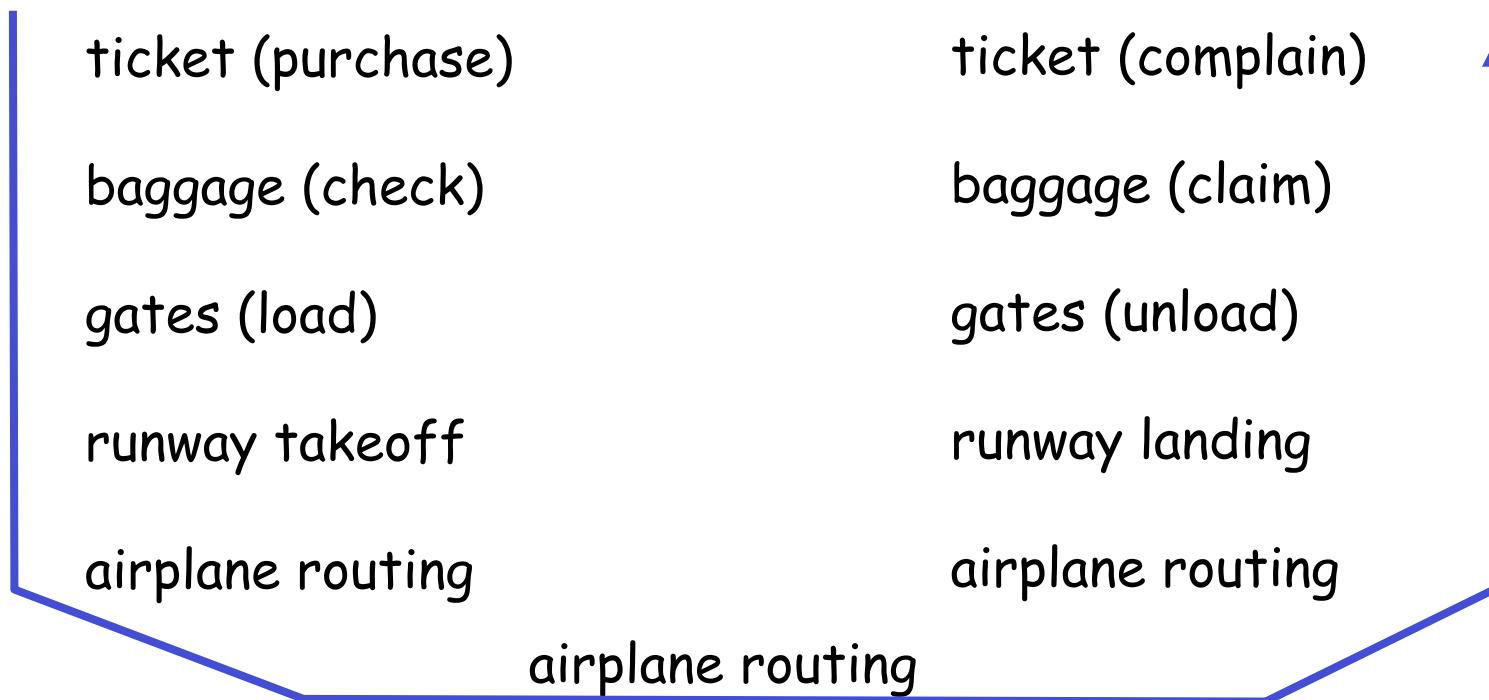
Is there any hope of organizing structure of network?

Or at least our discussion of networks?

Layering



Organization of air travel



- ❑ a series of steps

Organization of air travel: a different view

ticket (purchase)	ticket (complain)
baggage (check)	baggage (claim)
gates (load)	gates (unload)
runway takeoff	runway landing
airplane routing	airplane routing
	airplane routing

Layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Layered air travel: services

Counter-to-counter delivery of person+bags

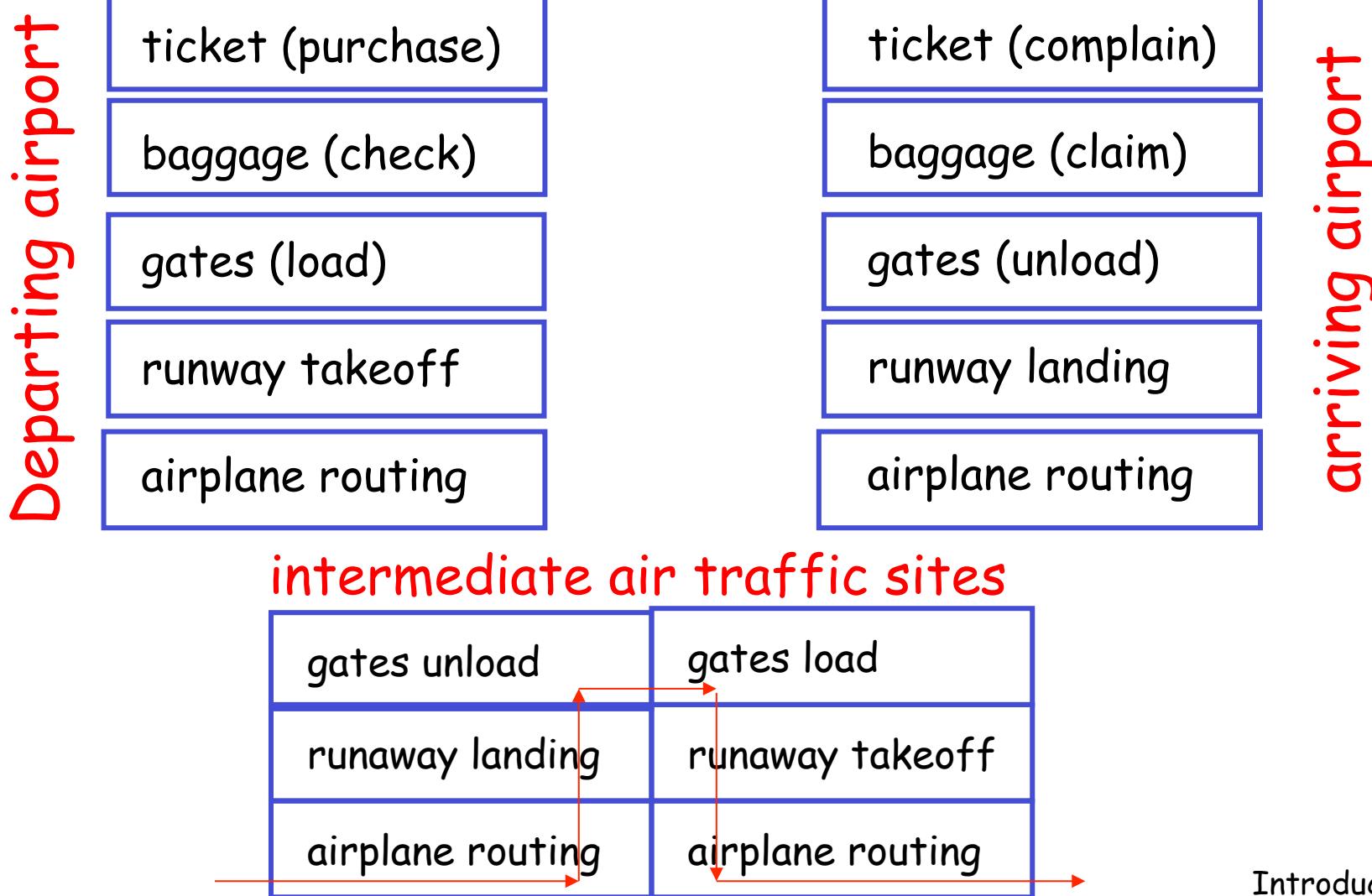
baggage-claim-to-baggage-claim delivery

people transfer: loading gate to arrival gate

runway-to-runway delivery of plane

airplane routing from source to destination

Distributed implementation of layer functionality



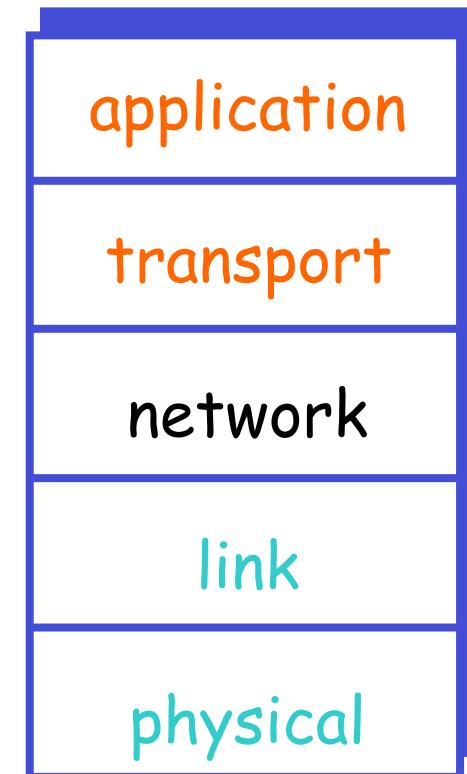
Why layering?

Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered **reference model** for discussion
- modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system (I.e. if baggage check and claim procedures changed due to Sept 11th or if the boarding rules change, boarding people by age)
- layering considered harmful?

Internet protocol stack

- **application:** supporting network applications
 - FTP, SMTP, HTTP
- **transport:** host-host data transfer
 - TCP, UDP
- **network:** routing of datagrams from source to destination
 - IP, routing protocols
- **link:** data transfer between neighboring network elements
 - PPP, Ethernet
- **physical:** bits “on the wire”



Typically in HW

Typically SW

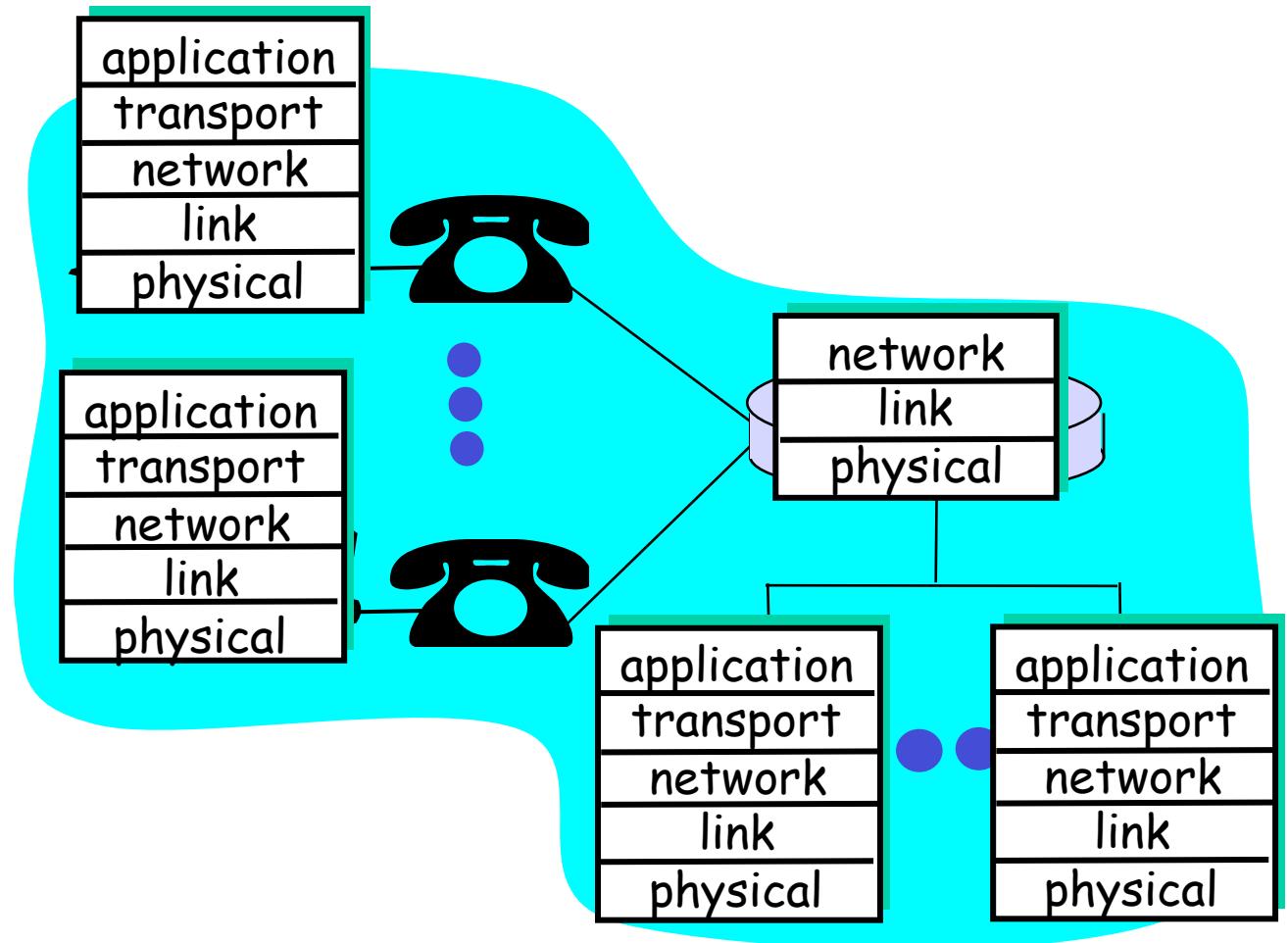
Introduction

1-16

Layering: logical communication

Each layer:

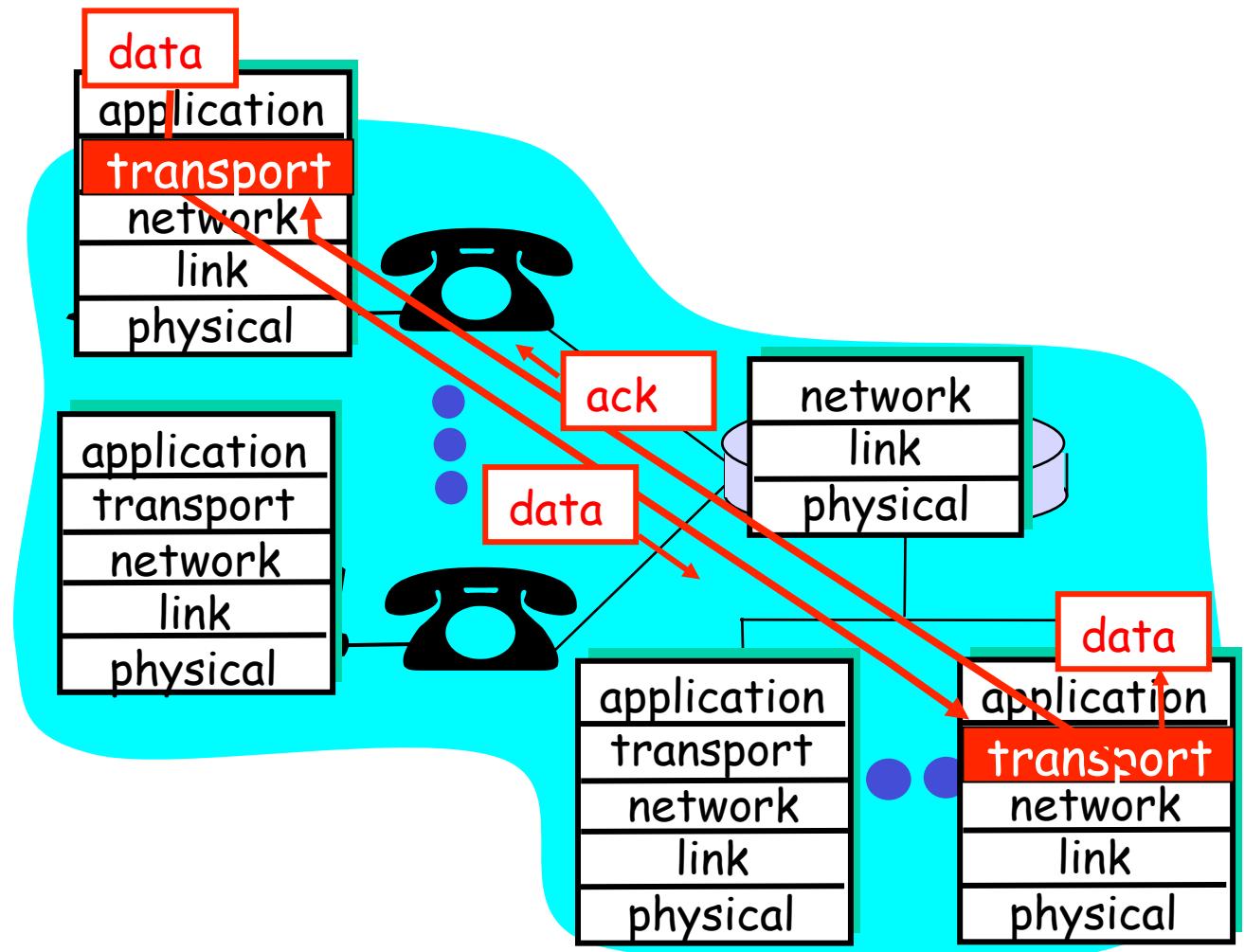
- distributed
- “entities” implement layer functions at each node
- entities perform actions, exchange messages with peers



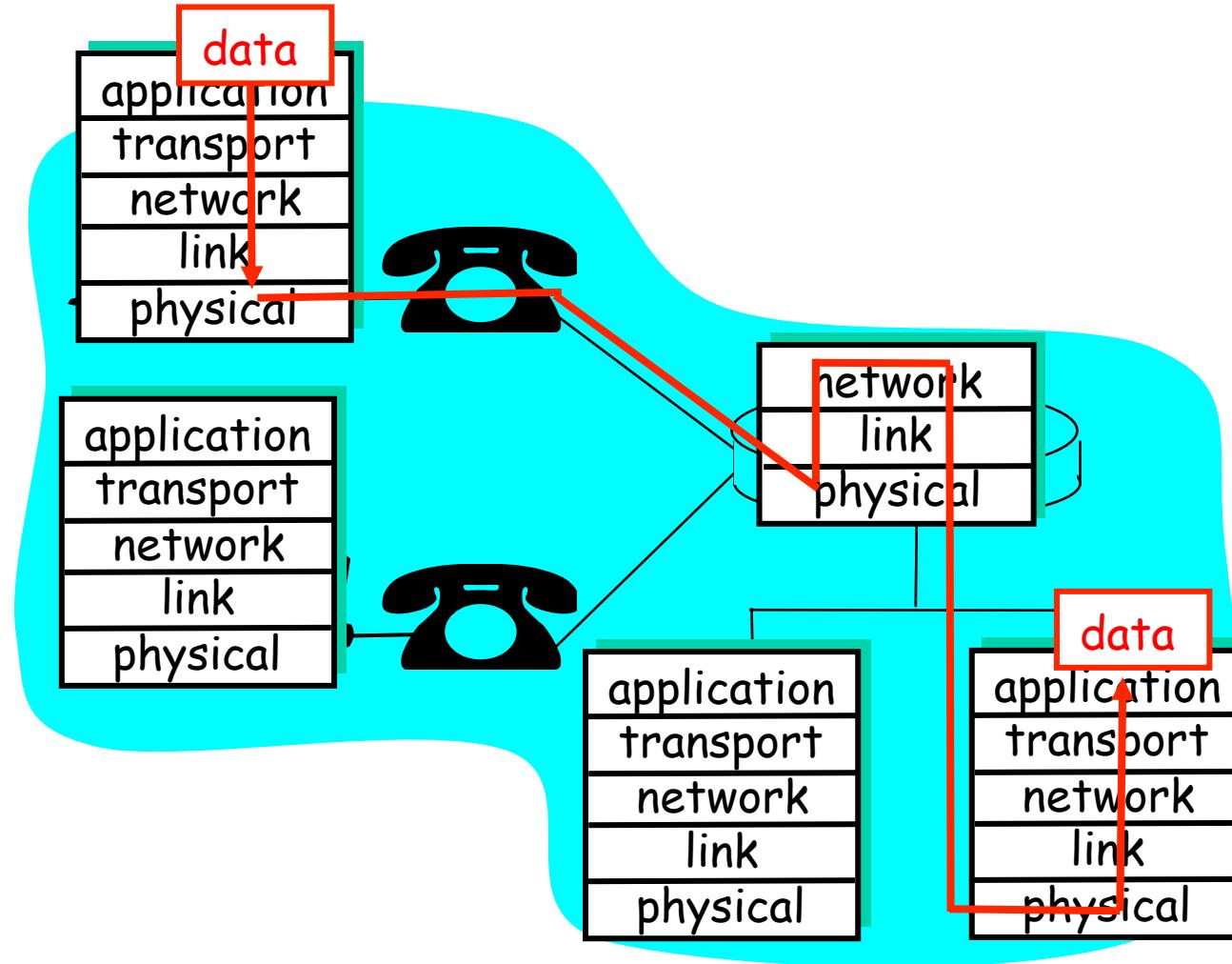
Layering: logical communication

E.g.: transport

- take data from app
- add addressing, reliability check info to form “datagram”
- send datagram to peer
- wait for peer to ack receipt
- analogy: post office



