

Reti di Elaboratori

Corso di Laurea in Informatica
Università degli Studi di Roma “La Sapienza”
Canale A-L e M-Z
Prof.ssa Chiara Petrioli

Parte di queste slide sono state prese dal materiale associato al libro
Computer Networking: A Top Down Approach , 5th edition.

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Francesco LoPresti, Un. di Roma Tor Vergata

Chapter 1: roadmap

1.1 What *is* the Internet?

1.2 Network edge

- end systems, access networks, links

1.3 Network core

- circuit switching, packet switching, network structure

1.4 Delay, loss and throughput in packet-switched networks

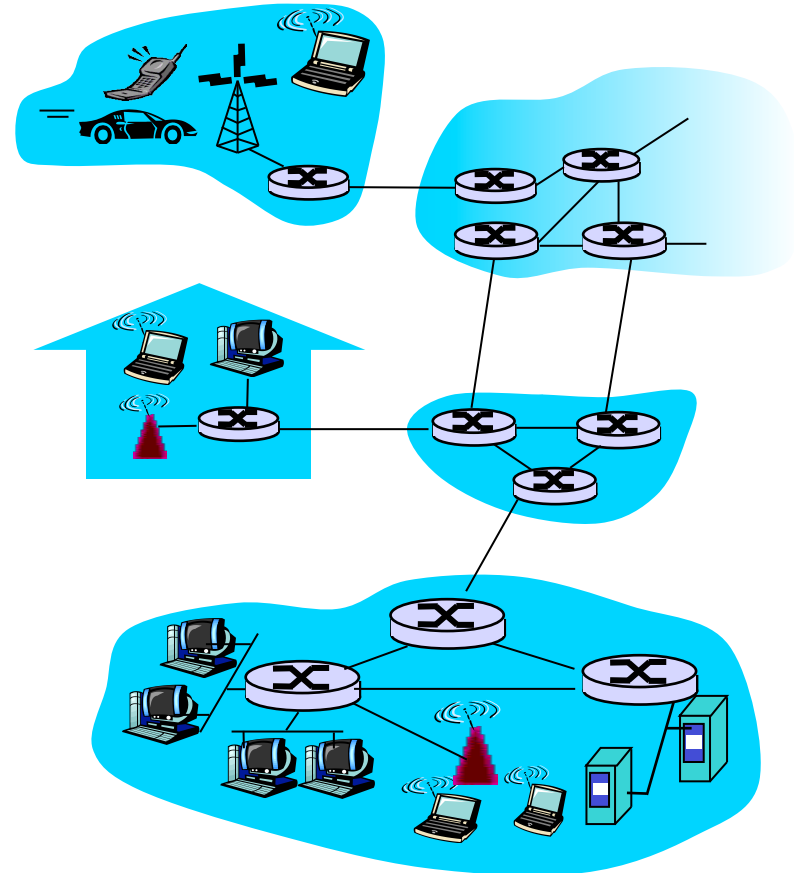
1.5 Protocol layers, service models

1.6 Networks under attack: security

1.7 History

A closer look at network structure:

- ❑ **network edge:**
applications and hosts
- ❑ **access networks,**
physical media: wired,
wireless
communication links
- ❑ **network core:**
 - ❖ interconnected routers
 - ❖ network of networks



The network edge:

□ end systems (hosts):

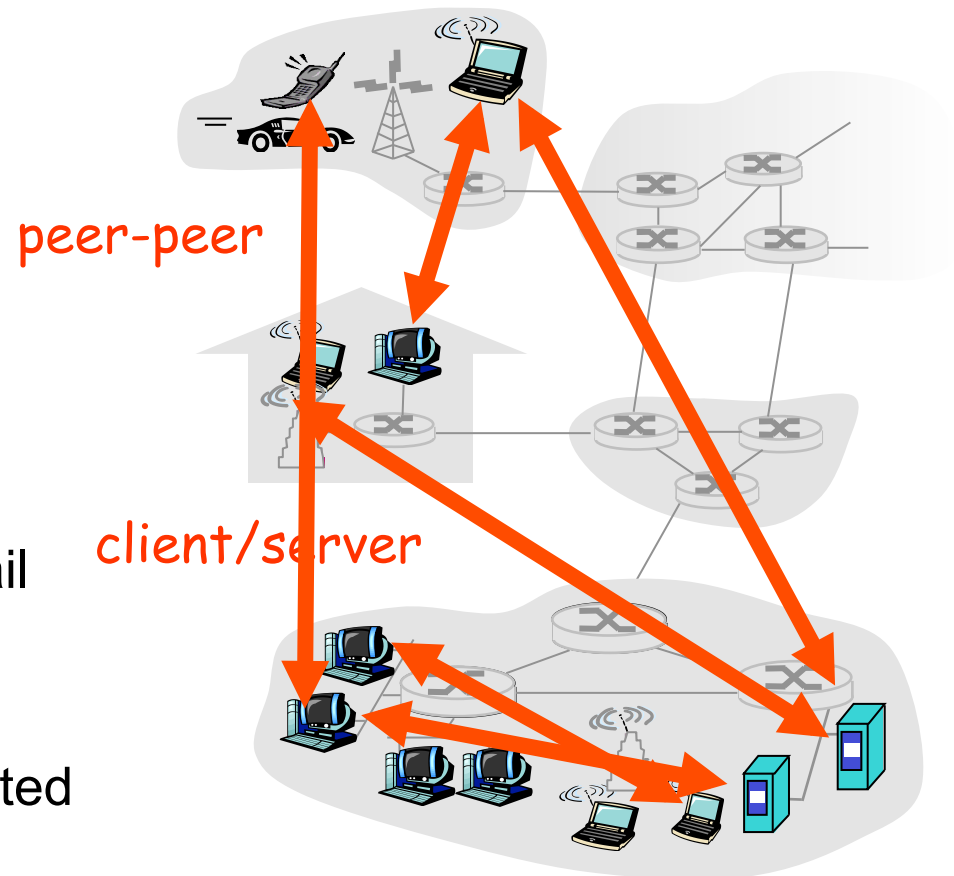
- run application programs
- e.g. Web, email
- at “edge of network”

□ client/server model

- ❖ client host requests, receives service from always-on server
- ❖ e.g. Web browser/server; email client/server

□ peer-peer model:

- ❖ minimal (or no) use of dedicated servers
- ❖ e.g. Skype, BitTorrent



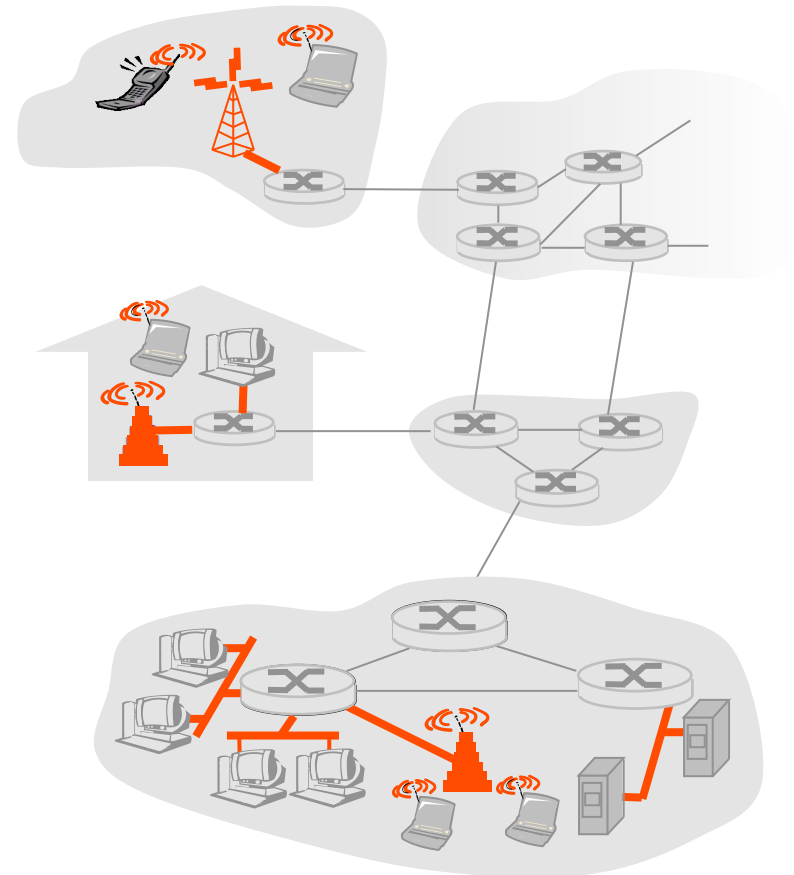
Access networks and physical media

Q: How to connect end systems to edge router?

- ❑ residential access nets
- ❑ institutional access networks (school, company)
- ❑ mobile access networks

Keep in mind:

- ❑ bandwidth (bits per second) of access network?
- ❑ shared or dedicated?
- ❑ reliable/unreliable (bit error rates)

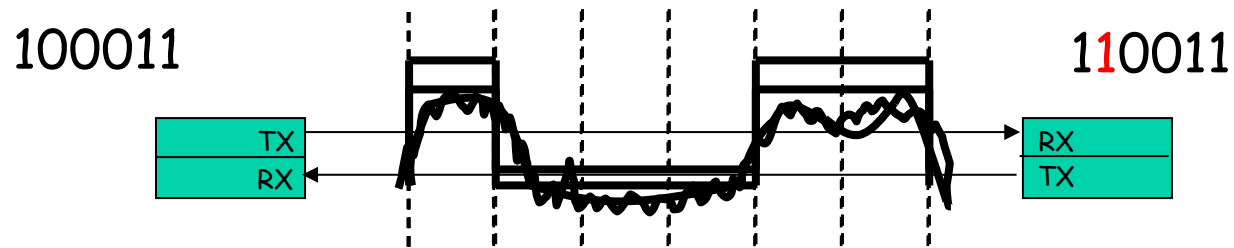


Transmission across a physical link



- ❑ **Bits:** propagate between transmitter and receiver
- ❑ **physical link:** what lies between transmitter & receiver
- ❑ **guided media:**
 - signals propagate in solid media: copper, fiber, coax
- ❑ **unguided media:**
 - signals propagate freely, e.g., radio

Transmission across a physical link

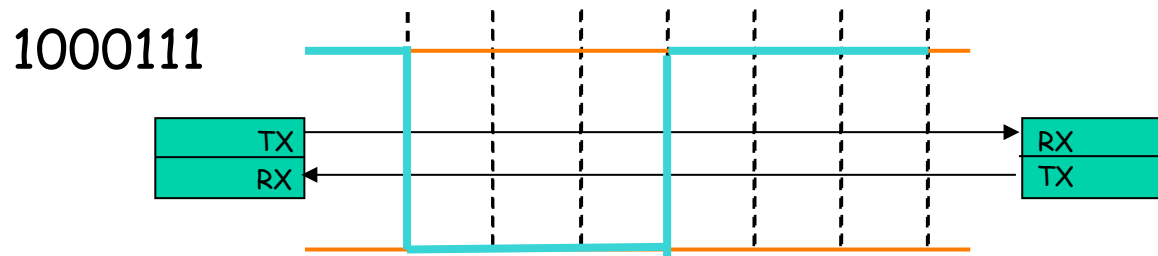


- ❑ Bit sequence modulates a suitable waveform which is sent across the link
 - How and which depends on the medium
- ❑ As the signal travels it experiences
 - **Attenuation** (absorption)
 - **Distortion** (limited bandwidth (frequency))
 - **Noise** (interference, thermal noise)
 - Influenced by medium, bit rate and distance
- ❑ Received sequence may be incorrect!!!