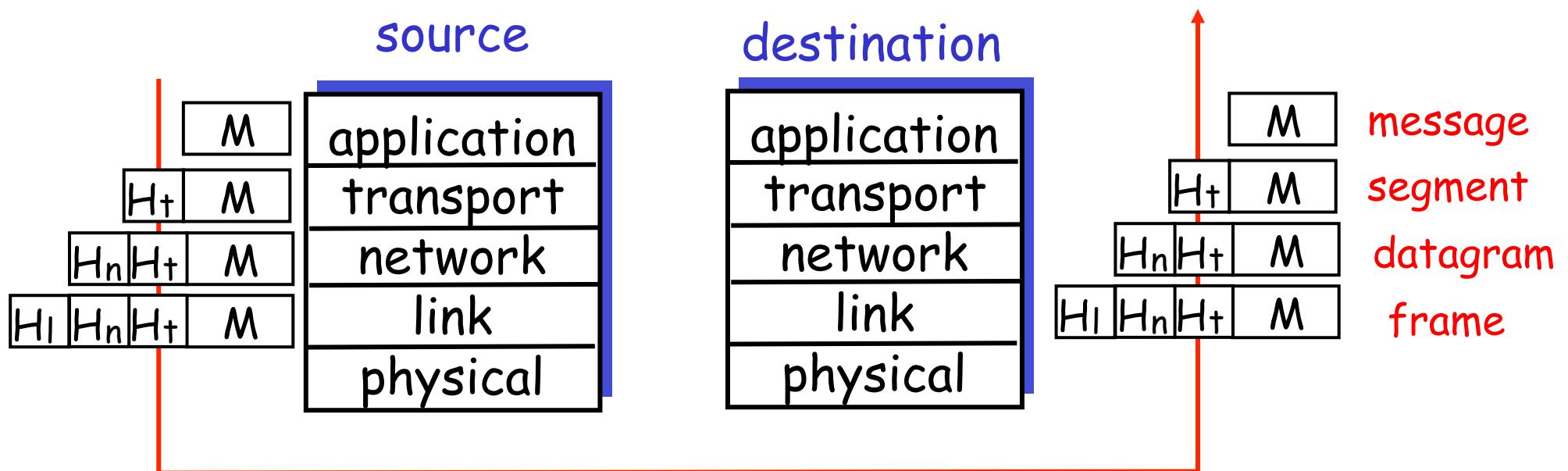
Layering: physical communication



Protocol layering and data

Each layer takes data from above

- adds header information to create new data unit
- passes new data unit to layer below



Layering: pros

□ Vantaggi della stratificazione

○ Modularita'

- Semplicita' di design
- Possibilita' di modificare un modulo in modo trasparente se le interfacce con gli altri livelli rimangono le stesse
- Possibilita' per ciascun costruttore di adottare la propria implementazione di un livello purché requisiti su interfacce soddisfatti

○ Gestione dell'eterogeneita'

- Possibili moduli 'diversi' per realizzare lo stesso insieme di funzioni, che riflettano l'eterogeneita' dei sistemi coinvolti (e.g. diverse tecnologie trasmissive, LAN, collegamenti punto-punto, ATM etc.)
- Moduli distinti possibili/necessari anche se le reti adottassero tutte la stessa tecnologia di rete perché ad esempio le applicazioni possono avere requisiti diversi (es. UDP e TCP). All'inizio TCP ed IP erano integrati. Perche' adesso sono su due livelli distinti?

Layering: cons

- Svantaggi della stratificazione
 - A volte modularita' inficia efficienza
 - A volte necessario scambio di informazioni tra livelli non adiacenti non rispettando principio della stratificazione

Internet History

1961-1972: Early packet-switching principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching (MIT)
- 1964: Baran - packet-switching in military nets
- Davies at the National Physical Laboratory, UK was also developing ideas on packet switching
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

Packet switches dubbed Interface Message Processors (IMP)

- 1972:
 - ARPAnet demonstrated publicly by Robert Kahn
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



- Leonard Kleinrock with first IMP

Kleinrock's students:
Vinton Cerf
John Postel...

Network measurement center UCLA

Internet History

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii (Abramson)
- 1973: Metcalfe's PhD thesis proposes Ethernet
- 1974: Cerf and Kahn - architecture for interconnecting networks
- late 70's: proprietary architectures, e.g. IBM SNA (Schwartz)
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

Internet History

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: SMTP e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: FTP protocol defined
- 1988: TCP congestion control
- new national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

Internet History

1990, 2000's: commercialization, the Web, new apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web

Late 1990's - 2000's:

- more killer apps: instant messaging, peer2peer file sharing (e.g., Napster)
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

Significant late developments: P2P, broadband access, wireless Internet

Recent trends (2000-2016)

- Intense evolution
- Aggressive deployment of broadband Internet access to homes
 - enabler of distribution of user generated videos, on demand streaming videos, multi-person video conferencing services
- Ubiquitous deployment of high speed wireless access
 - number of wireless devices connected to Internet > wired devices from 2011
- Development of social networks
- Companies such as Google and Microsoft have developed extensive private networks
- Internet commerce companies and institutions run their applications on the cloud

A short digression:

where is Internet standardized?

Who controls the Internet?

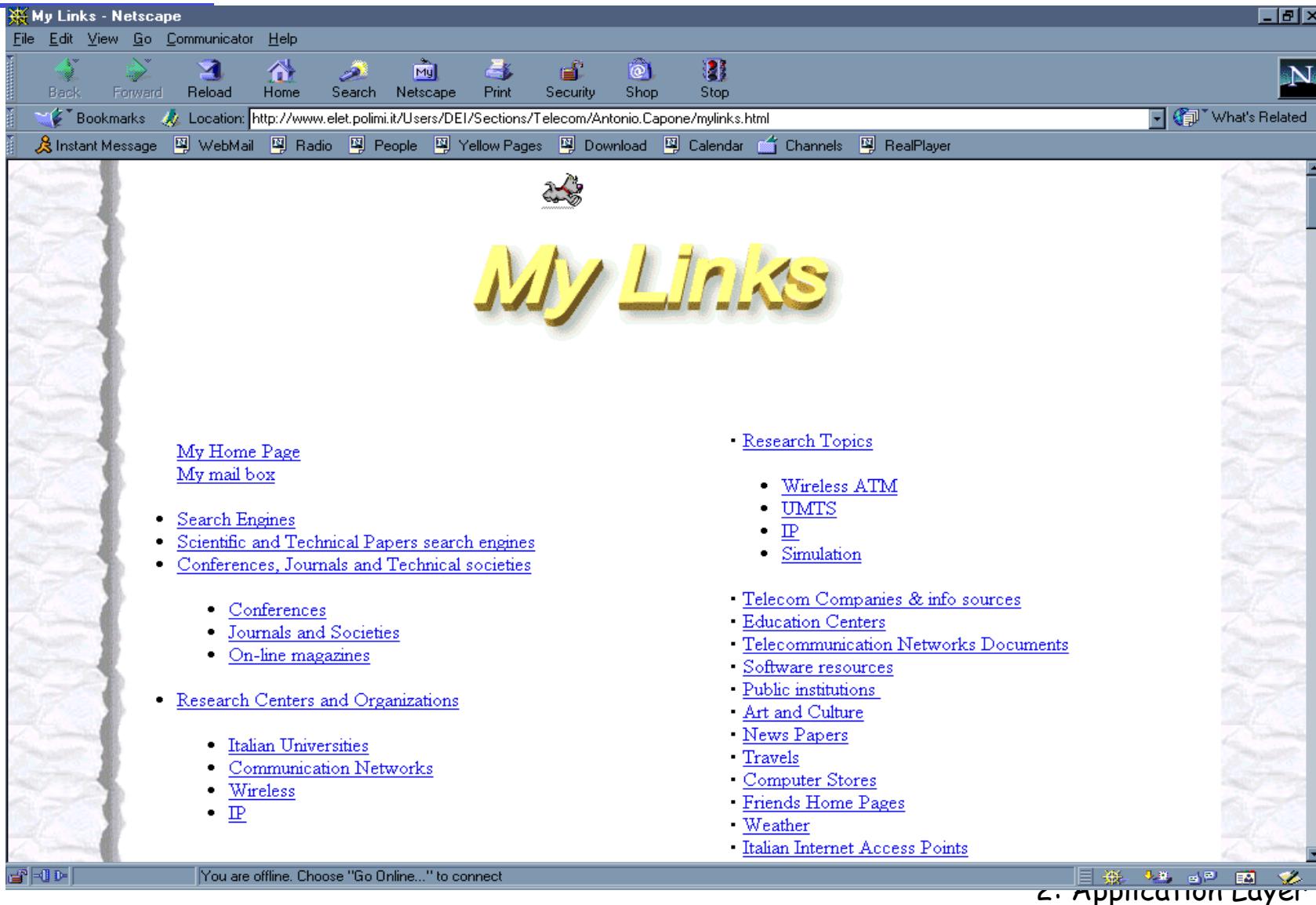
- No single administrative organization
- IETF - Internet Engineering Task Force (since 86)
 - Developement of current protocols and specifications for standardization.
 - International community, open to everyone
 - Most of the work via mailing lists
 - Meets three times/year
 - organized in areas and working groups
 - Dynamically activated & deactivated on need
 - group coordination: IESG (Internet Engineering Steering Group). Area directors are members of the IESG. Responsible for the actions associated with entry into and movement along the Internet "standards track," including final approval of specifications as Internet Standards.
- Industry also preemptively determine standards

Chapter 2 outline

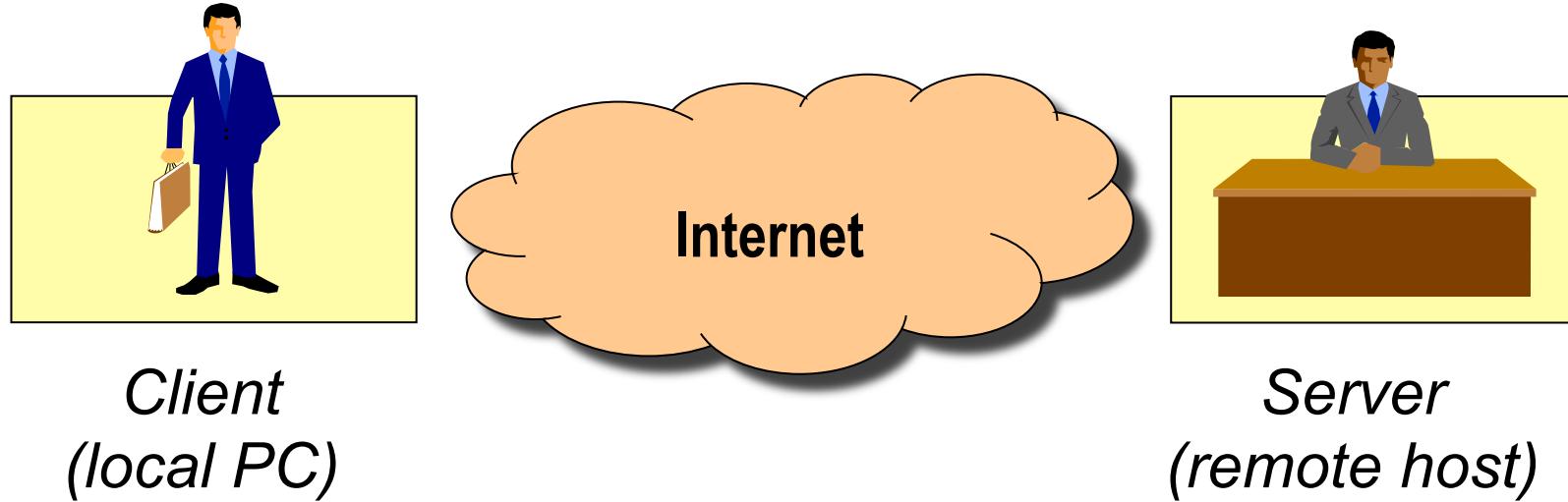
- 2.1 Principles of app layer protocols
 - clients and servers
 - app requirements
- 2.2 Web and HTTP
- 2.3 FTP
- 2.4 Electronic Mail
 - SMTP, POP3, IMAP
- 2.5 DNS
- 2.9 Content distribution
 - Network Web caching
 - Content distribution networks
 - P2P file sharing

I servizi di IP: esempio

il WEB



Basic scenario



Client wants to retrieve a web page. What happens?

What is a "page" on the web?

a resource (i.e a file), specified by a
URL: Uniform Resource Locator.

e.g. personal web pages:

`HTTP://cerbero.elet.polimi.it/people/bianchi/index.html`

The three components of an URL

1. Protocol (also called "scheme")

- *how can a page be accessed? (application protocol used)*
 - **http://cerbero.elet.polimi.it/people/bianchi/index.html**

2. Host name

- *Where is the page located? (symbolic or numeric location)*
 - **http://cerbero.elet.polimi.it/people/bianchi/index.html**

3. File (resource) name

- *What is the page called? (with full path)*
 - **http://cerbero.elet.polimi.it/people/bianchi/index.html**

Network applications: some jargon

Process: program running within a host.

- within same host, two processes communicate using **interprocess communication** (defined by OS).
- processes running in different hosts communicate with an **application-layer protocol** (defining message format, which message to exchange and in which order)

user agent: interfaces with user “above” and network “below”.

- implements user interface & application-level protocol
 - Web: browser
 - E-mail: mail reader
 - streaming audio/video: media player

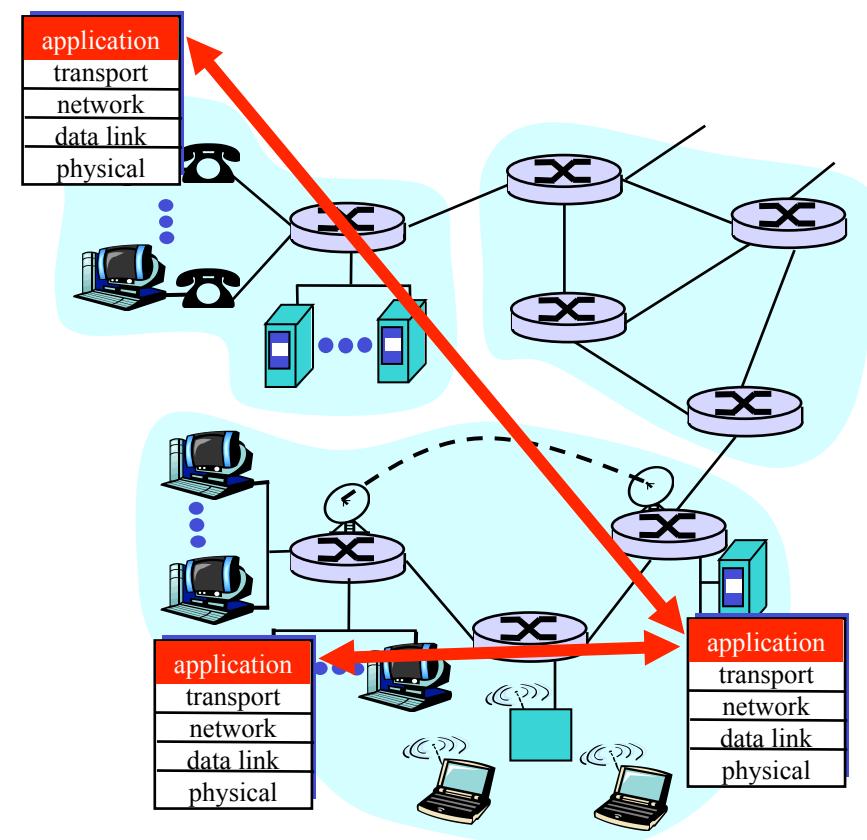
Applications and application-layer protocols

Application: communicating, distributed processes

- e.g., e-mail, Web, P2P file sharing, instant messaging
- running in end systems (hosts)
- exchange messages to implement application

Application-layer protocols

- one “piece” of an app
- define messages exchanged by apps and actions taken
- use communication services provided by lower layer protocols (TCP, UDP)



App-layer protocol defines

- Types of messages exchanged, eg, request & response messages
- Syntax of message types: what fields in messages & how fields are delineated
- Semantics of the fields, ie, meaning of information in fields
- Rules for when and how processes send & respond to messages

Public-domain protocols:

- defined in RFCs
- allows for interoperability
- eg, HTTP, SMTP

Proprietary protocols:

- eg, KaZaA

Client-server paradigm

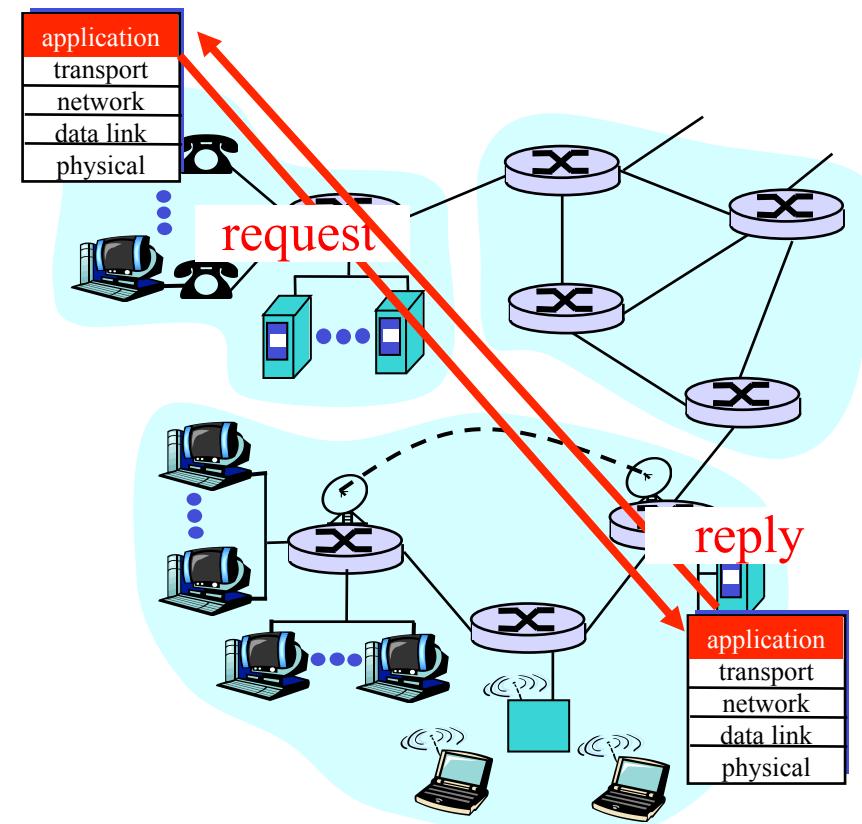
Typical network app has two pieces:
client and *server*
(NEXT SLIDE)

Client:

- initiates contact with server (“speaks first”)
- typically **requests service** from server,
- Web: client implemented in browser; e-mail: in mail reader

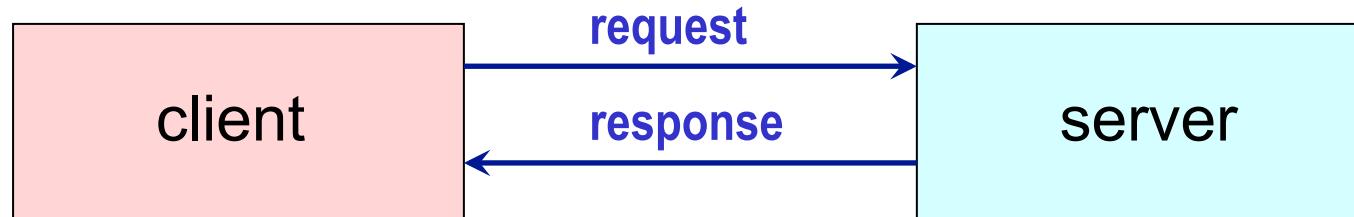
Server:

- provides requested service to client**
- e.g., Web server sends requested Web page, mail server delivers e-mail



Architettura Client-Server

- Un processo client è solo in grado fare richieste di servizio (informazioni) e di interpretare le risposte
- Un processo server ha solo il compito di interpretare le richieste e fornire le risposte
- Se è necessario nello stesso host sia fare richieste che fornire risposte vengono usati due processi, client e server.
- Un protocollo applicativo per un'architettura client-server rispecchia questa divisione di ruoli e prevede messaggi di richiesta (request) generati dal lato client e messaggi di risposta (response) generati dal server



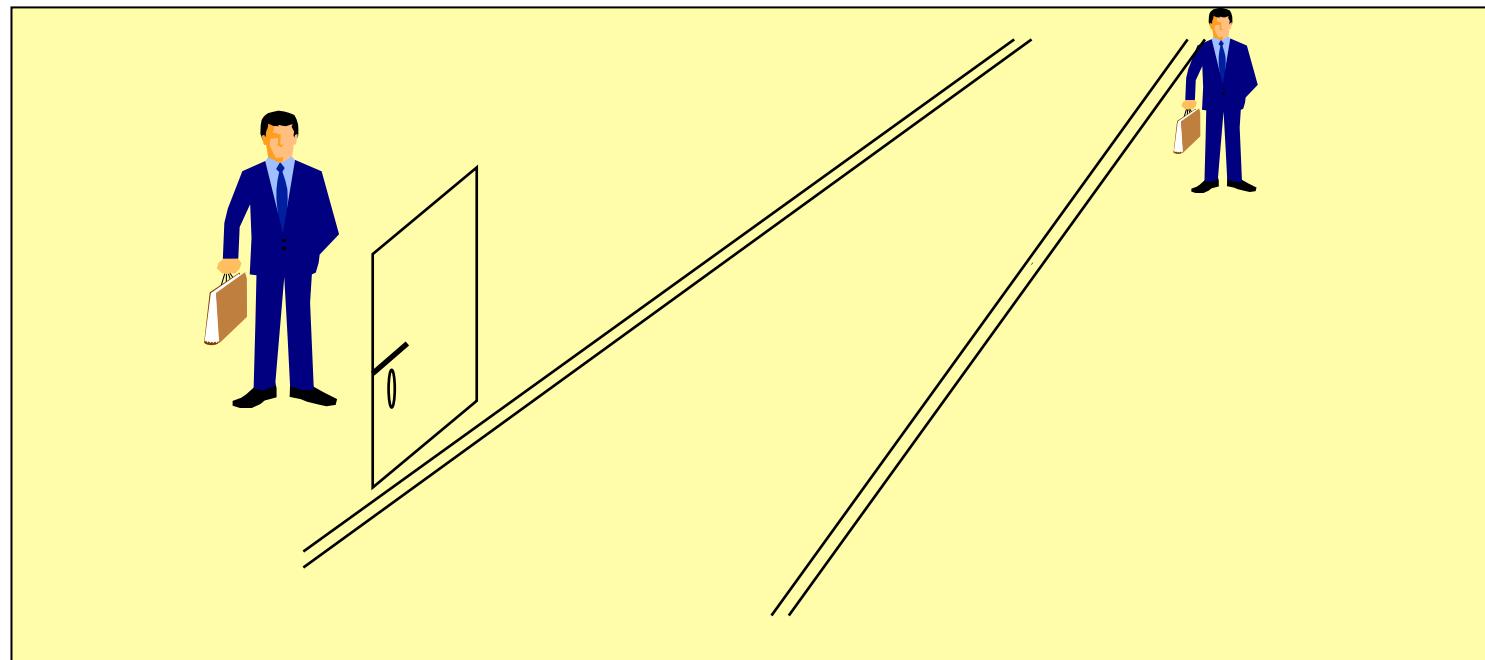
Programmi Client e Server

- Distinguiamo tra programma (software) e processo, istanza del programma in esecuzione su un host
- Un processo server è in esecuzione a tempo illimitato sul proprio host (daemon) e viene attivato mediante una *passive open*
- Un processo client viene attivato solo al momento di fare le richieste e viene attivato mediante una *active open* su richiesta dell'utente o di altro processo applicativo
- la *passive open* del server fa sì che da quel momento il server accetti richieste dai client
- l'*active open* del client richiede l'indicazione dell'indirizzo e della porta del client, ed inizia la connessione TCP per mettere in comunicazione il client e il server

Addressing processes:

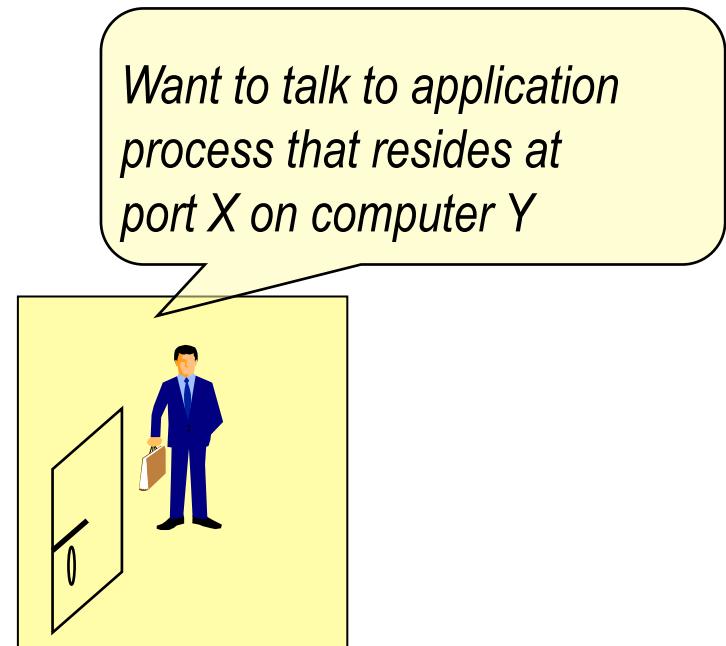
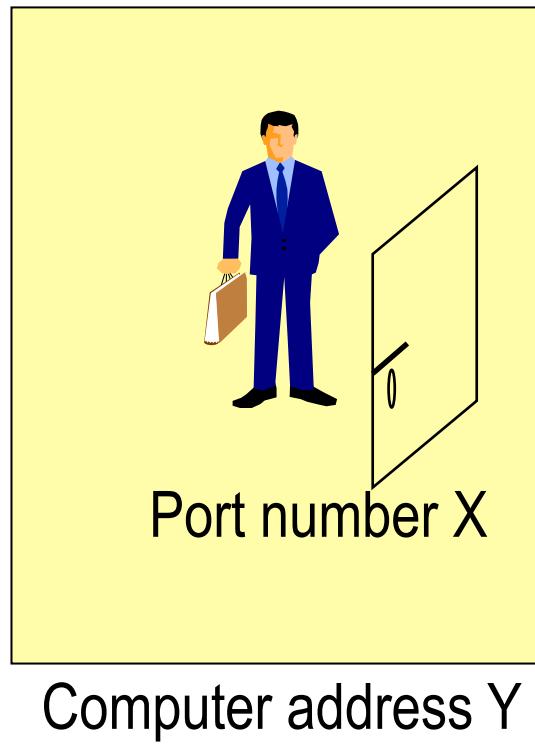
- For a process to receive messages, it must have an identifier
- Every host has a unique 32-bit IP address
- Q: does the IP address of the host on which the process runs suffice for identifying the process?
- Answer: No, many processes can be running on same host
- Identifier includes both the IP address and **port numbers** associated with the process on the host.
- Example port numbers:
 - HTTP server: 80
 - Mail server: 25

Port numbers



the “address” of the SW process inside the computer!

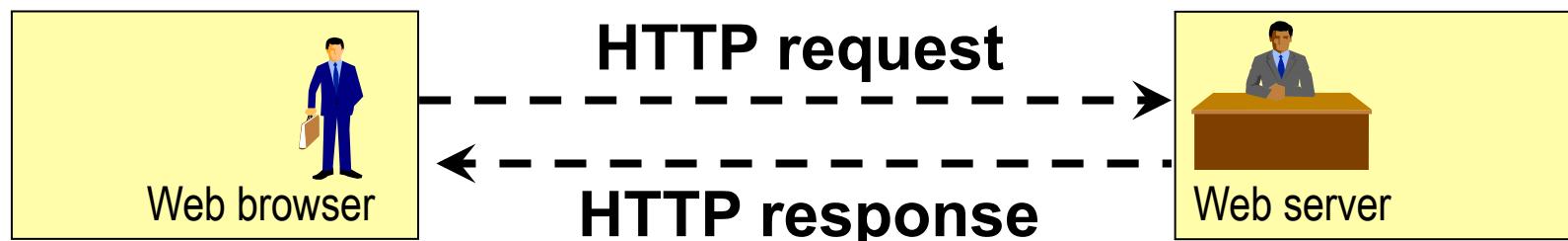
Same addressing scheme works for different machines!!



Re-understanding URLs

HTTP://cerbero.elet.polimi.it/people/bianchi/index.html

protocol location Filename



Web browser (SW process) needs to send HTTP Request to Web Server.
Location not enough (it is just the computer address!)

Addressing web servers: (wrong) idea

HTTP://cerbero.elet.polimi.it/....

The diagram shows the URL "HTTP://cerbero.elet.polimi.it/...." with three curly braces underneath it. The first brace covers "HTTP:", labeled "protocol". The second brace covers "cerbero.elet.polimi.it", labeled "location". The third brace covers the trailing slashes and dots, labeled "Filename".

protocol location Filename

Location = computer address!

Protocol = SW process address (HTTP = goto Web Server)

.... **No need for port numbers??**

...

What if more than one Web Server installed???

URL structure (corrected!)

HTTP://cerbero.elet.polimi.it :*portnumber*/....

The URL is divided into three main parts by curly braces: 'protocol' covers 'HTTP:', 'location' covers the host name 'cerbero.elet.polimi.it', and 'Filename' covers the trailing slash and any path components.

Default value for HTTP: 80

HTTP://cerbero.elet.polimi.it/...

equal to

HTTP://cerbero.elet.polimi.it:80/...

different from

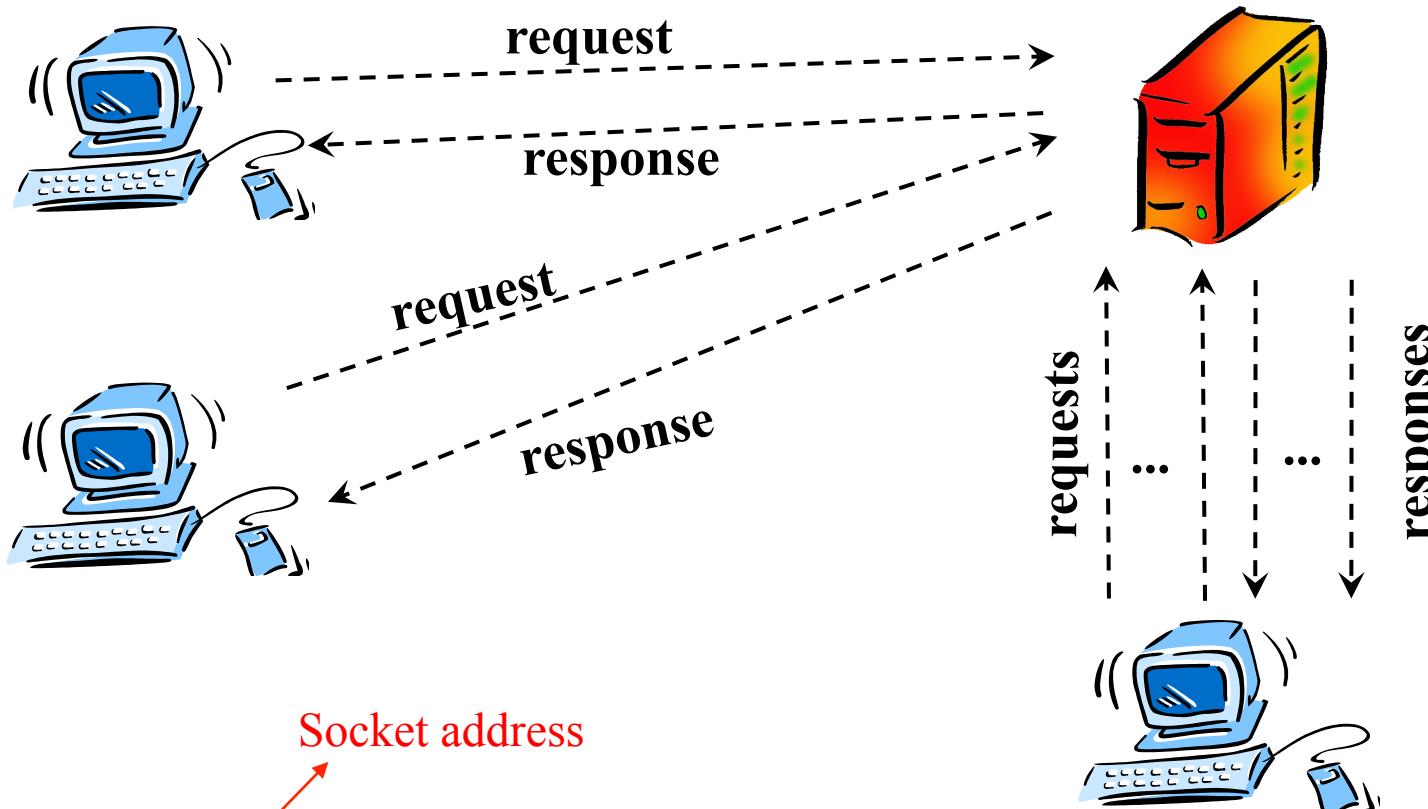
HTTP://cerbero.elet.polimi.it:8080/... (if exists!)

Port numbers

- 16 bit address (0-65535)
- well known port numbers for common servers
 - FTP 20, TELNET 23, SMTP 25, HTTP 80, POP3 110,
... (full list: RFC 1700)
- number assignment (by IANA)
 - 0 not used
 - 1-255 reserved for well known processes
 - 256-1023 reserved for other processes
 - 1024-65535 dedicated to user apps

Programmi Client e Server

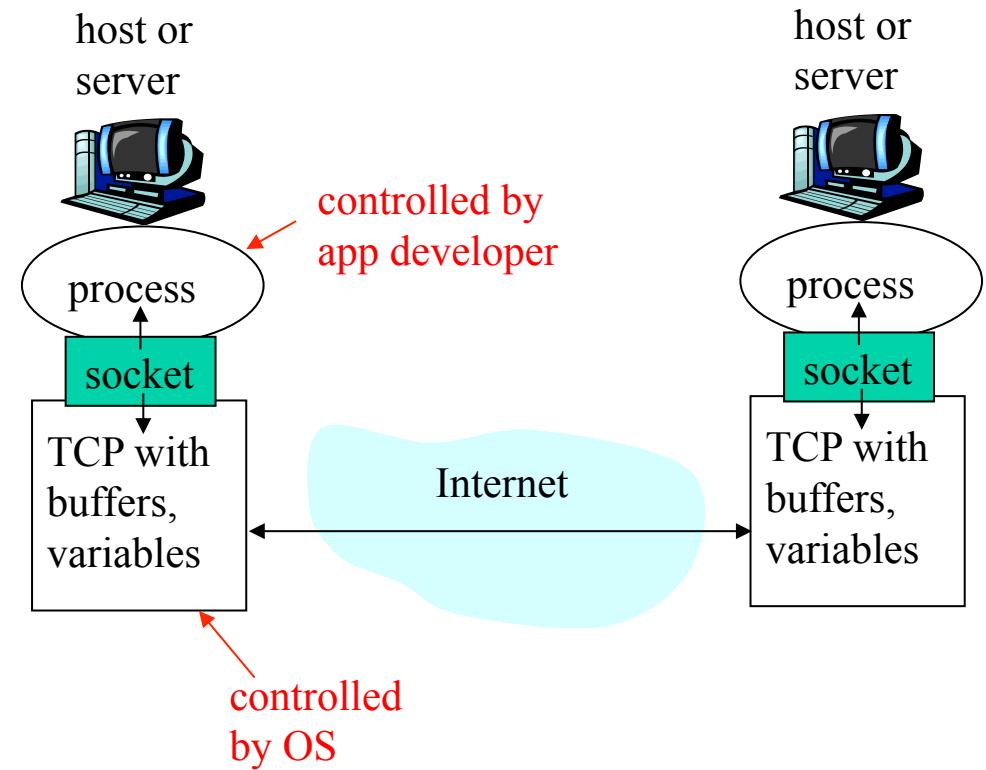
- Normalmente più client possono inviare richieste ad uno stesso server
- Un client può fare più richieste contemporanee



- Un flusso di informazioni tra due processi identificato da una quadrupla (*indirizzo IP sorgente, porta sorgente, indirizzo IP destinazione, porta destinazione*)
- Di solito il client NON USA un well known port #
 - OS assegna un num. di porta disponibile