Codifica NRZ

❑ Codifica NRZ (Non Return to Zero)

Ogni bit ha associato un valore stabile per la sua intera durata (1: High; 0: Low)

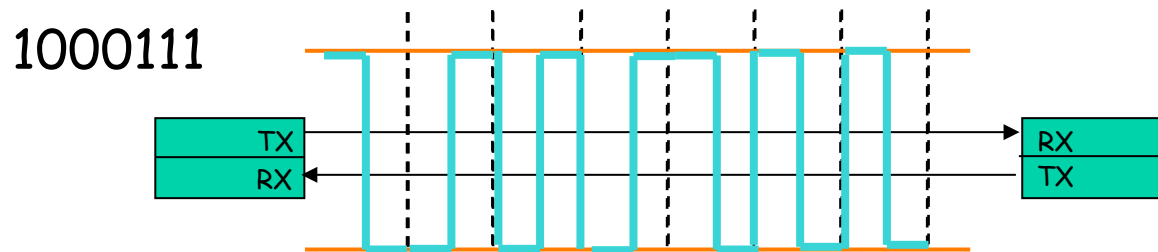


Problemi di sincronizzazione del ricevitore (nessuna transizione nel caso di sequenze di zeri o di uni) → NRZ 5B6B o 4B5B

Codifica Manchester

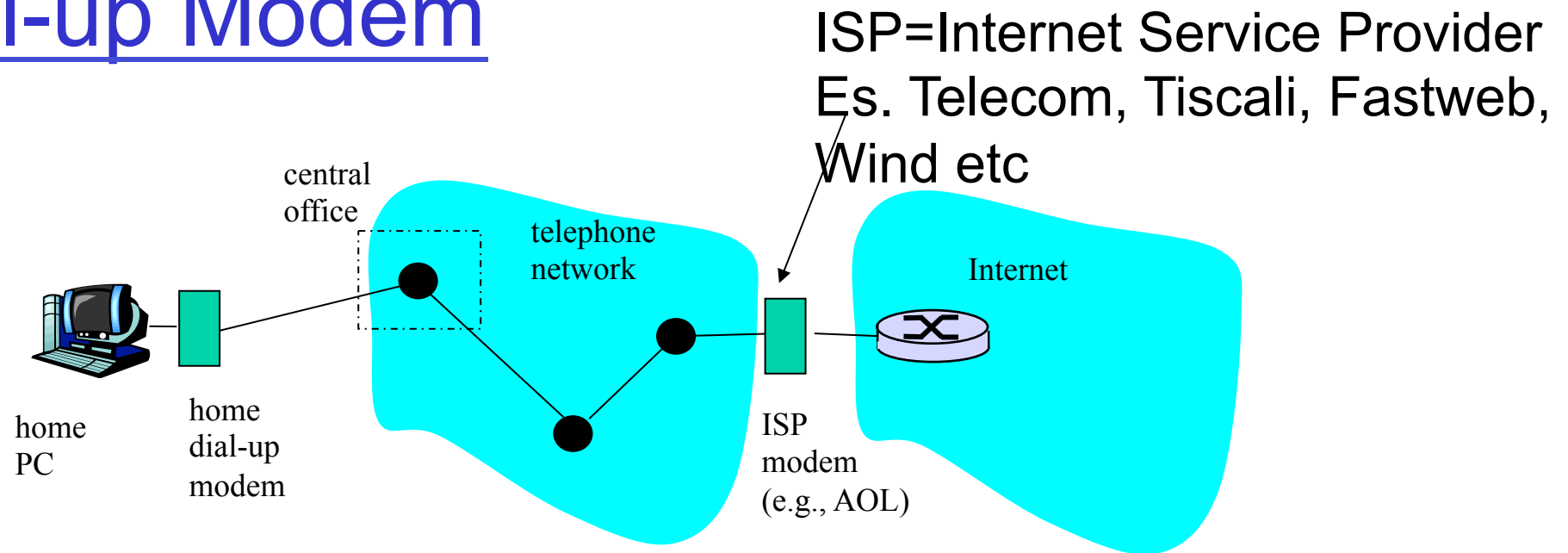
❑ Codifica Manchester

Una transizione basso-alto (codifica dello zero) o alto-basso (codifica del valore uno) in corrispondenza di ogni bit



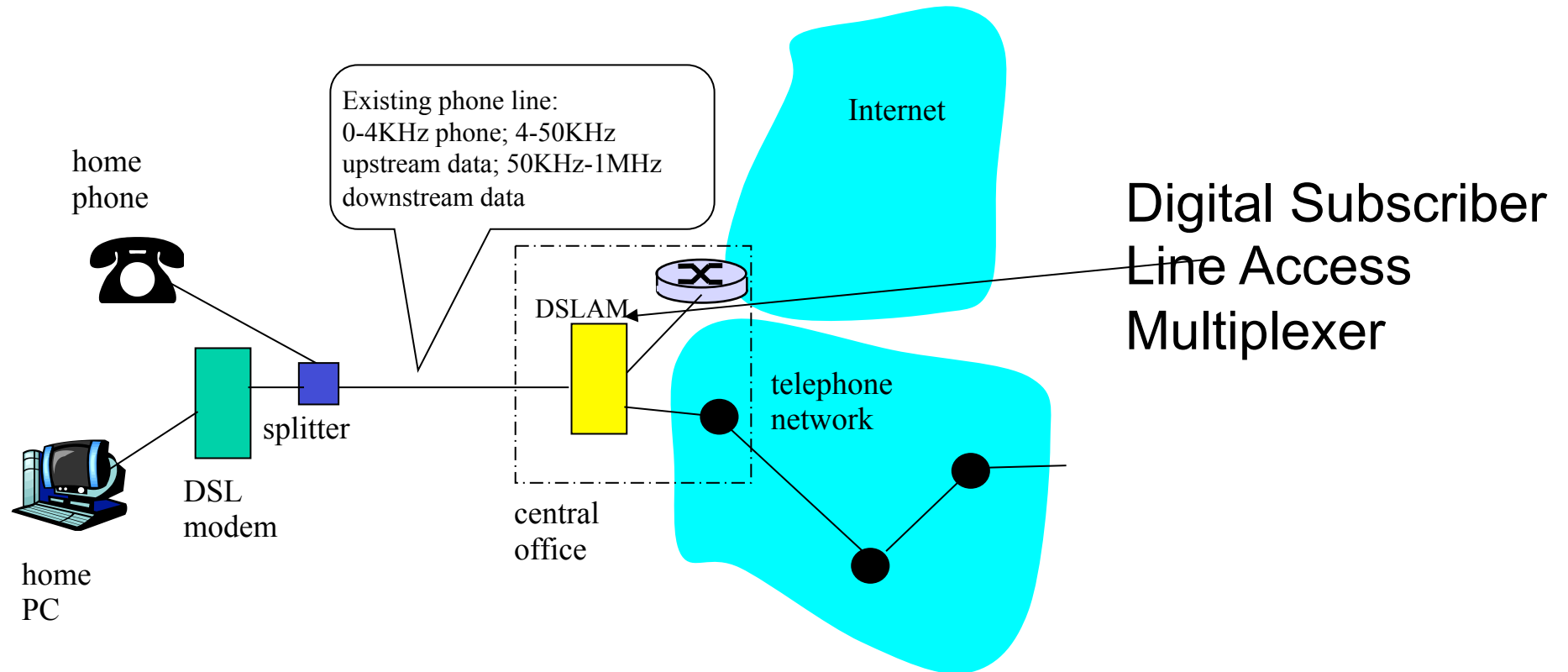
Usato in Ethernet 10Mbps e Token Ring

Dial-up Modem



- ❖ Uses existing telephony infrastructure
 - ❖ Home is connected to **central office**
- ❖ up to 56Kbps direct access to router (often less)
- ❖ Can't surf and phone at same time: not **"always on"**