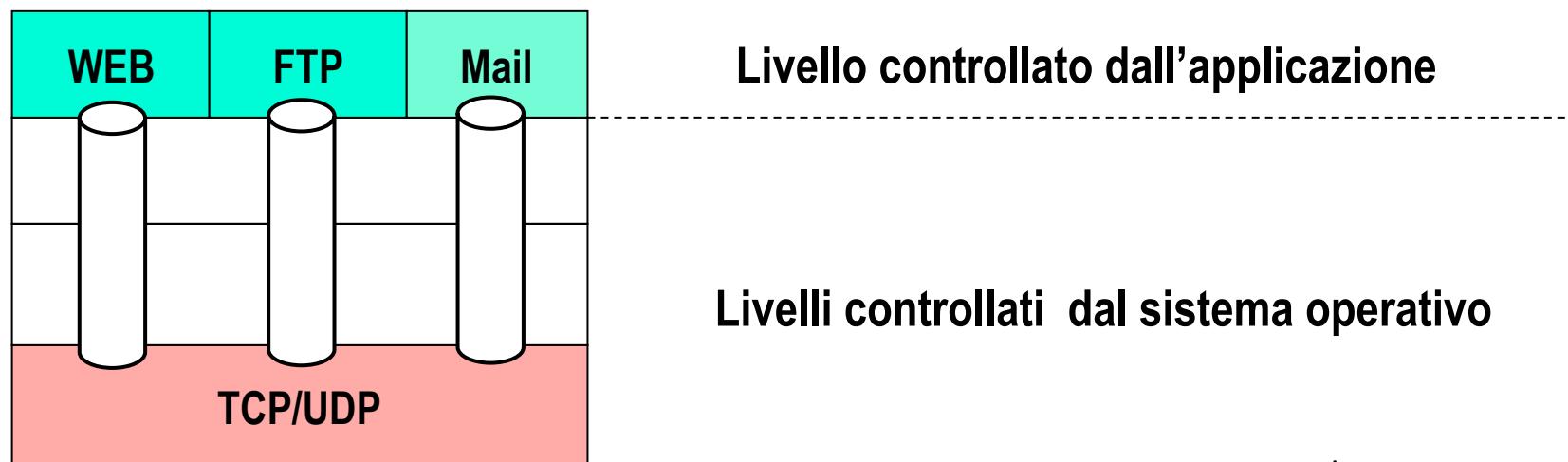
Processes communicating across network

- process sends/receives messages to/from its socket
- socket analogous to door
 - sending process shoves message out door
 - sending process assumes transport infrastructure on other side of door which brings message to socket at receiving process
- API: (1) choice of transport protocol; (2) ability to fix a few parameters (lots more on this later)



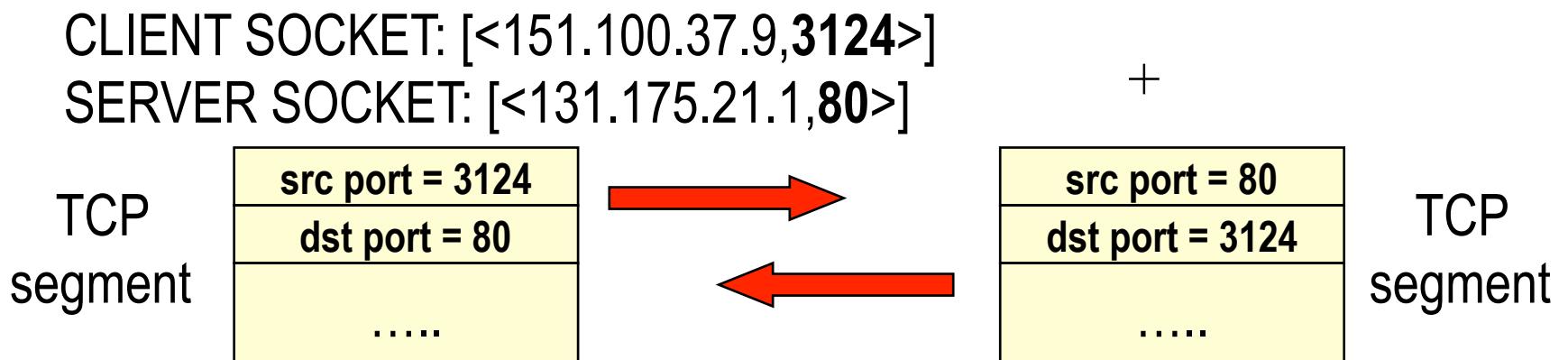
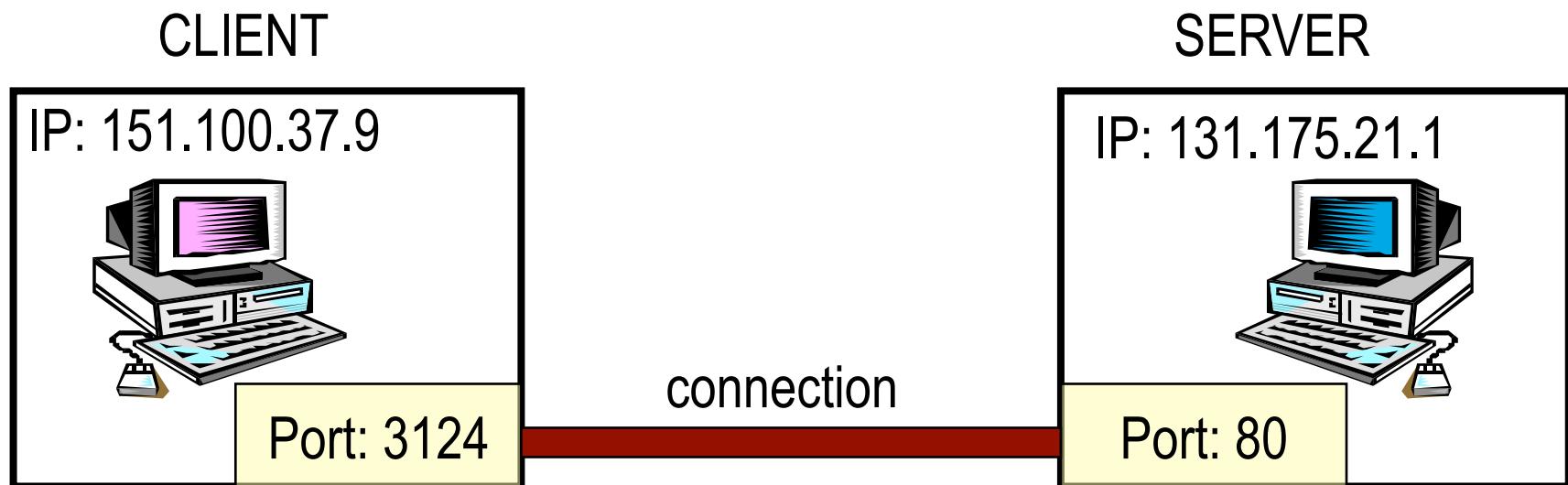
Interazione coi livelli inferiori: architettura TCP/IP

- I protocolli applicativi che si appoggiano sull' Internet Protocol si appoggiano direttamente sul protocollo di trasporto TCP/UDP
- Lo scambio di messaggi fra i processi applicativi avviene utilizzando i servizi dei livelli inferiori attraverso i SAP (Service Access Point).
- Qualsiasi protocollo di livello applicativo accede al servizio di traporto mediante socket.
- I protocolli di trasporto estendono il servizio di trasporto in rete dell'info tra due end systems offerto da IP consentendo il trasporto dell'info tra due PROCESSI APPLICATIVI in esecuzione sui due end systems
- Come si distingue a quale protocollo di trasporto passare il segmento? Protocol number

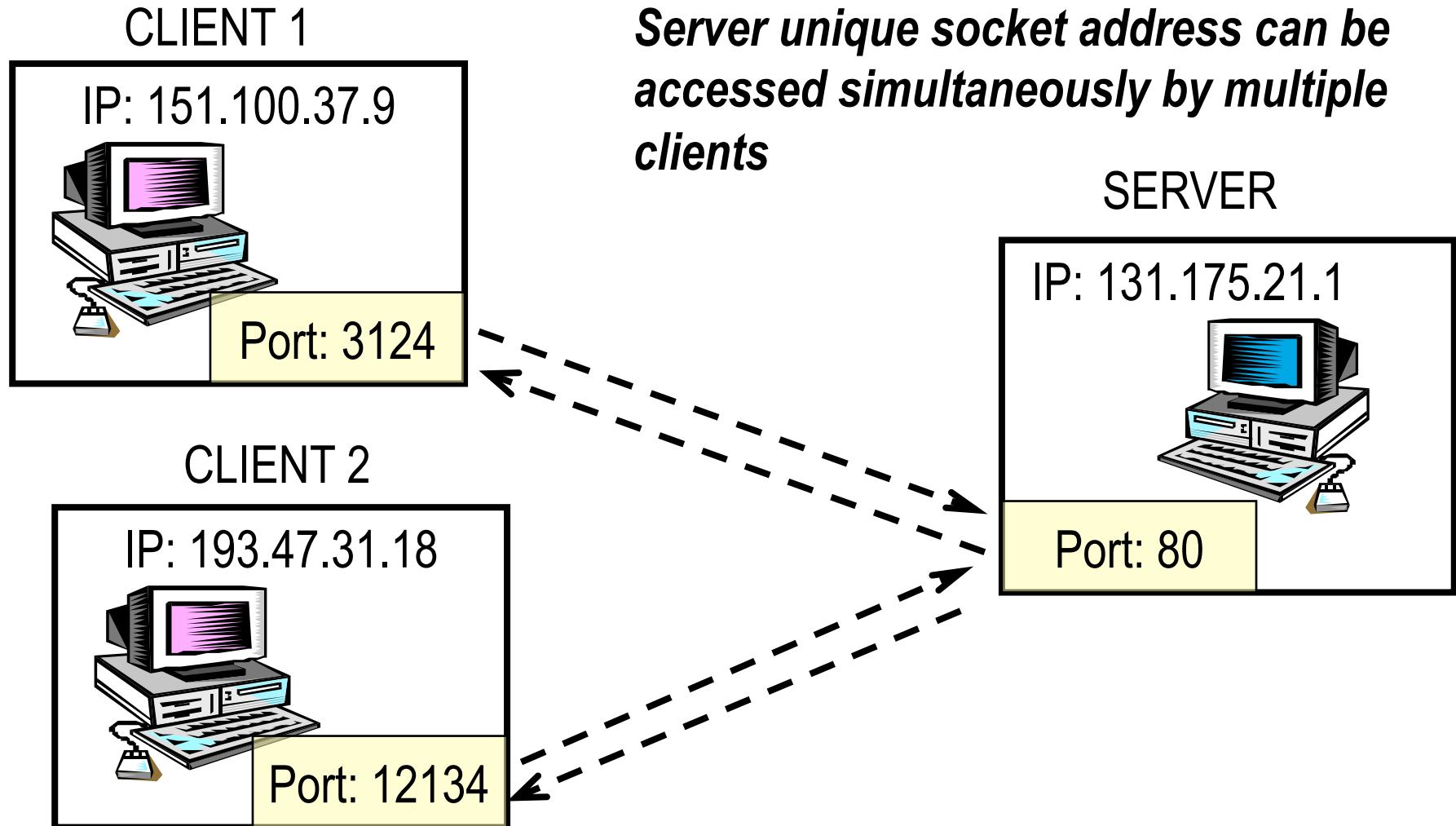


Connections

identified by socket addresses at its ends

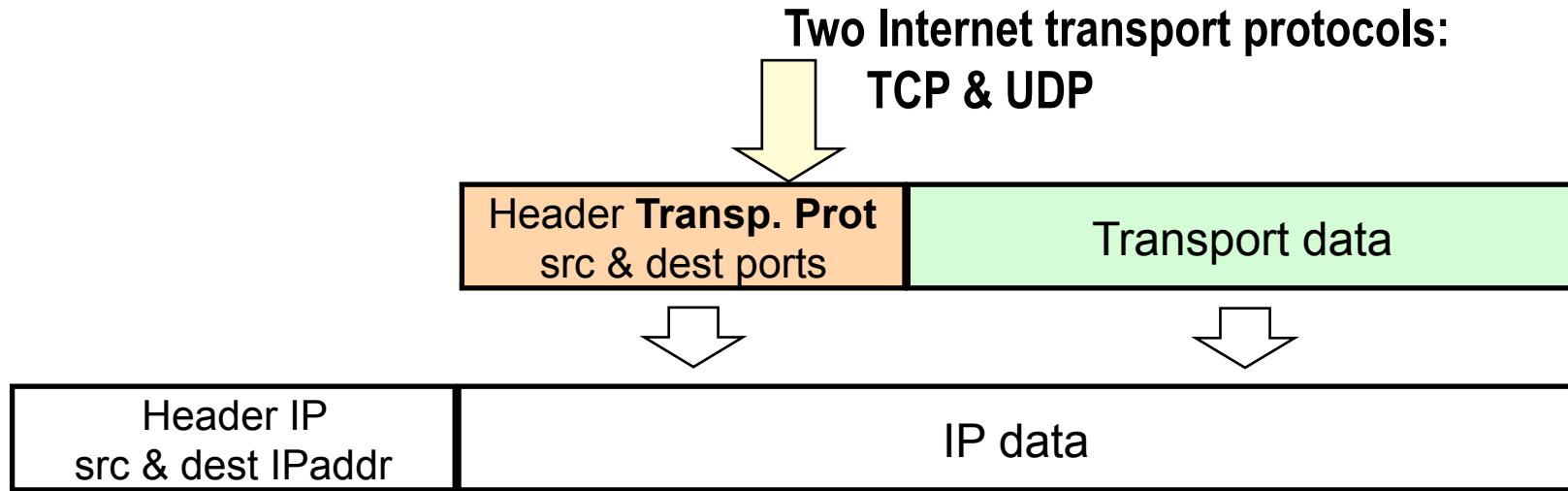


Managing multiple connections



Socket addresses: where?

- ❑ Ports: in the Internet Transport protocol header (TCP or UDP)
- ❑ IP addresses: In the IP packet header



IP packets travel in the network based on IP address;

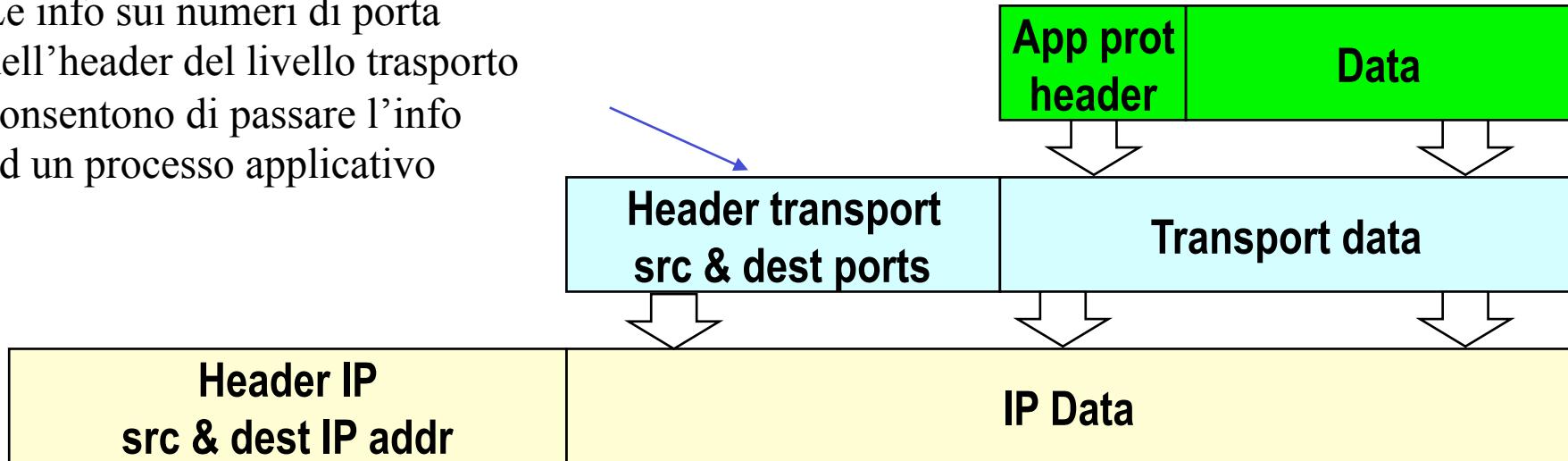
Once they reach destination, they are delivered to app based on port #

Sockets and Internet transport protocols

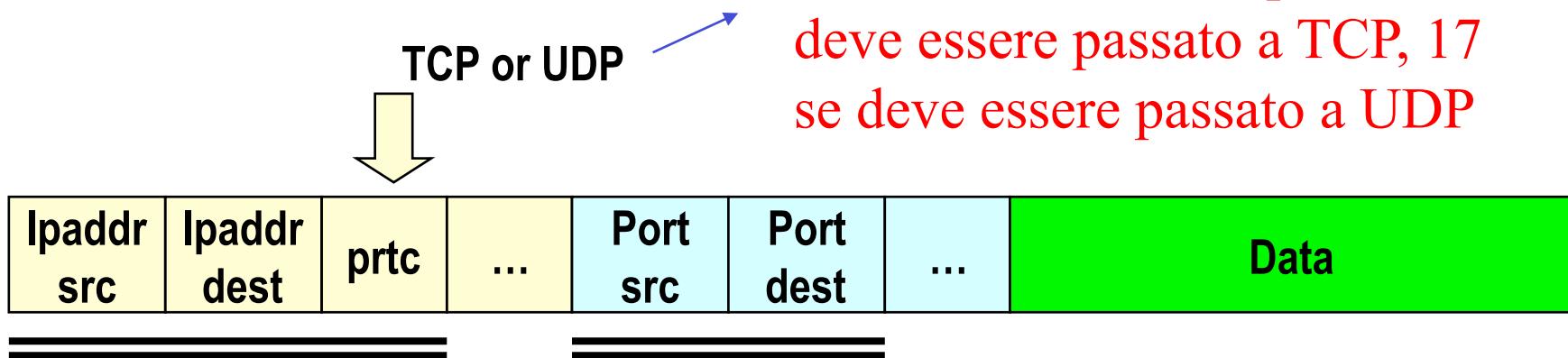
- Two transport protocols in Internet:
 - TCP (reliable, connection-oriented)
 - UDP (unreliable, connectionless)
- When opening socket, needs to specify which transport to use!

Socket addresses (refined)

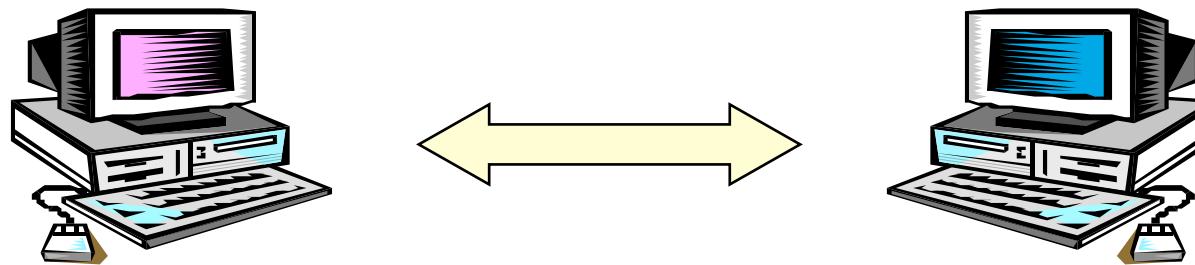
Le info sui numeri di porta
dell'header del livello trasporto
consentono di passare l'info
ad un processo applicativo



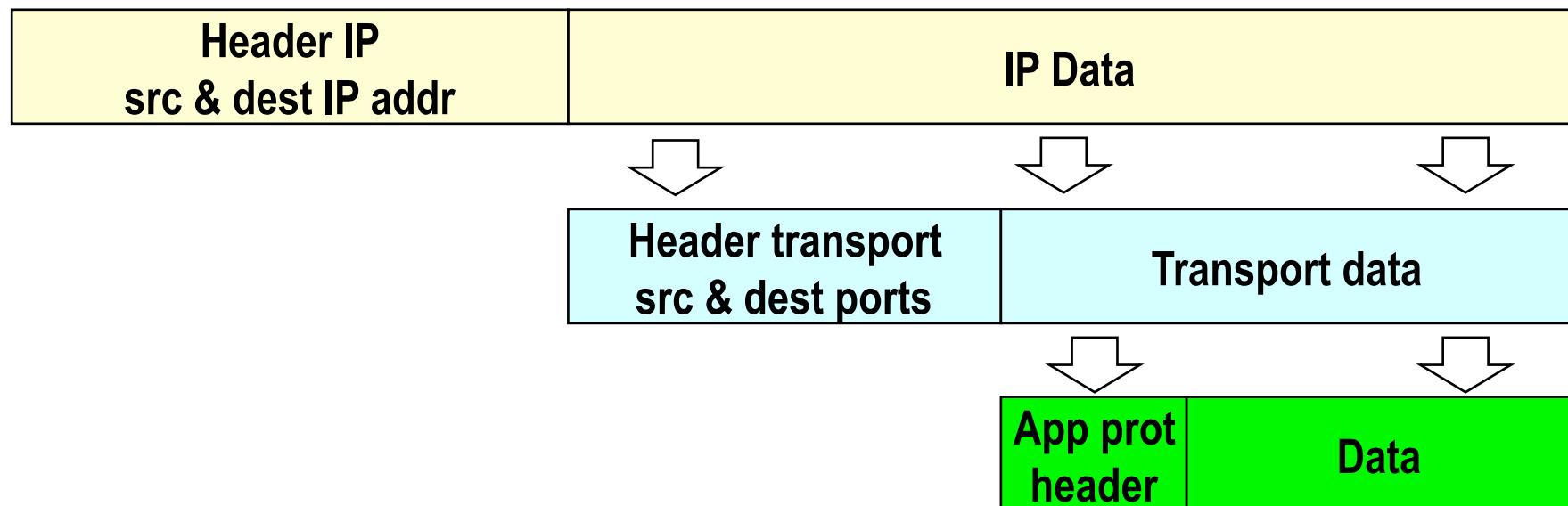
Nell'header IP il campo
Protocol vale 6 se il pacchetto
deve essere passato a TCP, 17
se deve essere passato a UDP



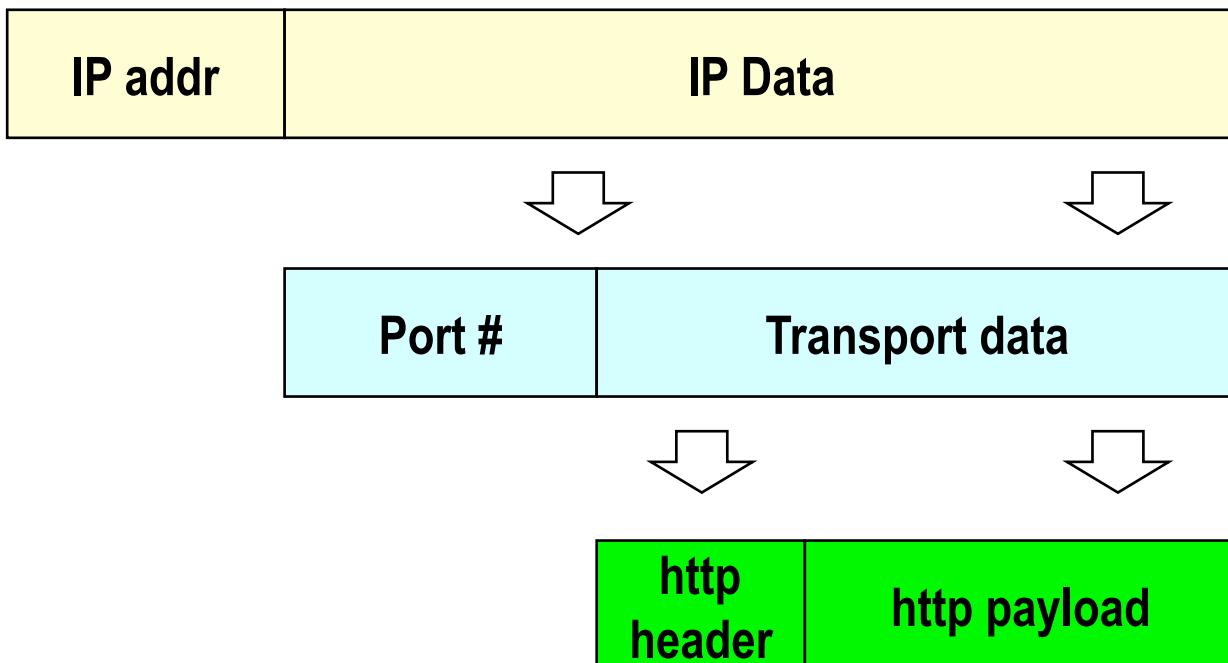
The Internet level view



Information units travelling in the network: IP packets



Demultiplexing at receiver



IP SW: checks IP packet;
Sends to transport sw
selects whether UDP or TCP
Transport demux

Transport SW: checks segment;
Sends to application sw based on Port number

Programmi Client e Server

- Un client può essere eseguito in modalità parallela o seriale
 - esempio: invia più richieste in parallelo per i file che compongono una pagina web
- Anche un server può essere eseguito in modalità parallela o seriale
- Normalmente i server che usano TCP vengono eseguiti in modalità parallela e sono dunque in grado di rispondere a più richieste contemporaneamente
- Con ognuno dei client viene aperta una connessione TCP che viene mantenuta per il tempo necessario a scambiare richieste e risposte
- La gestione delle procedure per ciascun client collegato avviene mediante la generazione di processi figli (per clonazione del processo o uso di processi multi-thread)

What transport service does an app need?

Data loss

- some apps (e.g., audio) can tolerate some loss
- other apps (e.g., file transfer, telnet) require 100% reliable data transfer

Bandwidth

- some apps (e.g., multimedia) require minimum amount of bandwidth to be “effective”
- other apps (“elastic apps”) make use of whatever bandwidth they get

Timing

- some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

Transport service requirements of common apps

Application	Data loss	Bandwidth	Time Sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps video:10kbps-5Mbps	yes, 100's msec
stored audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	few kbps up	yes, 100's msec
instant messaging	no loss	elastic	yes and no

Internet transport protocols services

TCP service:

- ❑ *connection-oriented*: setup required between client and server processes
- ❑ *reliable transport* between sending and receiving process
- ❑ *flow control*: sender won't overwhelm receiver
- ❑ *congestion control*: throttle sender when network overloaded
- ❑ *does not providing*: timing, minimum bandwidth guarantees

UDP service:

- ❑ unreliable data transfer between sending and receiving process
- ❑ does not provide: connection setup, reliability, flow control, congestion control, timing, or bandwidth guarantee

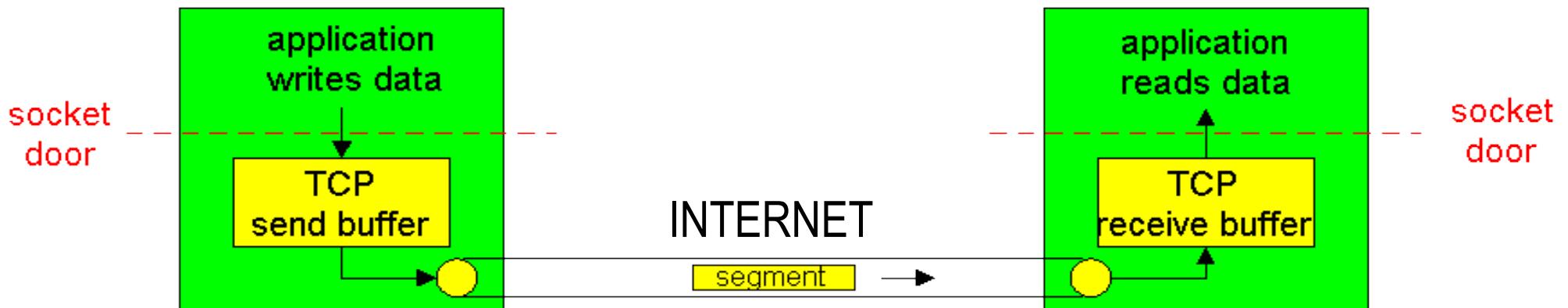
Q: why bother? Why is there a UDP? Which service does it add to IP?

Multiplexing/Demultiplexing
Error Control

Internet apps: application, transport protocols

Application	Application layer protocol	Underlying transport protocol
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	proprietary (e.g. RealNetworks)	TCP or UDP
Internet telephony	proprietary (e.g., Dialpad)	typically UDP

Why it is trivial (!) to write networking apps?



- Application software duties:
 - open socket (e.g. C, C++, JAVA function call, OS call, external library primitive)
 - Injects message in its own socket
 - being confident message is received on the other side
- TCP software: in charge of managing segments!
 - reliable message transport when TCP used
 - Segmentation performed by TCP transmitter
 - Receive buffer necessary to ensure proper packet's order & reassembly

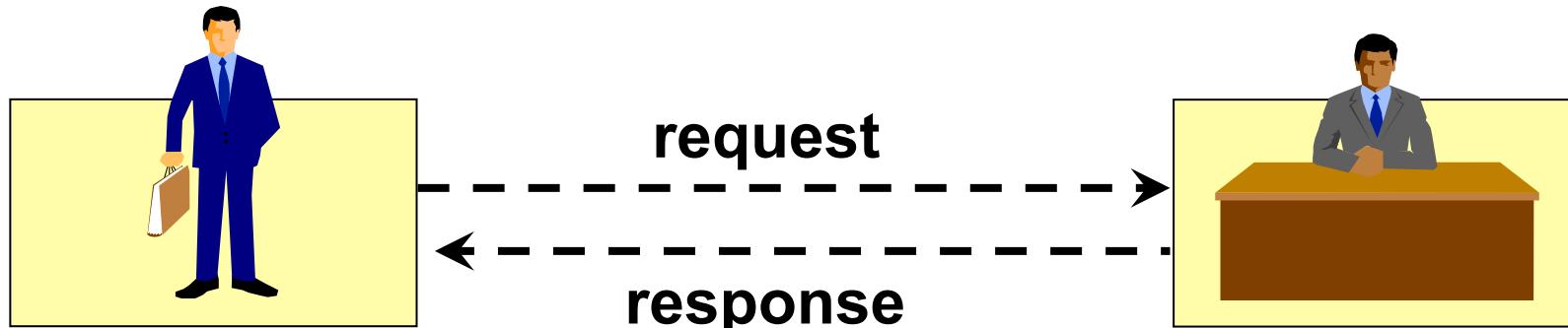
1. Protocol

□ HTTP: the protocol of the WWW

- version 1.0, RFC 1945, may 1996
- version 1.1, RFC 2068 (jan97), RFC 2616 (jun99)

(but also **FTP**: file transfer protocol, **TELNET**: opens a telnet window, **FILE**: access local file etc.)

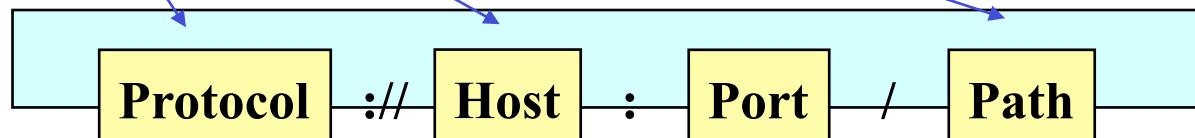
□ HTTP: Request-response protocol. A request message is sent from the client to a recipient server. The recipient server send a response message back.



□ Different requests for the same URL can result in different responses due to the time of request, changes in the resource, the request header fields.

The three components of an URL

1. *Protocol (also called "scheme")*
 - *how can a page be accessed? (application protocol used)*
 - **http://www.dsi.uniroma1.it/people/petrioli/index.html**
2. *Host name*
 - *Where is the page located? (symbolic or numeric location)*
 - **http://www.dsi.uniroma1.it/people/petrioli/index.html**
3. *File (resource) name*
 - *What is the page called? (with full path)*
 - **http://www.dsi.uniroma1.it/people/petrioli/index.html**

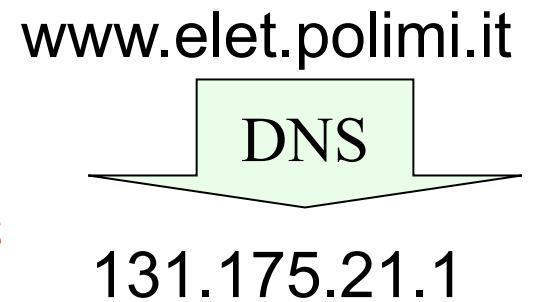


2. Location - host name

Specifies where is the page located:

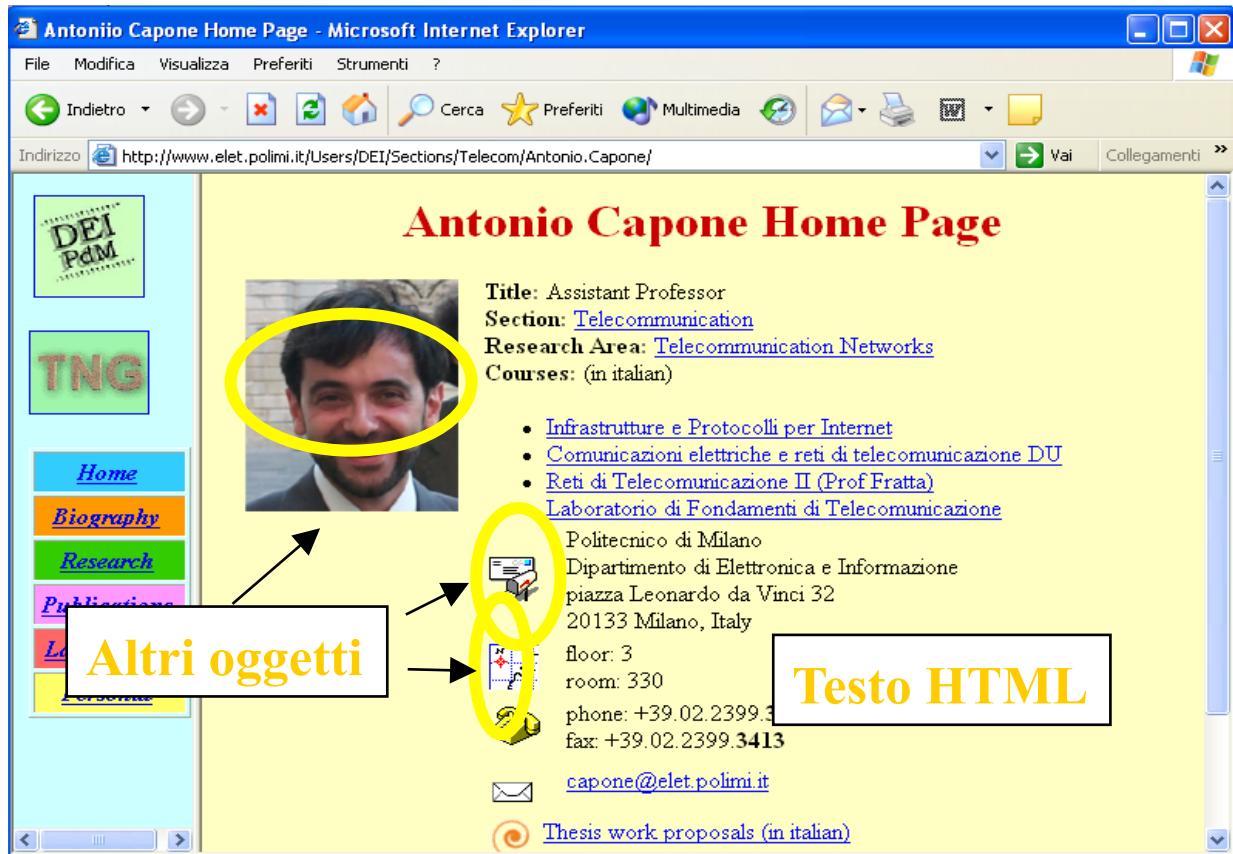
- on which host

- Humans understand names;
- Machines prefer numbers!
- Domain Name System (DNS) protocol:
 - translates names in numbers



Trasporto dei messaggi

- Supponiamo che un client richieda una pagina HTML di un server al cui interno sono contenuti i riferimenti ad altri oggetti (ad esempio 10 figure che compongono la pagina e che occorre visualizzare insieme al testo HTML).



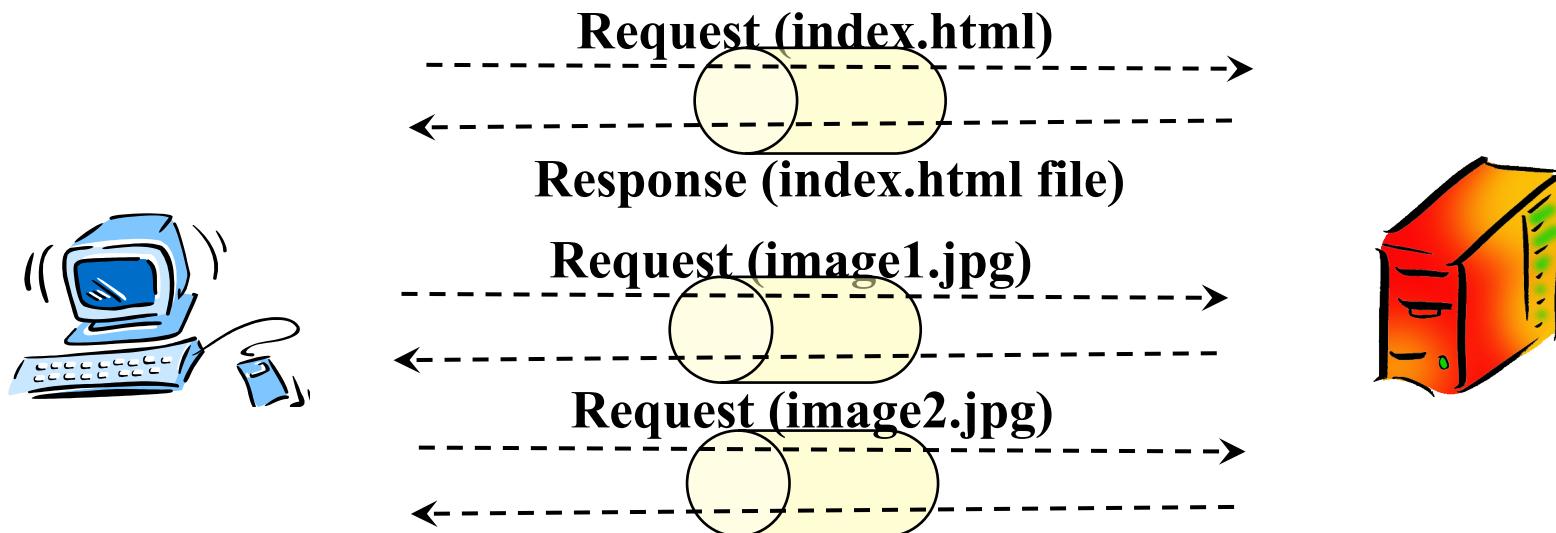
■ nel trasferimento dell'insieme di oggetti sono possibili 2 modalità

■ Non-persistent connection (default mode di HTTP 1.0)

■ Persistent connection (default mode di HTTP 1.1)