Digital Subscriber Line (DSL)

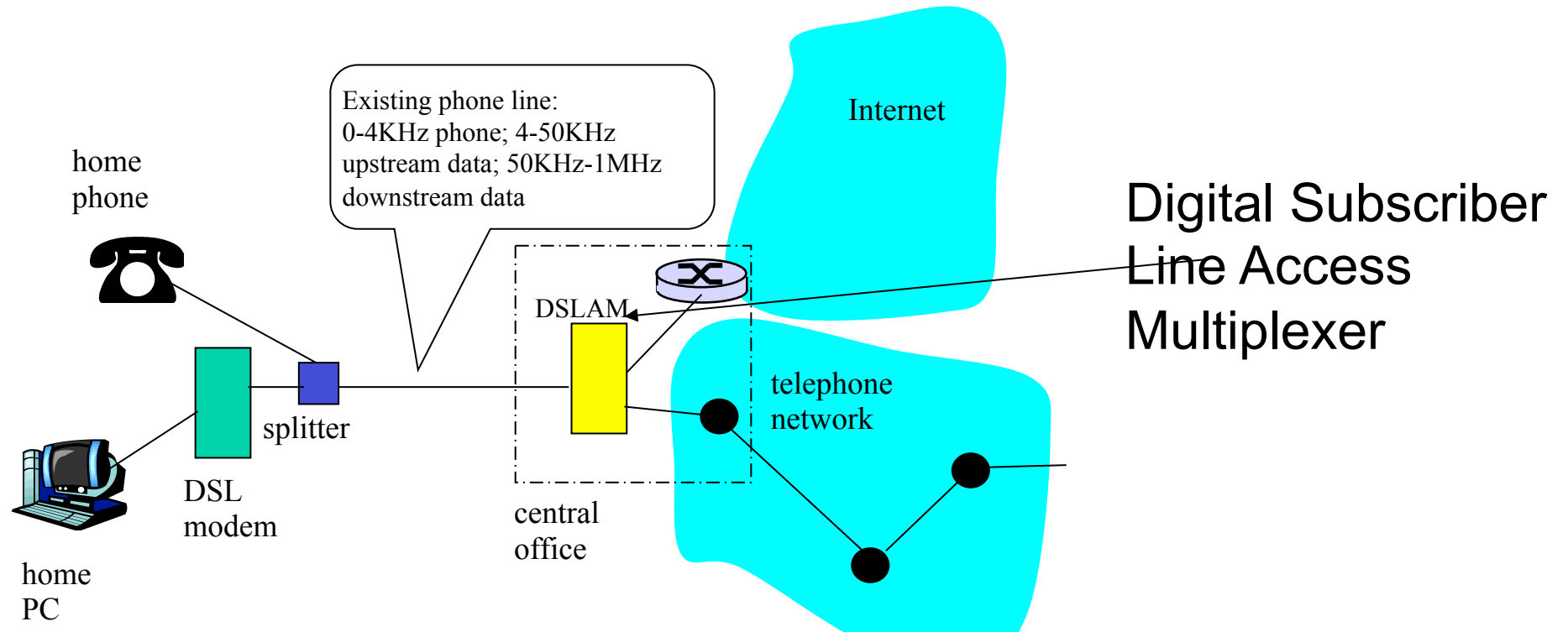


- ❖ Also uses existing telephone infrastructure
- ❖ up to 1 Mbps upstream (typically < 256 kbps)
- ❖ up to 8 Mbps downstream (typically < 1 Mbps)
- ❖ dedicated physical line to telephone central office

ADSL loops extender

- ❑ An **ADSL loop extender** or **ADSL repeater** is a device placed midway between the subscriber and central office by the telephone company to extend the distance and increase the channel capacity of their DSL connection.
- ❑ In some cases, service can now be established as far as 10 miles from the Central Office (factor of 2 improvement)

Digital Subscriber Line (DSL)