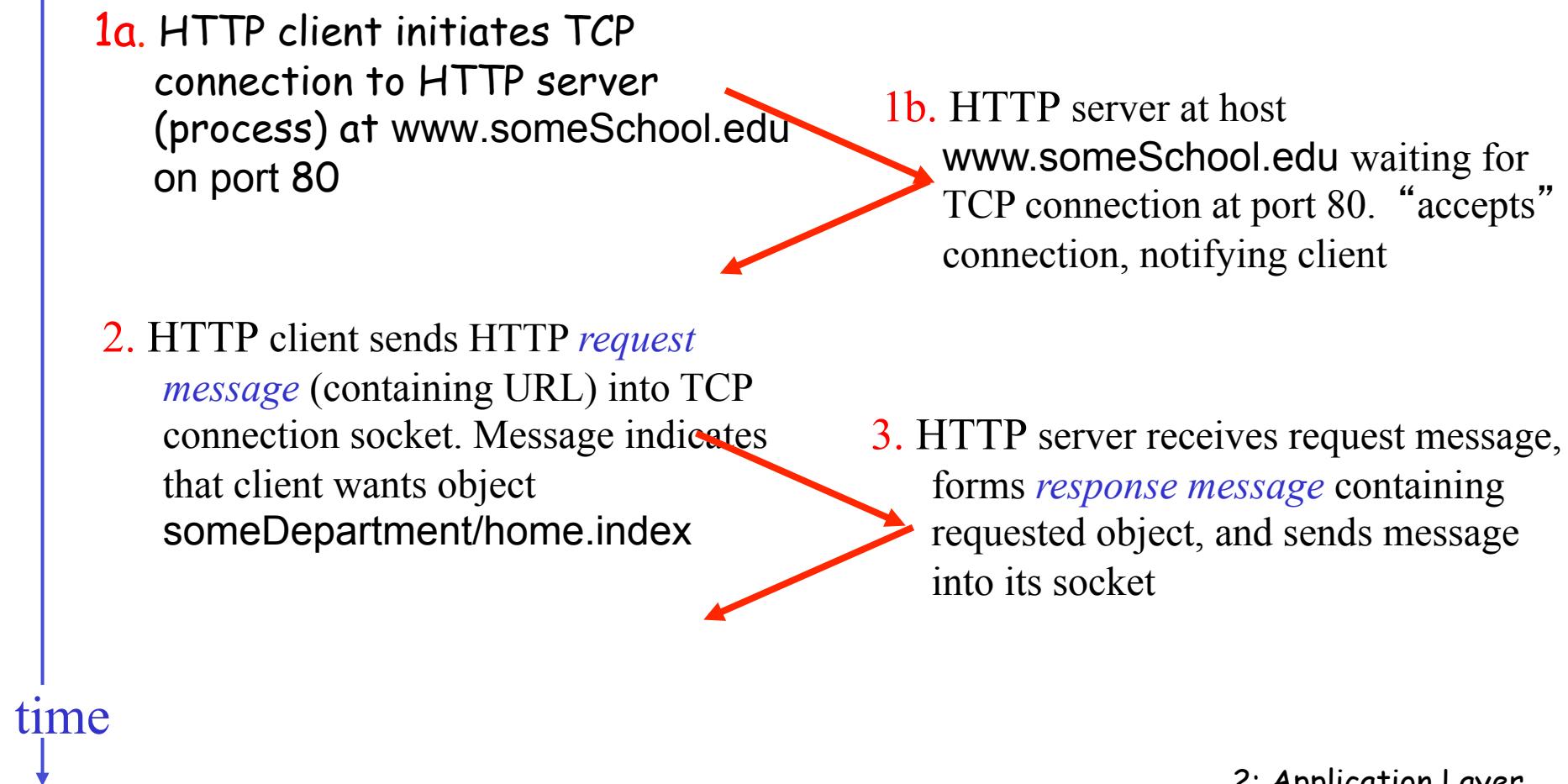
Non persistent

- Viene aperta una connessione per una sola request-response: inviato l'oggetto, il server chiude la connessione
- La procedura viene ripetuta per tutti i file collegati al documento HTML base



Nonpersistent HTTP (HyperText Transfer Protocol)

Suppose user enters URL `www.someSchool.edu/someDepartment/home.index` (contains text, references to 10 jpeg images)



Nonpersistent HTTP (cont.)

time ↓

4. HTTP server closes TCP connection.
5. HTTP client receives response message containing html file, displays html. Parsing html file, finds 10 referenced jpeg objects
6. Steps 1-5 repeated for each of 10 jpeg objects

Response time modeling

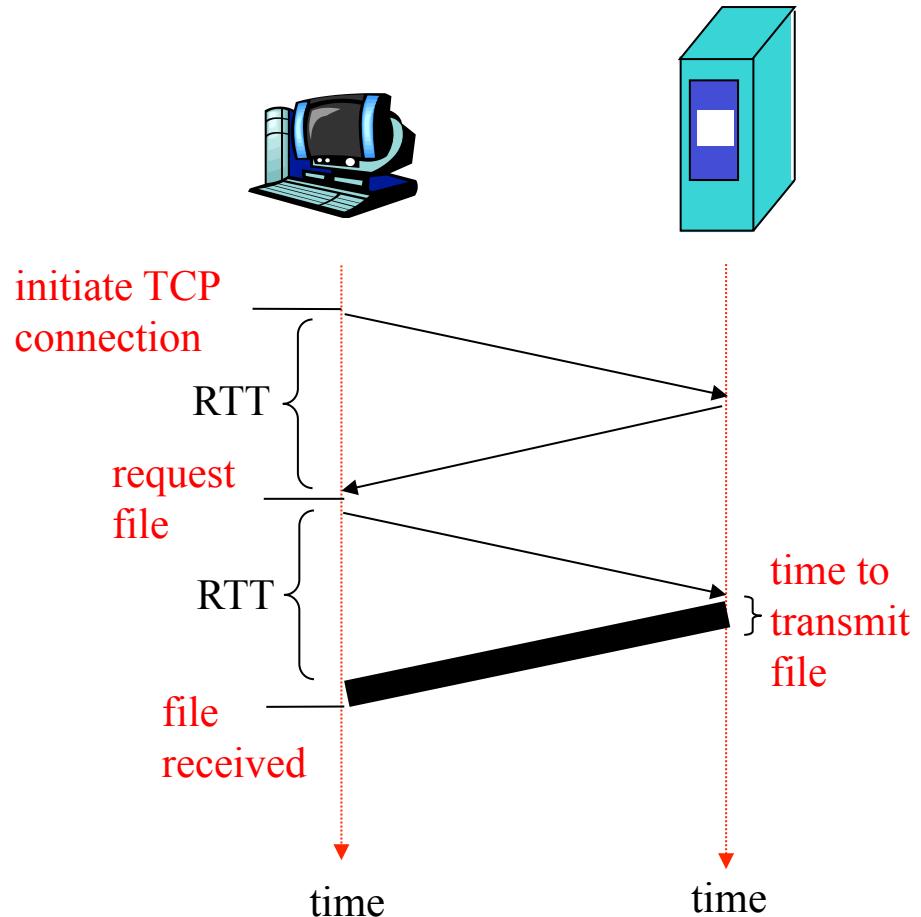
Definition of RRT: time to send a small packet to travel from client to server and back.

Response time:

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time

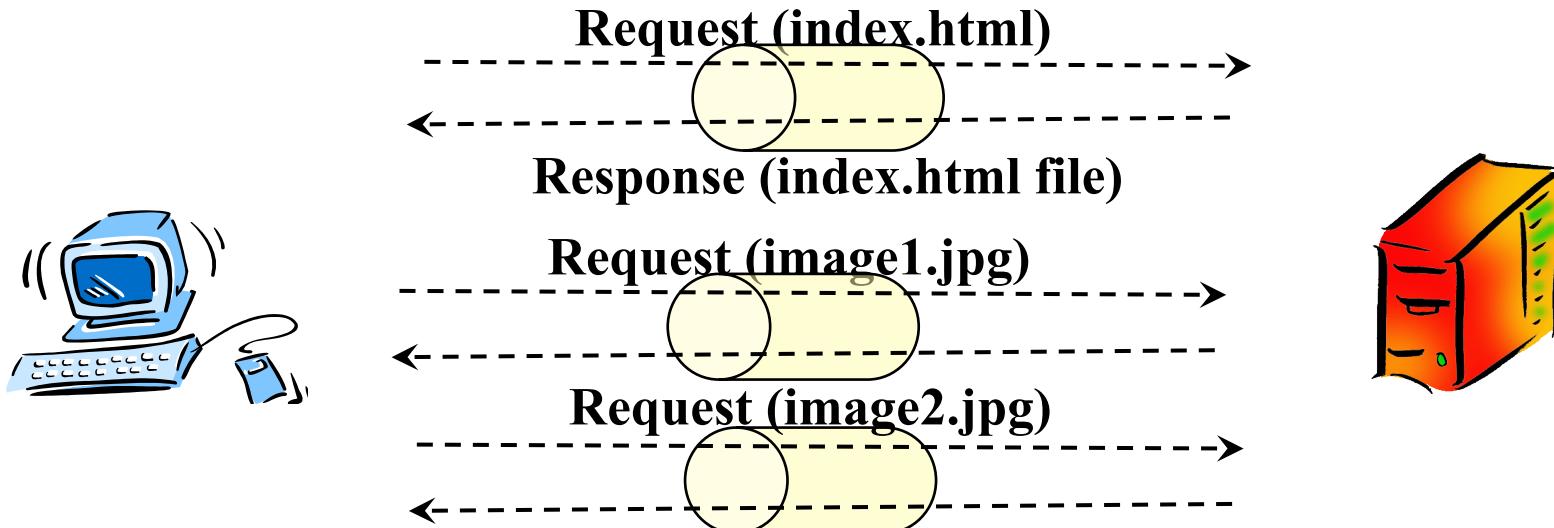
$$\text{total} = 2\text{RTT} + \text{transmit time}$$

Problema: OGNI VOLTA CHE SI APRE UNA CONNESSIONE SI PAGA 2RTT PRIMA DI SCAMBIARE DATI



Non persistent

- Viene aperta una connessione per una sola request-response: inviato l'oggetto, il server chiude la connessione
- La procedura viene ripetuta per tutti i file collegati al documento HTML base
- Le connessioni TCP per più oggetti possono essere aperte in parallelo per minimizzare il ritardo
- Il numero max di connessioni è di solito configurabile nel browser



Un esempio...

- Esempio: **di richiesta oggetto**

```
GET /ntw/index.html HTTP/1.0
Date: Wed, 22 Mar 2000 09:09:01 GMT
Pragma: No-cache
From: gorby@moskvax.com
User-agent: Mozilla/4.3
```

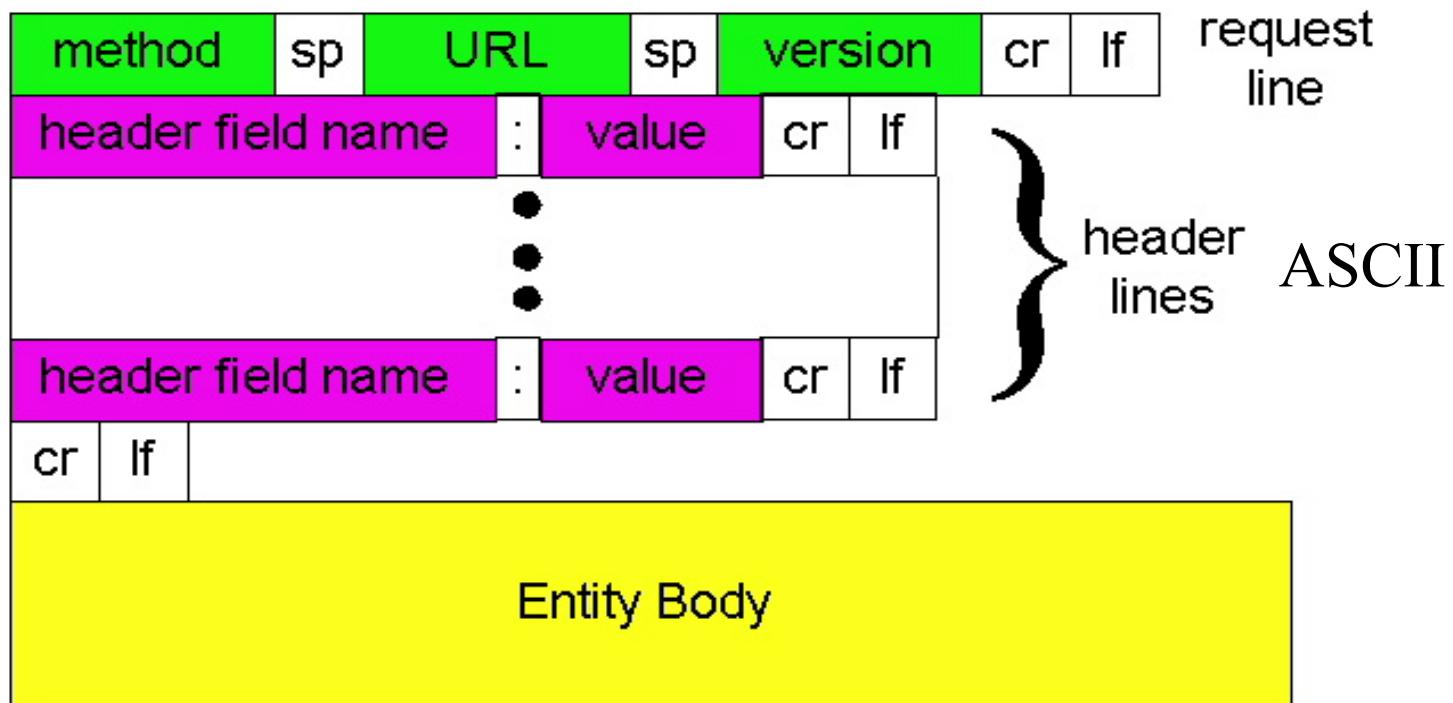
HTTP è testuale (ASCII)

- Esempio: **risposta**

```
HTTP/1.0 200 OK
Date: Wed, 22 Mar 2000 09:10:01 GMT
Server: Apache/1.3.0 (Unix)
Content-Length: 6821
Last-Modified: Mon, 22 Jun 1998 09:23:24 GMT
Content-Type: text/html
data data data data data ...
```

HTTP Messages

Request



HTTP Message

Methods:

GET	E' usato quando il client vuole scaricare un documento dal server. Il documento richiesto è specificato nell'URL. Il server normalmente risponde con il documento richiesto nel corpo del messaggio di risposta.
HEAD	E' usato quando il client non vuole scaricare il documento ma solo alcune informazioni sul documento (come ad esempio la data dell'ultima modifica). Nella risposta il server non inserisce il documento ma solo degli header informativi.
POST	E' usato per fornire degli input al server da utilizzare per un particolare oggetto (di solito un applicativo) identificato nell'URL.
PUT	E' utilizzato per memorizzare un documento nel server. Il documento viene fornito nel corpo del messaggio e la posizione di memorizzazione nell'URL.

■ Altri methods:

■ **DELETE, LINK, UNLINK, ...**

Per inviare info al server

E' possibile anche usare GET

GET /sercah.cgi?string=greek-architects HTTP/1.0

Un esempio...

- Esempio: **di richiesta oggetto**

```
GET /ntw/index.html HTTP/1.0
Date: Wed, 22 Mar 2000 09:09:01 GMT
Pragma: No-cache
From: gorby@moskvax.com
User-agent: Mozilla/4.3
```

HTTP è testuale (ASCII)

- Esempio: **risposta**

```
HTTP/1.0 200 OK
Date: Wed, 22 Mar 2000 09:10:01 GMT
Server: Apache/1.3.0 (Unix)
Content-Length: 6821
Last-Modified: Mon, 22 Jun 1998 09:23:24 GMT
Content-Type: text/html
data data data data data ...
```

Header

Header name	:	Header value
--------------------	----------	---------------------

- Gli header servono per scambiare informazione di servizio aggiuntiva
- E' possibile inserire più linee di header per messaggio
- Esempi

Cache-control	Informazione sulla cache
Accept	Formati accettati
Accept-language	Linguaggio accettato
Authorization	Mostra i permessi del client
If-modified-since	Invia il doc. solo se modificato
User-agent	Tipo di user agent

Header Types

- General: used in request and response messages
- Request header (e.g. to express preferences on the nature of the response, include additional info with the request, to specify a constraint on the server in handling the request)
- Response header: to provide additional info about the response or to request additional info from the user
- Entity header (both in request and response messages) to provide info on the entity such as the last time it was modified

Header Types-Request

- General: used in request and response messages

Date	Indicates the date and time of the message origination, e.g. Date: Tue 16 May 2000 11:29:32 GMT
Pragma	Permits to send directives to the recipient, requesting it to behave in a particular way while handling a request or response. Pragma: no-cache informs proxies on the path not to return a cached copy

- Request header (e.g. to express preferences on the nature of the response, include additional info with the request, to specify a constraint on the server in handling the request)

Authorization	Includes credentials required to access a resource (e.g. Authorization: Basic YXZpYXRpS29IDizDizM1NA== where Basic is the authentication scheme and YXZpYXRpS29IDizDizM1NA== an encoding of user id and password)
From	The user may include e-mail address for identification (From: gorby@moskvax.com)
If-Modified-Since	Indicates not to return a copy of the resource if it has not been modified after the specified date. Used for caching at a proxy or browser (saves the time needed for downloading of the resource). GET /foo.html HTTP/1.0 If-Modified-Since: Sun, 21 May 2000 07:00:25 GMT
Referer	Includes the URL from which the requested URL was obtained. Can be used for locating obsolete links. Referer: http://www.cnn.com
User-Agent	User browser used, client machine OS, etc User-Agent: Mozilla/4.04 [en]C-WorldNet (win95;I)

Header Types

- Entity header (both in request and response messages) to provide info on the entity such as the last time it was modified

Allow	Indicates the list of valid methods that can be applied to a resource PUT /foo.html HTTP/1.0 Allow: HEAD, GET, PUT
Content-Type	Media type of the entity body (e.g. image/gif, text/html etc.)
Content-Encoding	Indicates how the resource representation can be decoded into the format indicated in the Content-Type field (e.g. if the file has been compressed with gzip) Content-Encoding:gzip
Content-Length	Length of the entity body (bytes). Allow to check whether all the body was received
Expires	The entity should be considered stale after the time specified in this header (after this time the content could still be stored in cache but should not be returned to users without first validating with the origin server)
Last-Modified	Specify the time at which the resource was modified last

Un esempio...