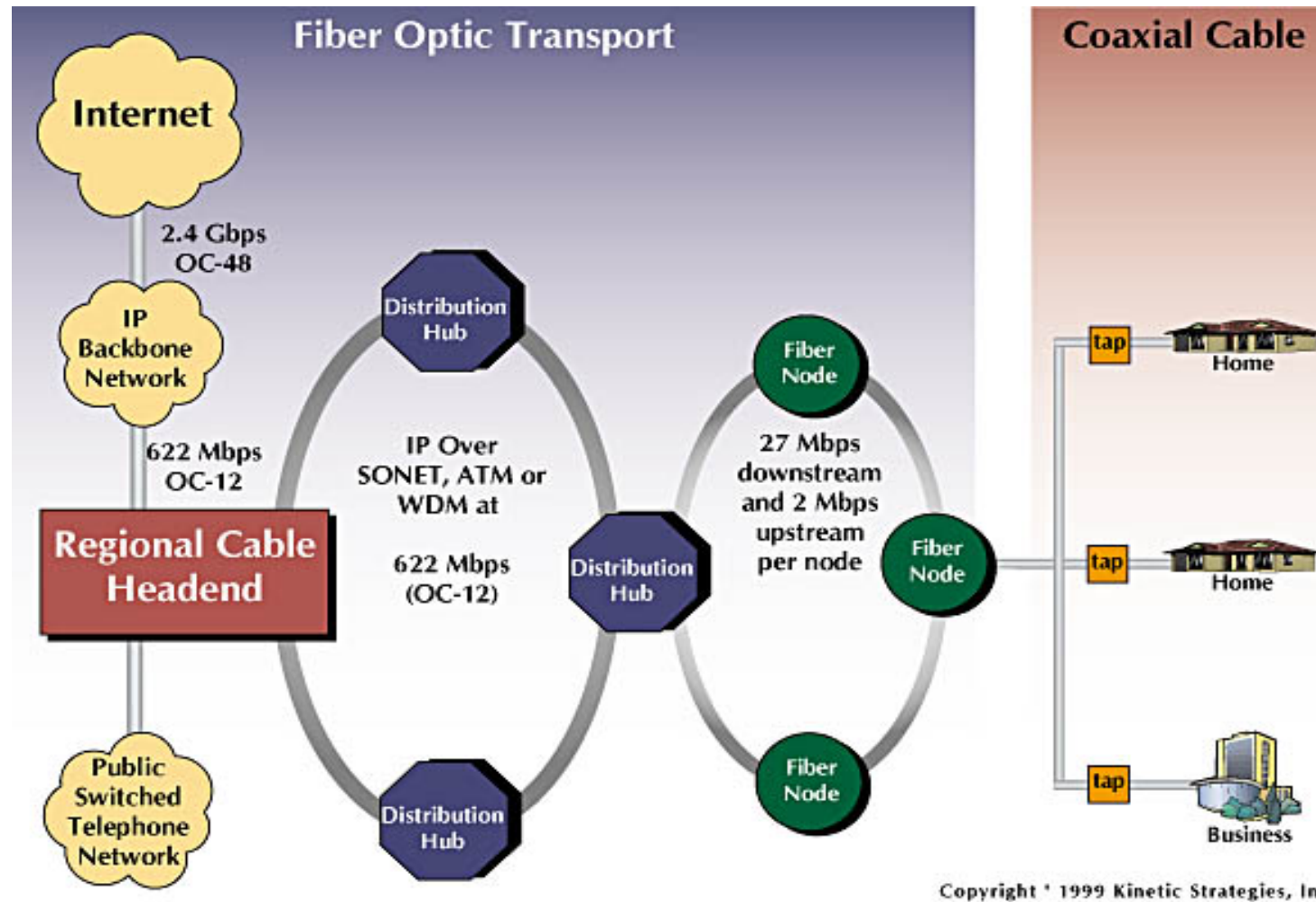


- ❖ Also uses existing phone lines
 - ❖ up to 1 Mbps
 - ❖ up to 8 Mbps
 - ❖ dedicated physical line
- Speed significantly increased in the last few years**
- technologies more robust to interference;
 - lower distance from DSM modem to DSLAM
- is expected to raise speed to 1Gbps by 2016)**

Residential access: cable modems

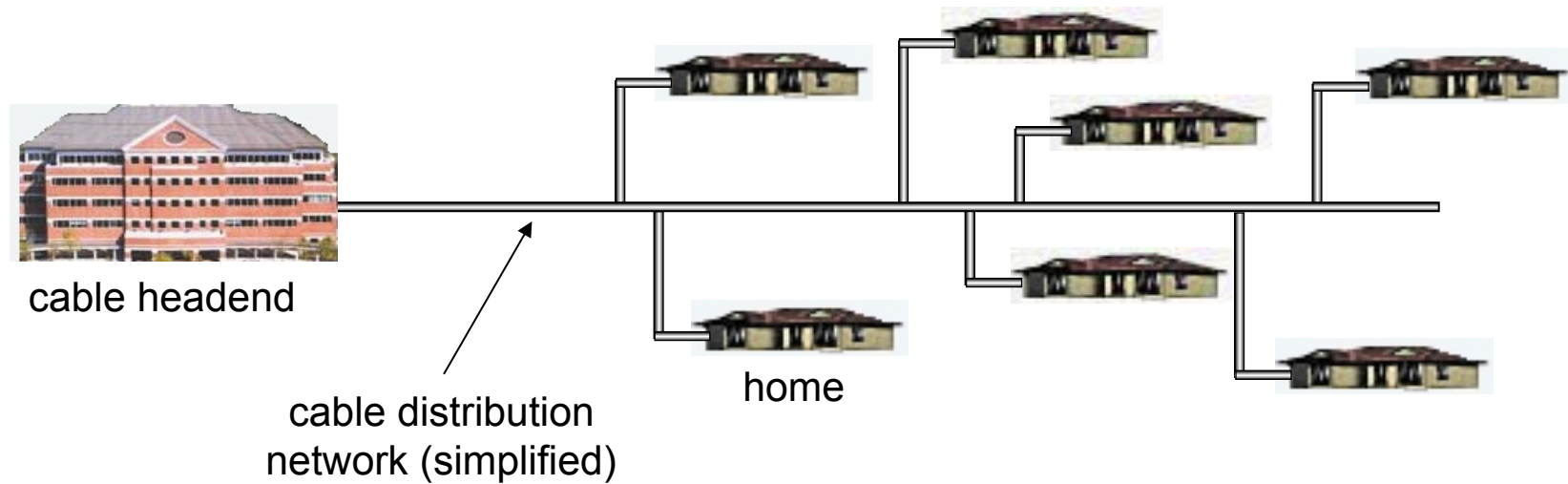
- ❑ Does not use telephone infrastructure
 - Instead uses cable TV infrastructure
- ❑ HFC: hybrid fiber coax
 - asymmetric: up to 30Mbps downstream, 2 Mbps upstream
- ❑ network of cable and fiber attaches homes to ISP router
 - homes share access to router
 - unlike DSL, which has dedicated access