- Esempio: **di richiesta oggetto**

```
GET /ntw/index.html HTTP/1.0
Date: Wed, 22 Mar 2000 09:09:01 GMT
Pragma: No-cache
From: gorby@moskvax.com
User-agent: Mozilla/4.3
```

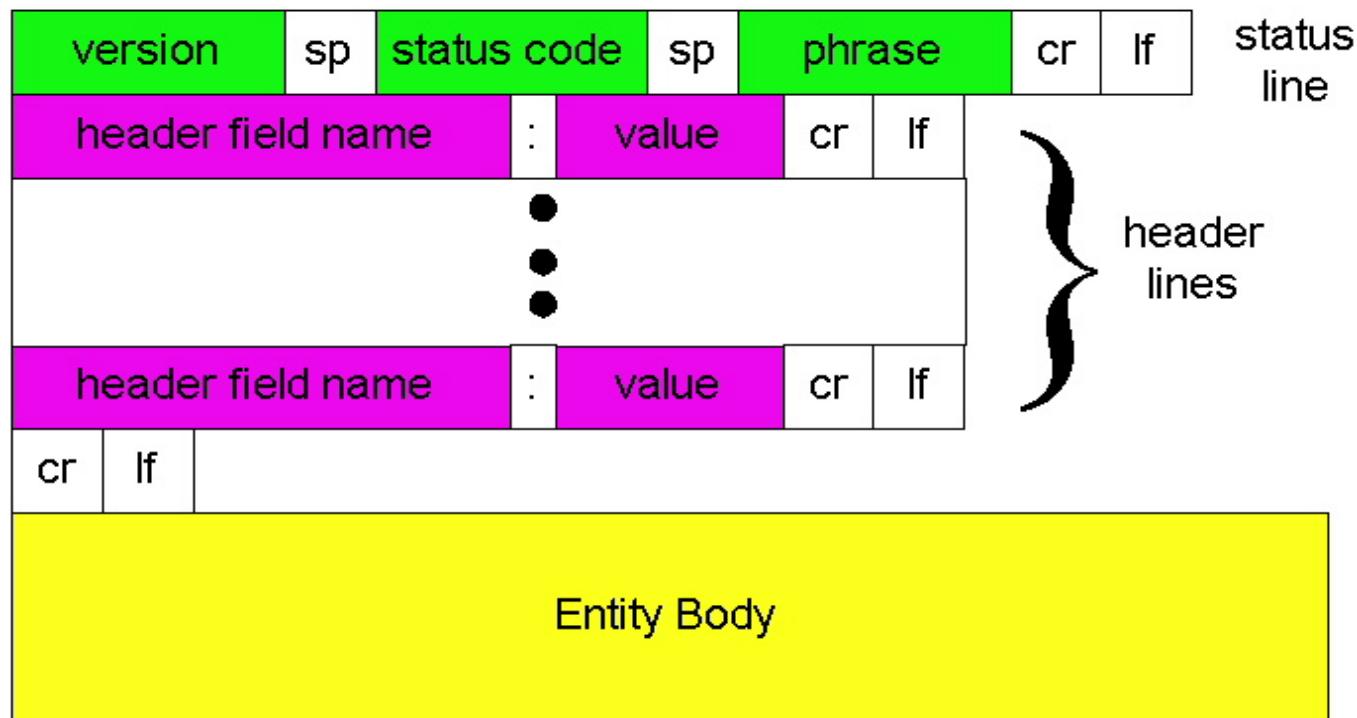
HTTP è testuale (ASCII)

- Esempio: **risposta**

```
HTTP/1.0 200 OK
Date: Wed, 22 Mar 2000 09:10:01 GMT
Server: Apache/1.3.0 (Unix)
Content-Length: 6821
Last-Modified: Mon, 22 Jun 1998 09:23:24 GMT
Content-Type: text/html
data data data data data ...
```

Message

Response



Header Types-Response

- Response header: to provide additional info about the response or to request additional info from the user

Location	Used to redirect the request to where the resource can be found Location: http://www.foo.com/level1/twosdown/Location.html
Server	Origin server SW version number and configuration related info Server: Apache/1.2.6 Red Hat
WWW-Authenticate	Request to transmit the request with appropriate credentials according to a given scheme

Messaggi

■ 1xx Informational

Status codes:

■ 2xx Success

200 OK:	La richiesta ha avuto successo; l'informazione è inclusa
201 Created:	La risorsa e' stata creata con successo in seguito ad UN POST (se non puo' essere creata subito 202 Accepted)
202 Accepted:	The request has been received but has not yet been handled in full (e.g., since a program must run which requires a long time → the user agent can continue with its task without waiting for the action to complete at the origin server)
204 No content:	Successfully received, no change required in what the user actually sees

Messaggi

■ 3xx Redirection

301 Moved Permanently:	L'oggetto è stato spostato nell'URL indicato
302 Moved Temporally:	L'oggetto è stato spostato nell'URL indicato
304 Not Modified:	L'oggetto non modificato dal tempo incluso nella richiesta

■ 4xx Client error

400 Bad Request:	errore generico (sintassi errata(irriconoscibile)
401 Unauthorized:	Autenticazione fallita (e.g., Accesso senza necessari account e password)
403 Forbidden:	La richiesta e' stata ricevuta e compresa ma il server ha stabilito di non servirla. Le ragioni possono essere indicate nell'entity body.
404 Not Found:	l'oggetto non esiste sul server

■ 5xx Server error

500 Internal server error	Generico. Errore o guasto nel server
501 Not implemented	Funzione non implementata
503 Service unavailable	Servizio non disponibile

Un esempio

□ Esempio: di richiesta oggetto

```
PUT /motd HTTP/1.0
Date: Wed, 22 Mar 2000 08:10:07 GMT
From: gorby@moskvax.com
User-agent: Mozilla/4.0
Content-length:23
Allow:GET,HEAD,PUT
data data data data data ...
```

HTTP è testuale (ASCII)

□ Esempio: risposta

```
HTTP/1.1 200 OK
Date: Wed, 22 Mar 2000 09:10:07 GMT
Server: Apache/1.3.0 (Unix)
...
```

Cache locale: get condizionato

- E' possibile evitare di scaricare oggetti memorizzati nella memoria locale se non sono stati modificati

Client:

```
GET /fruit/kiwi.gif HTTP/1.0
User-agent: Mozilla/4.0
Accept: text/html, image/gif, image/jpeg
If-modified-since: Mon, 22 Jun 1998 09:23:24
```

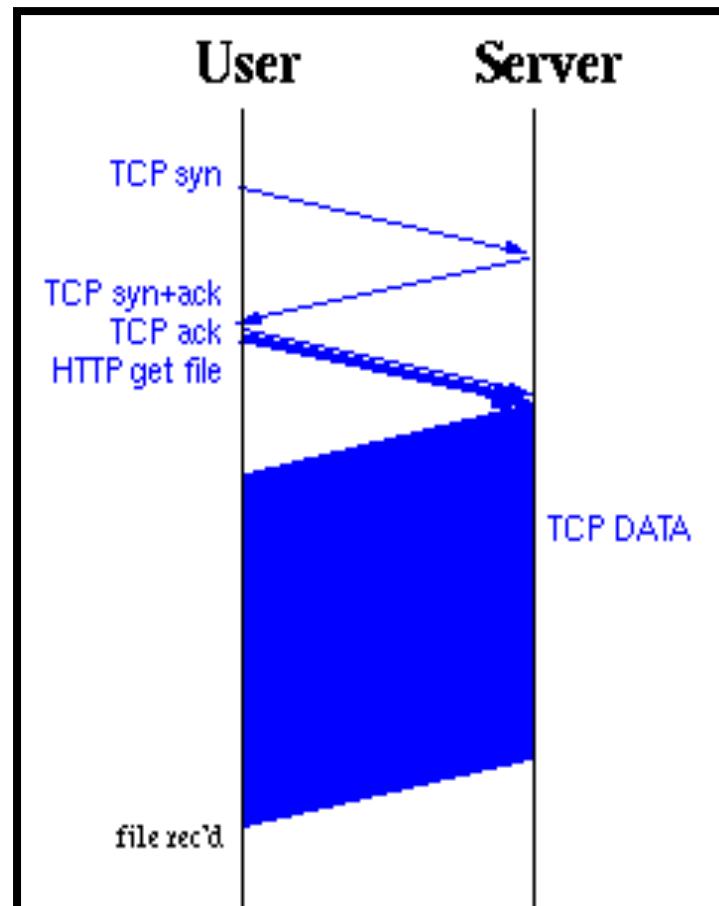
Server:

```
HTTP/1.0 304 Not Modified
Date: Wed, 19 Aug 1998 15:39:29
Server: Apache/1.3.0 (Unix)
(empty entity body)
```

- E' possibile anche usare il metodo HEAD

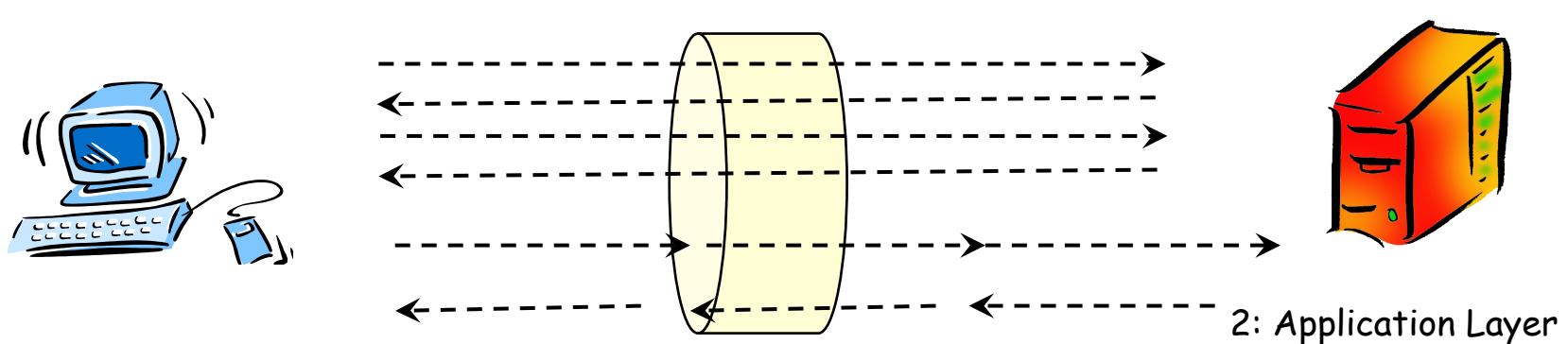
Performance drawbacks

- Mandatory roundtrips
 - TCP three-way handshake
 - get request, data return
 - new connections for each inlined image (parallelize)



Persistent connection (HTTP v1.1)

- Nel caso *persistent* il server non chiude la connessione dopo l'invio dell'oggetto
- La connessione rimane aperta e può essere usata per trasferire altri oggetti della stessa pagina web o anche più pagine
- I server chiudono di solito la connessione sulla base di un *time-out*
- **without pipelining:** il client invia una nuova richiesta solo dopo aver ricevuto la risposta per la precedente
 - *with pipelining:* più richieste vengono inviate consecutivamente dal client



Persistent HTTP

Nonpersistent HTTP issues:

- ❑ requires 2 RTTs per object
- ❑ OS must work and allocate host resources for each TCP connection
- ❑ but browsers often open parallel TCP connections to fetch referenced objects

Persistent HTTP

- ❑ server leaves connection open after sending response
- ❑ subsequent HTTP messages between same client/server are sent over connection

Persistent without pipelining:

- ❑ client issues new request only when previous response has been received
- ❑ one RTT for each referenced object

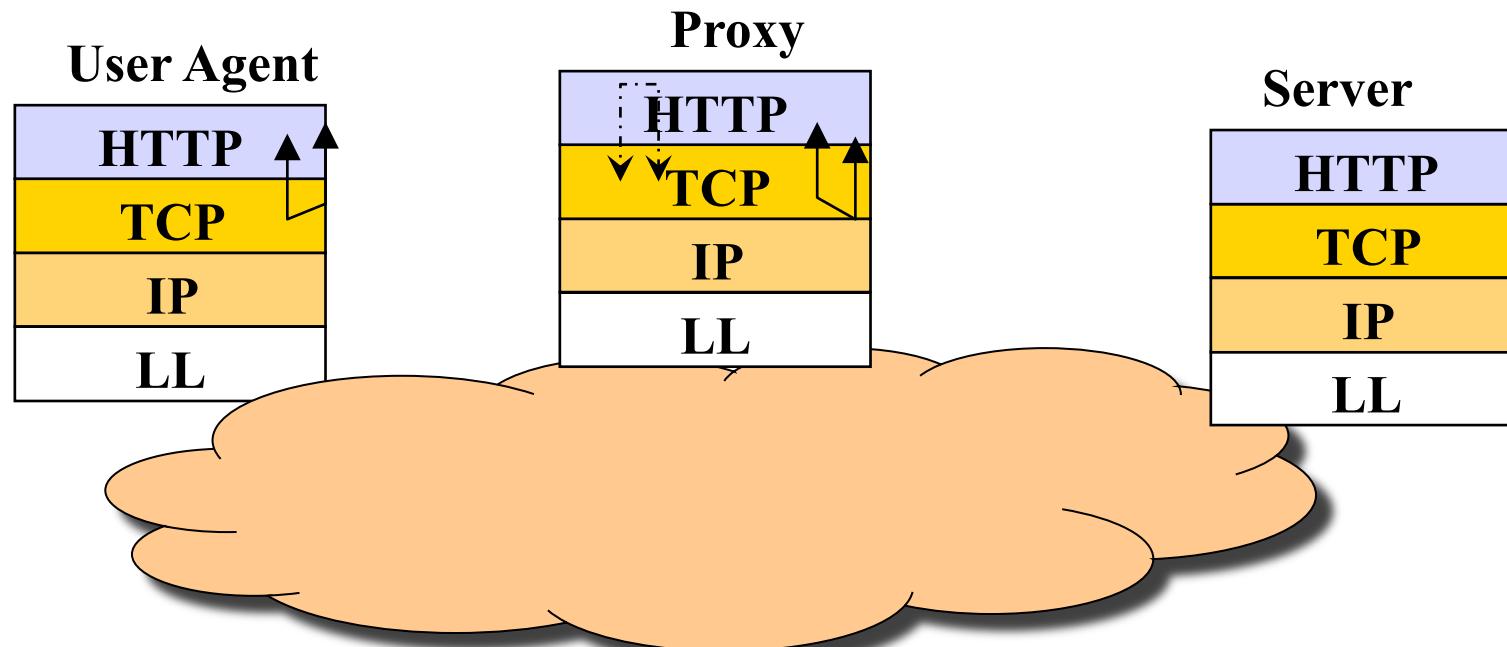
Persistent with pipelining:

- ❑ default in HTTP/1.1
- ❑ client sends requests as soon as it encounters a referenced object
- ❑ as little as one RTT for all the referenced objects

Proxy

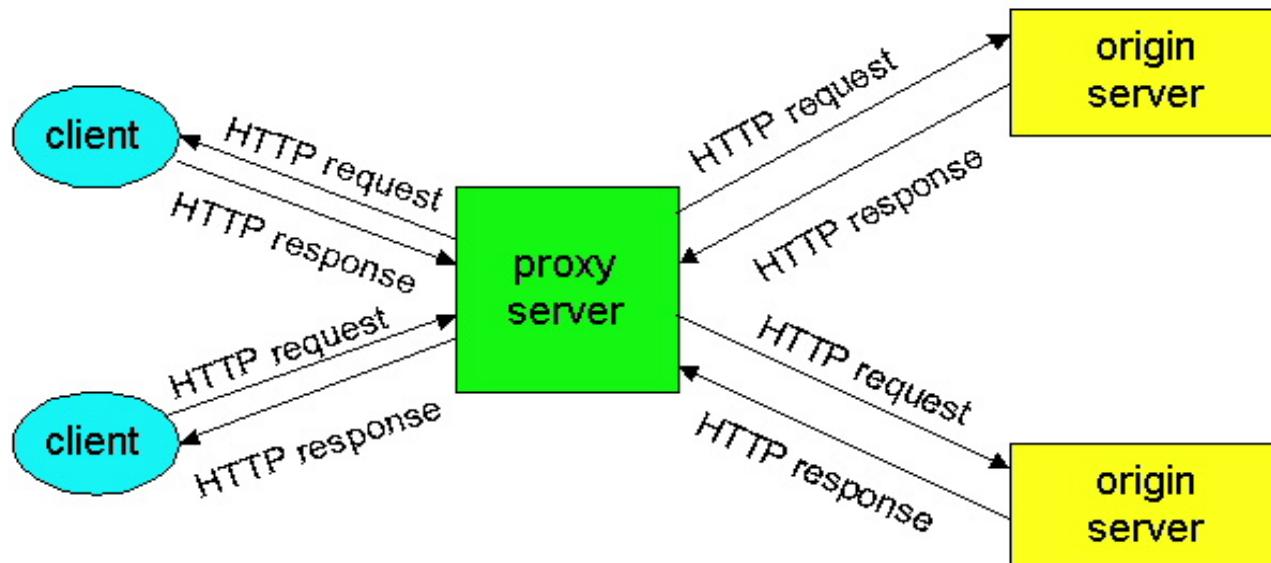
An intermediary program which acts as both a server and a client for the purpose of making requests on behalf of other clients.

- I proxy sono degli instradatori di messaggi di livello applicativo
- Devono essere sia client che server
- Il server vede arrivare tutte le richieste dal proxy (mascheramento degli utenti del proxy)



Cache di rete: uso dei proxy

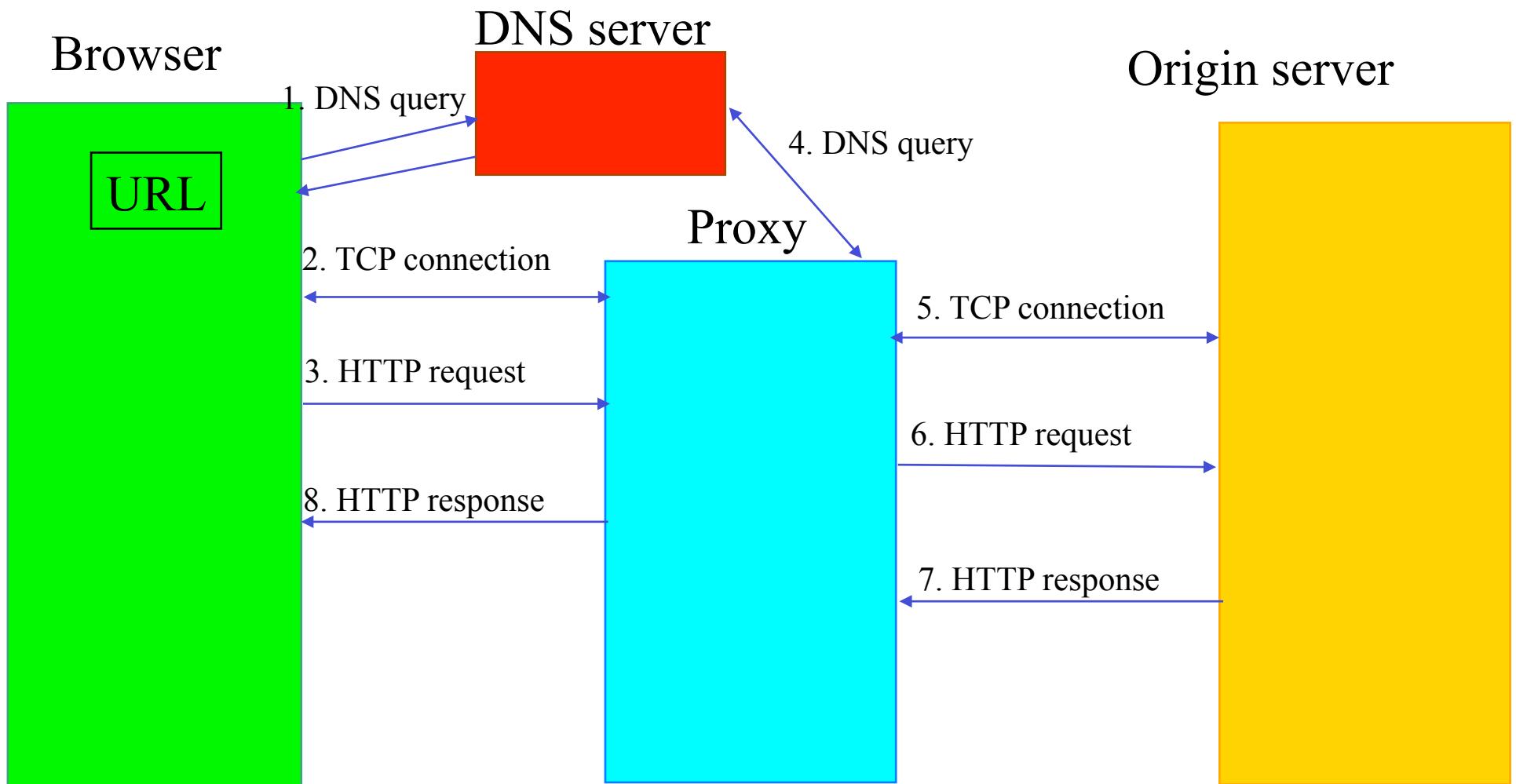
- Compito principale dei proxy è fornire una grande memoria di cache
- Se un documento è contenuto nella cache viene scaricato più velocemente sul client
- Razionale: i server forniscono le stesse info o info fortemente ridondanti a comunità locali di utenti



□ Consistency: cached pages might not be the most recent...