Residential access: cable modems

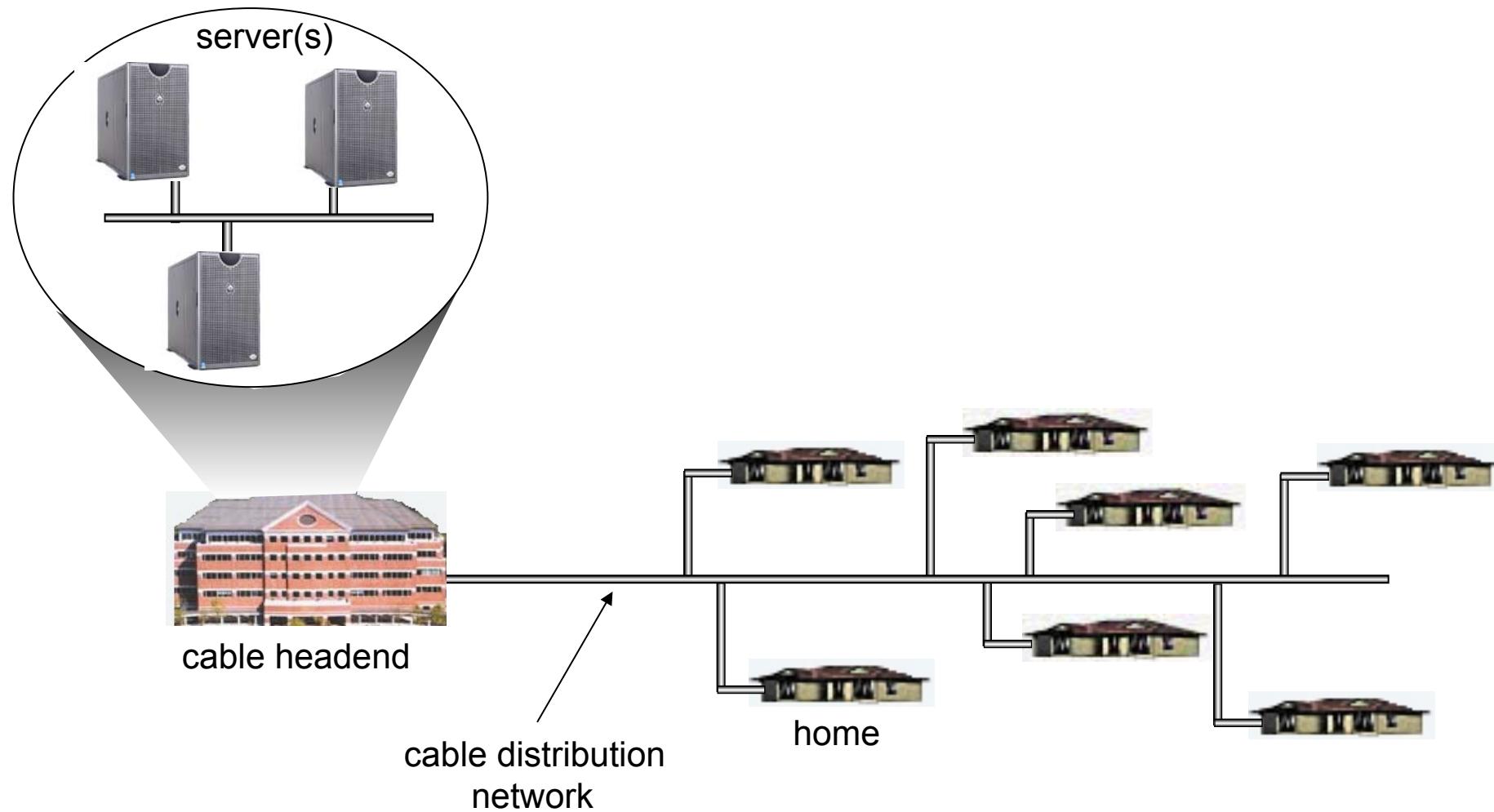


Cable Network Architecture: Overview

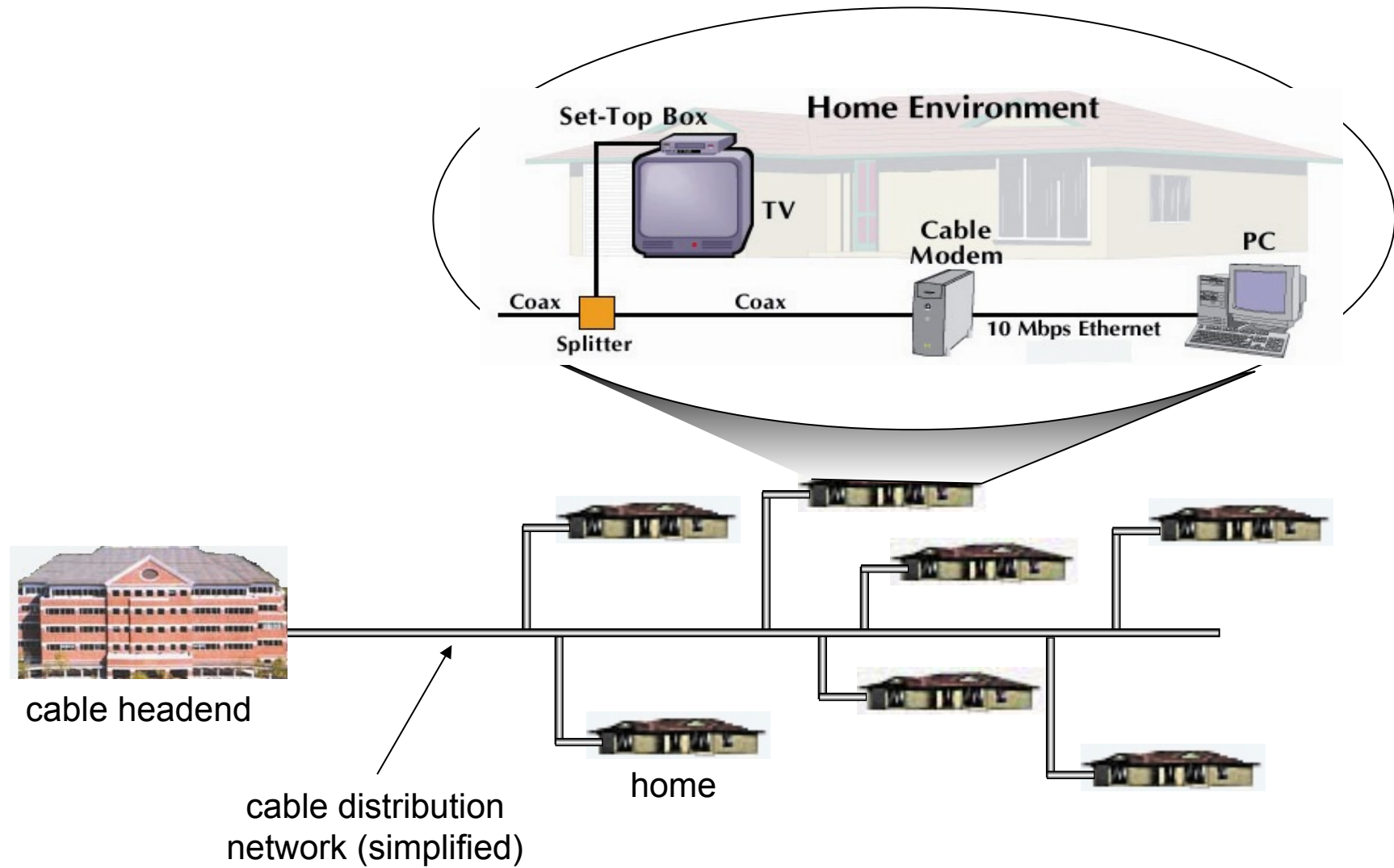
Typically 500 to 5,000 homes



Cable Network Architecture: Overview

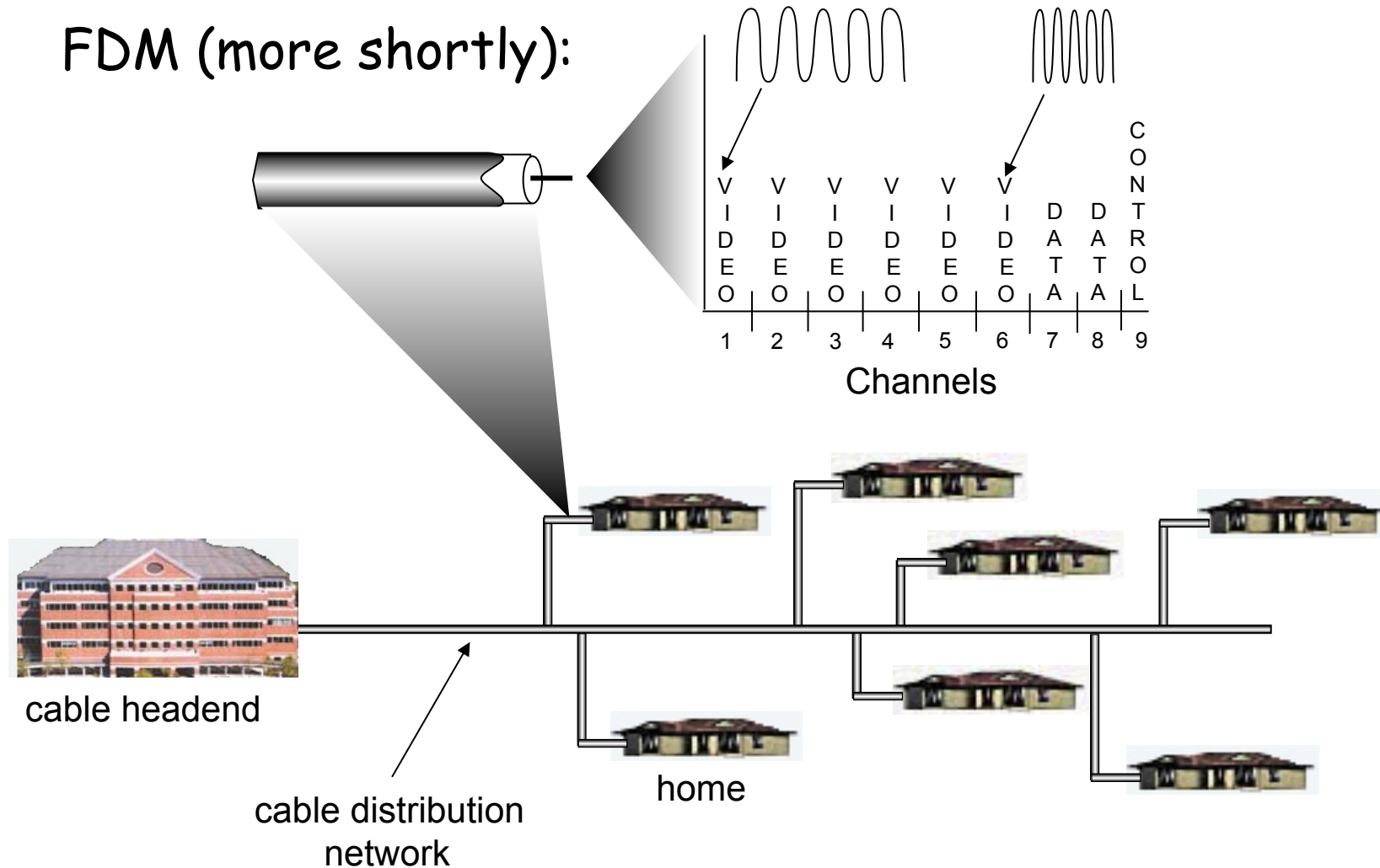