□ Ad clickthrough counts: how does Yahoo know how many times you accessed their pages, or **more importantly**, their ads?

Download attraverso un proxy...



Tipi di proxy

- Regular proxies
 - forwards requests and responses
- Caching proxy
 - has a cache of responses received in the past
 - when the proxy receives a request that can be satisfied by the cache the request message is not forwarded and the response is returned directly by the proxy
- Forwardng behavior
 - Transparent proxy
 - Only superficially changes the request (e.g., adding its ID)
 - Non transparent proxy
 - Can modify request/response
 - Anonymization
 - Media type transformation (e.g. changing image format to reduce size)
 - Language tranformation ...
 - Entrambi i transparent/non-transparent proxy possono avere associata una cache

Uso dei Proxy(1)

- Caching: ‘*storage of a response obtained earlier for later use, when clients request the same resource*’. Se la pagina richiesta e’ nella cache il proxy puo’ inviarla immediatamente al client purche’ sia fresh (ovvero non sia stata modificata rispetto alla copia in cache)
 - latency minore MA cache consistency importante. Strong consistency/weak consistency
- Agisce come front-end per un gruppo di utenti che condividono l’accesso al web (caching ma anche suddivisione delle risorse—pro: se piu’ utenti richiedono la stessa risorsa ad un server un’unica connessione deve essere stabilita; -- con: possibili ritardi dovuti al fatto che piu’ utenti condividono le stesse risorse)

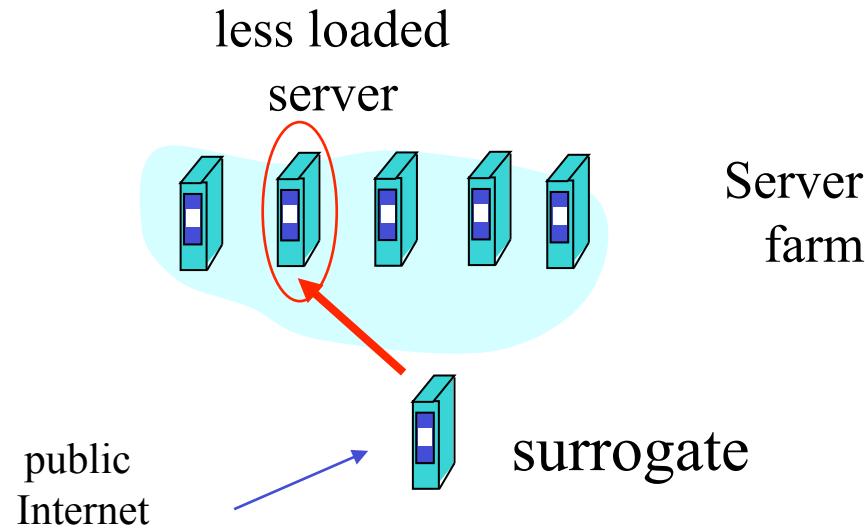
Uso dei Proxy(2)

- Mascheramento degli utenti del proxy: Il server vede arrivare tutte le richieste dal proxy e quindi non e' in grado di riconoscere l'utente. Il proxy conosce le informazioni di accesso relative agli utenti ma e' considerata un'entita' trusted.
 - Alcuni proxy non permettono l'anonymizzazione e indicano nell'header informazioni sul client
 - Mascheramento non completo: e.g., user agent e cookies
- Customizzazione delle risposte
- Filtraggio delle richieste e delle risposte
 - per limitare l'accesso a certi siti
 - filtrare richieste a siti non adatti ai minori
 - verificare la presenza di virus prima di inviare al client le risposte
 - togliere alcune informazioni dall'header prima di inviare la richiesta (es. e-mail dell'utente)

Other types of proxies

□ Reverse proxy (surrogates)

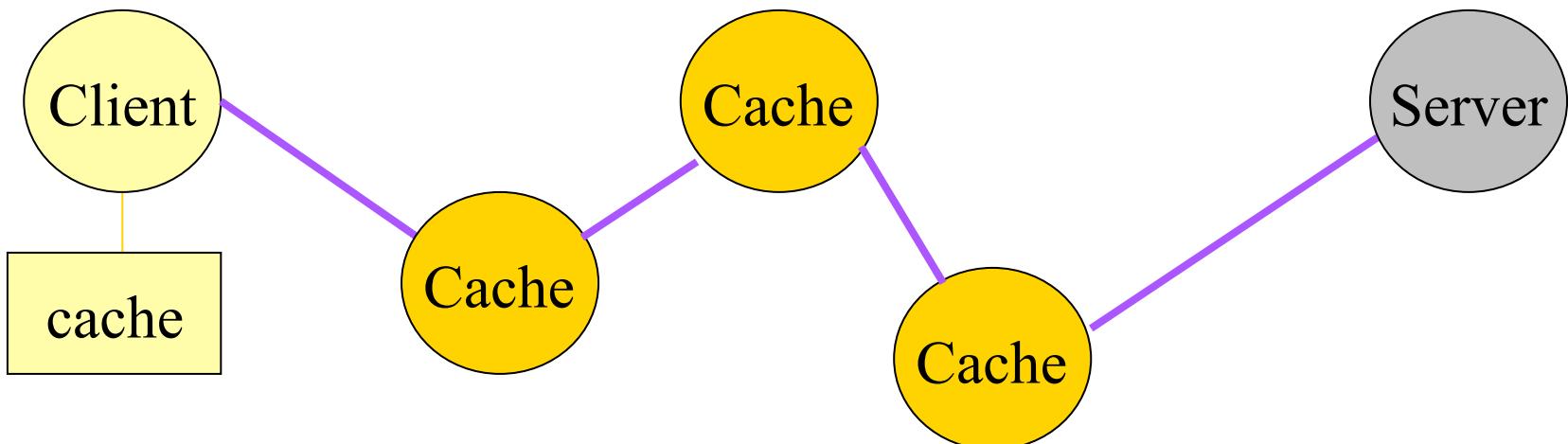
- Proxy placed close to the origin server to reduce the load on them.
Load balancing among servers of a server farm



- ## □ Interception proxy.
- Are invisible to the user. All messages pass through the interception proxy that can intercept them, examine the request, generate a response locally or redirect the request to another place (e.g. to the '**best replica**')

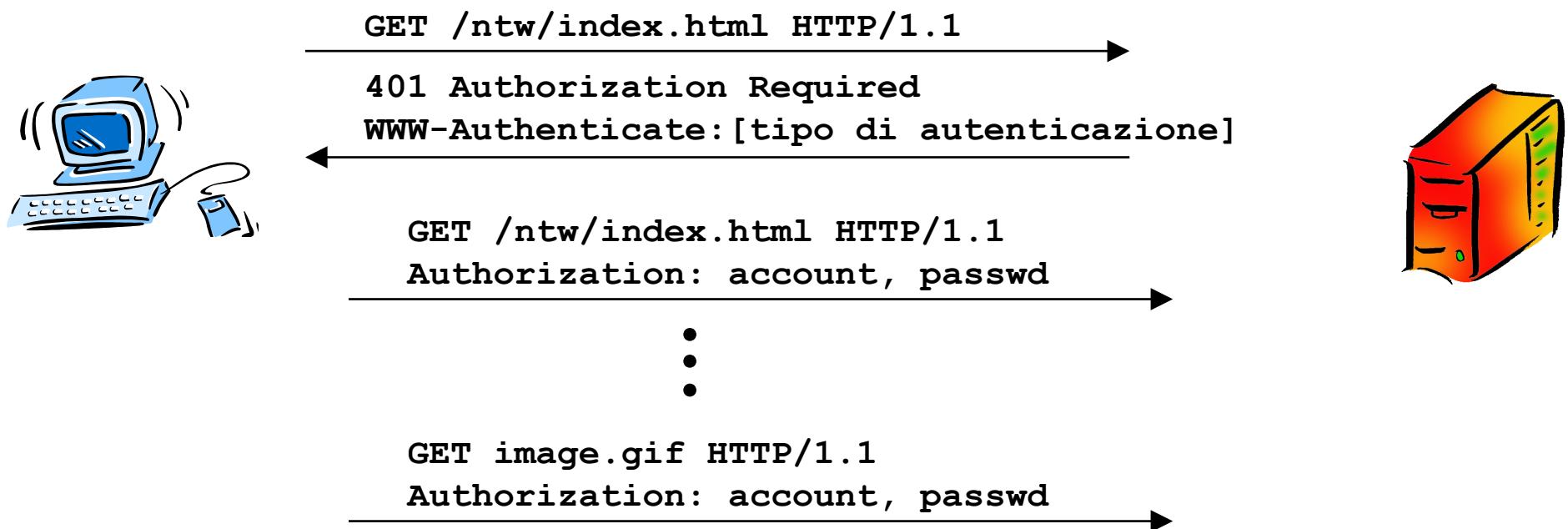
Hierarchical Caching

- The message being transported can cross several caches (e.g. at the department, ISP side, in front of the origin server -reverse proxies)
- Decreases communication costs
- Regional/national cache → high hit ratio



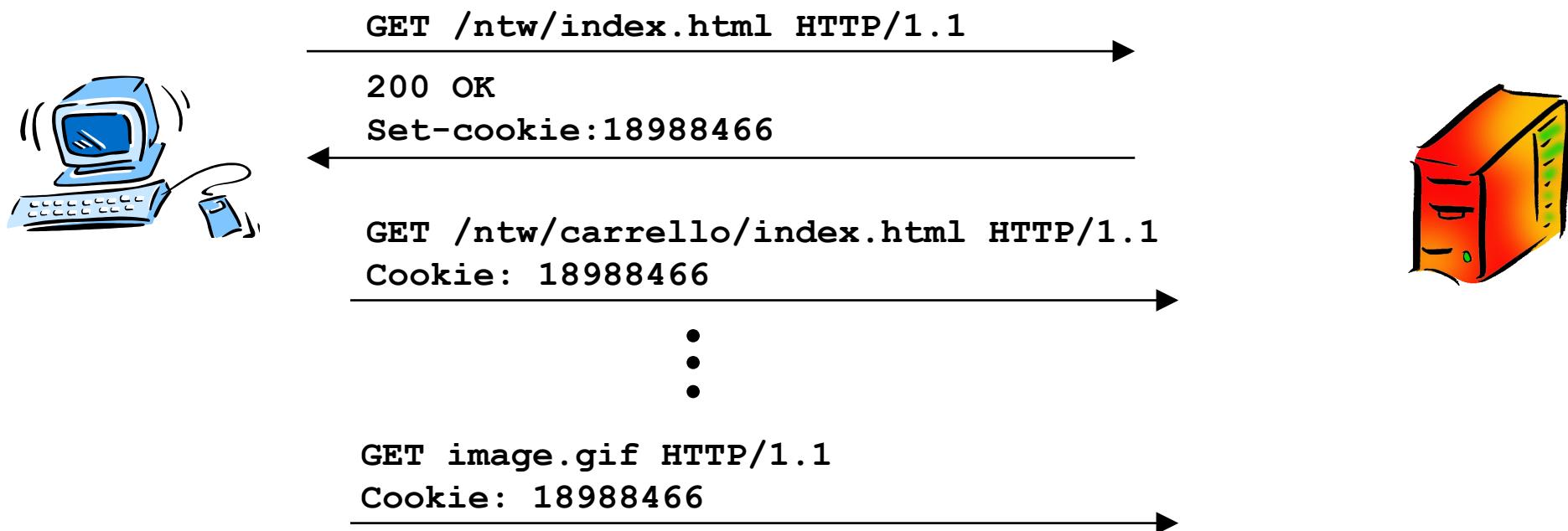
Autenticazione

- HTTP è stateless e quindi non si possono riconoscere richieste successive dello stesso utente
- in HTTP esiste un elementare meccanismo di autenticazione (account e password) che serve a riconoscere gli utenti
- Normalmente il browser memorizza passwd e account in modo da non richiedere la digitazione ogni volta



Cookie

- ❑ Esiste anche un altro modo per riconoscere richieste successive di uno stesso utente che non richiede di ricordare password
- ❑ Il numero di cookie inviato dal server viene memorizzato in un opportuno file
- ❑ Tramite i cookie si può mantenere uno stato “virtuale” per ciascun utente.

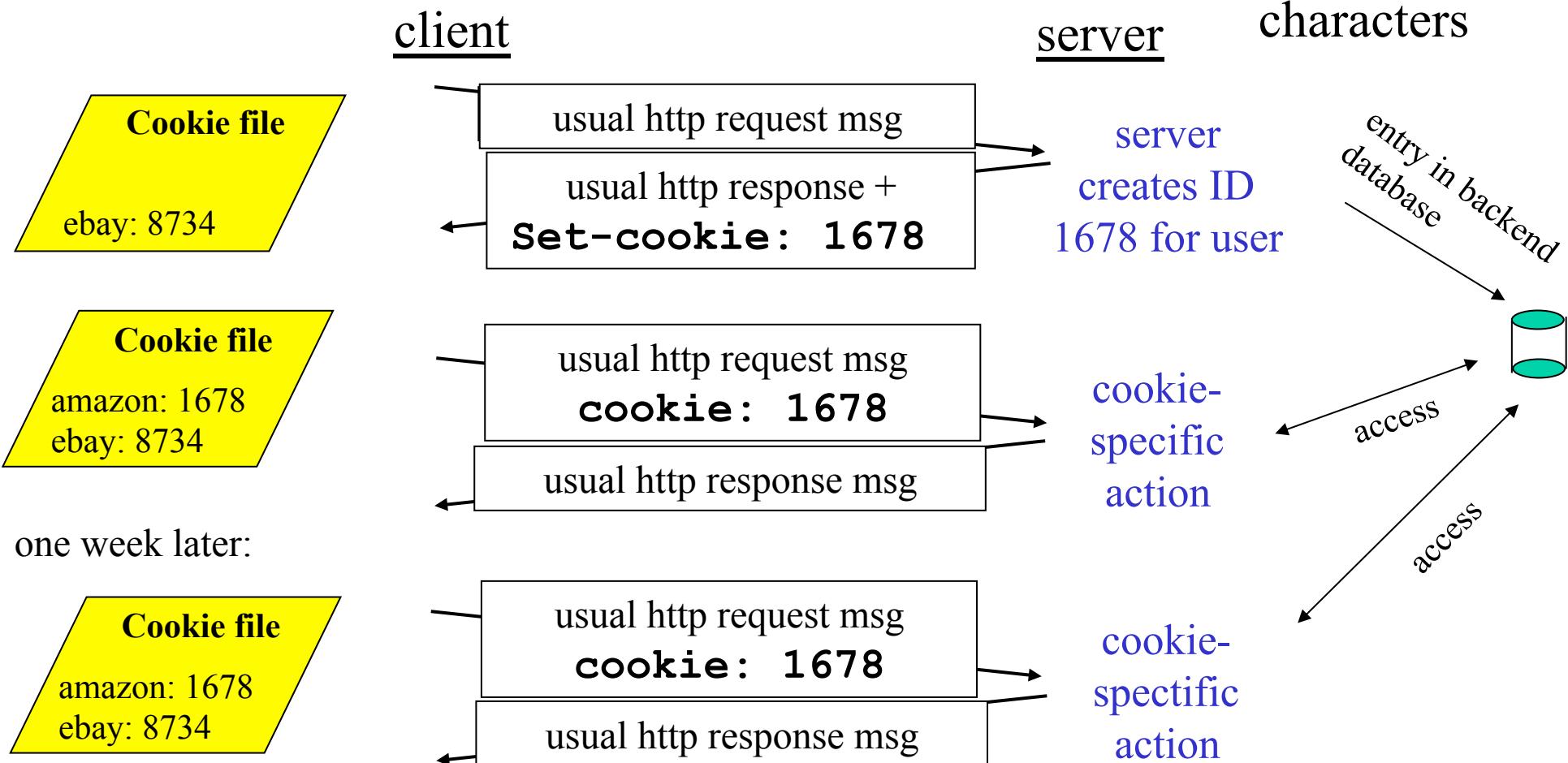


Cookie

- Permette di identificare gli utenti e mantenere delle info di stato su una transazione che richiede vari scambi di messaggi HTTP
 - Personalizzazione
 - un sito di e-commerce può personalizzare la risposta con il nome dell'utente e suggerimenti di acquisto personalizzati
 - Mantenere informazioni tra sessioni o interazioni
 - esempio: shopping cart

Cookies: keeping “state” (cont.)

Cookie=
string of
characters



Identifier to find user info on the back end database

Cookie

- User control on cookies
 - Whether to accept any cookies at all
 - Set a limit on the size and number of cookies
 - Limit the sites/domains from which a cookie can be accepted
 - Limit to a specific session acceptance of cookies
 - Require that cookies must originate from the same server as the current page being viewed
- Privacy concerns
 - Allow to track user behavior
 - Users typically not aware of when a cookie is sent
 - and on the use of the information he/she is providing through use of cookies
 - user profiling, shared between companies
 - Could also be sent from a Web Server different from the one user is aware to being connected to
 - e.g. due to redirection to download some objects referred to in the Main Server web page

Breve storia del Web

- 1945: Vannevar Bush in ‘As we may think’ proposes Memex to extend human memory via mechanical means. V. Bush wanted to make the existing store of knowledge (fastly growing) more accessible to mankind. ‘A memex is a device in which an individual stores all his books, records, and communications, and which is mechanized so that it maybe consulted with exceeding speed and flexibility. It is an enlarged supplement to his memory. ... Wholly new forms of encyclopedias will appear ready made with a mesh of **associative trails** running through them’
- 1965 Ted Nelson coins the term ‘hypertext’ to describe nonsequential writing that presents information as a collection of linked nodes. ‘Readers can pursue the information in a variety of ways by navigating from one node to another’
- 1970s scientific community communicate through ARPANET (exchanging e-mails, file transfers..)
- 1989 Tim Berners-Lee creates the World Wide Web (networked application that links users and services distributed across computers around the world. Users can view pages, search for information, send and receive e-mails, initiate business transactions etc. Tim Berners Lee proposed linking info presented on various machines at CERN (European Laboratory for Particle Physics, Geneva). Other systems had been created to download files (FTP)from remote computers, search for information (Gopher), WAIS (Wide Area Information Servers) etc. Why did the Web succed? 1)SIMPLE INTERFACE 2) ENHANCED FEATURES 3) USE OF HYPERTEXT

Breve storia del Web

- 1991 First browser and server introduced
- 1993 The Web consisted of around 50 servers
- 1993 First release of Mosaic browser for windows (written by Marc Andreesen and Eric Bina). The web accounted for 1% of the traffic of the Internet
- Late 1990s → Web responsible for over 75% of the Internet traffic!! Hundreds millions of users; millions of web sites. Reasons for success: graphical user interface, ease of publishing new content
- Web used for e-business, business to business, interactions between users (chatrooms and games). Web increasingly used also to access private or proprietary data.
 - Web traffic: small number of Web pages accounts for the majority of web accesses
 - from text to images, animations, video files (increased storage needed at the Web servers, increased time for downloading, increased load on the network)
 - Dynamically generated content (e.g. search engines generates pointers based on keywords)
 - Web used for banking, e-business; important personal info (e.g. credit card info transferred) → security, authentication needed