Cable Network Architecture: Overview

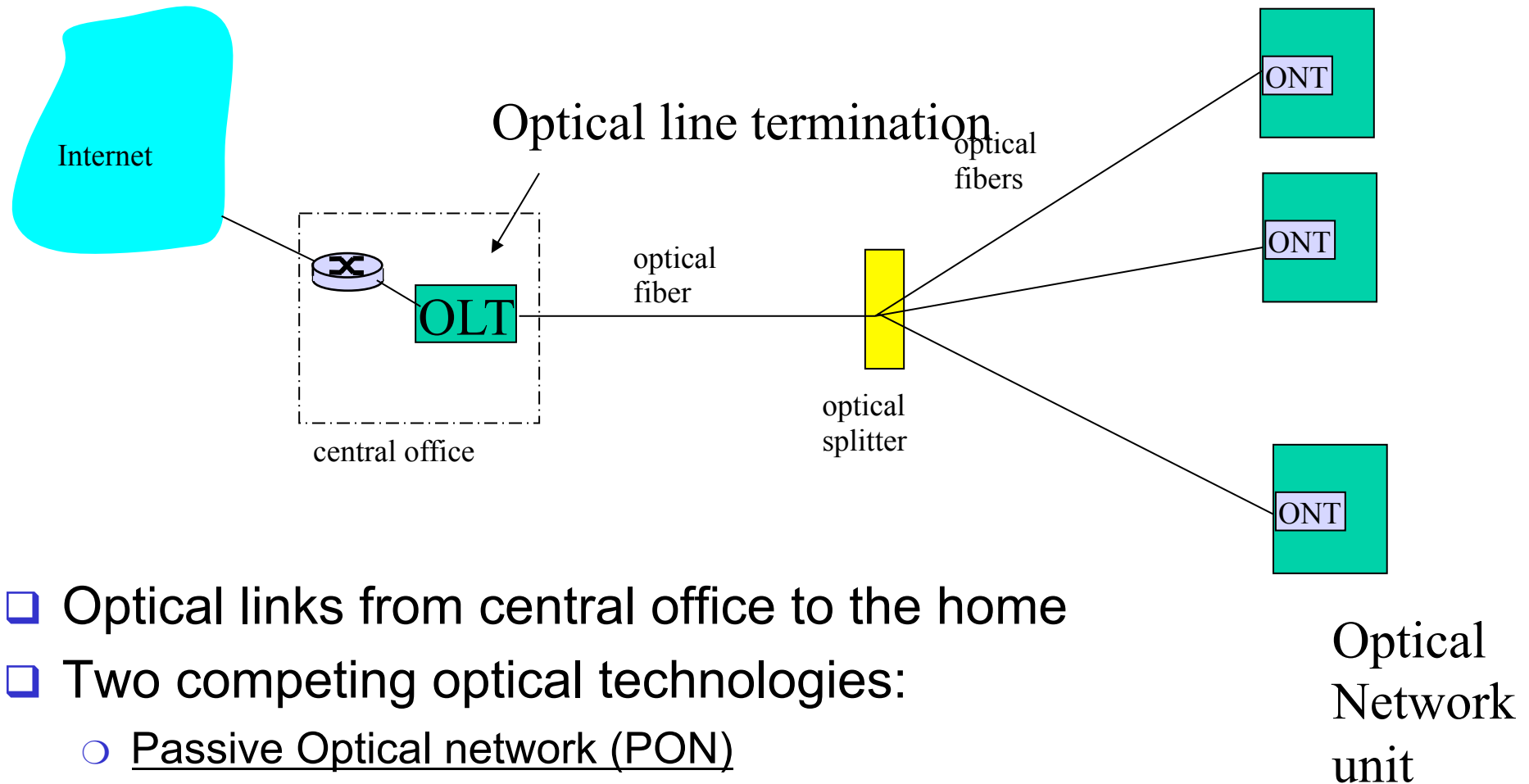


Cable Network Architecture: Overview

FDM (more shortly):

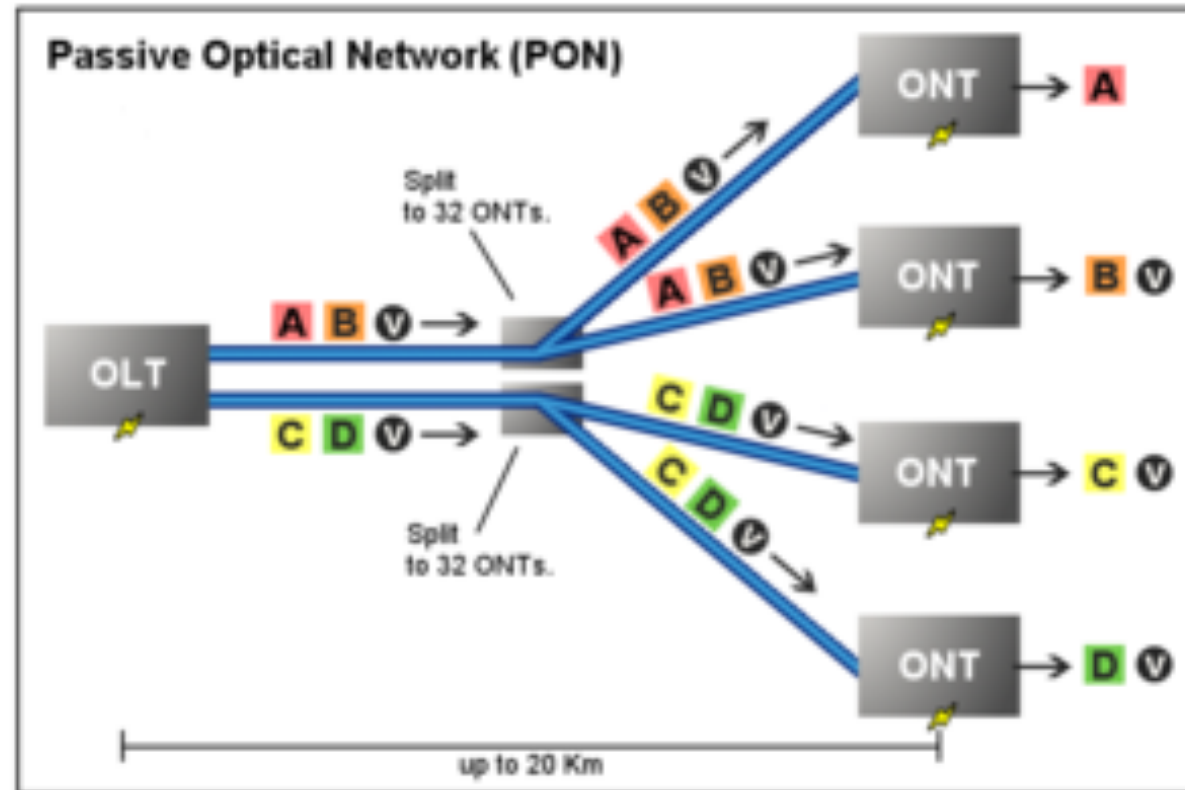
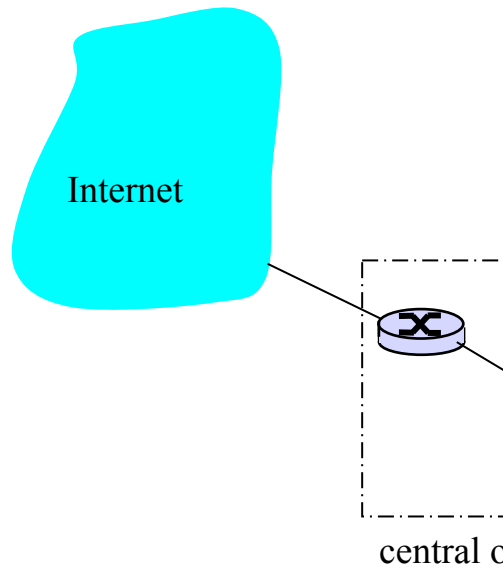


Fiber to the Home



- ❑ Optical links from central office to the home
- ❑ Two competing optical technologies:
 - Passive Optical network (PON)
 - Active Optical Network (PAN)
- ❑ Much higher Internet rates; fiber also carries television and phone services

Fiber to the Home



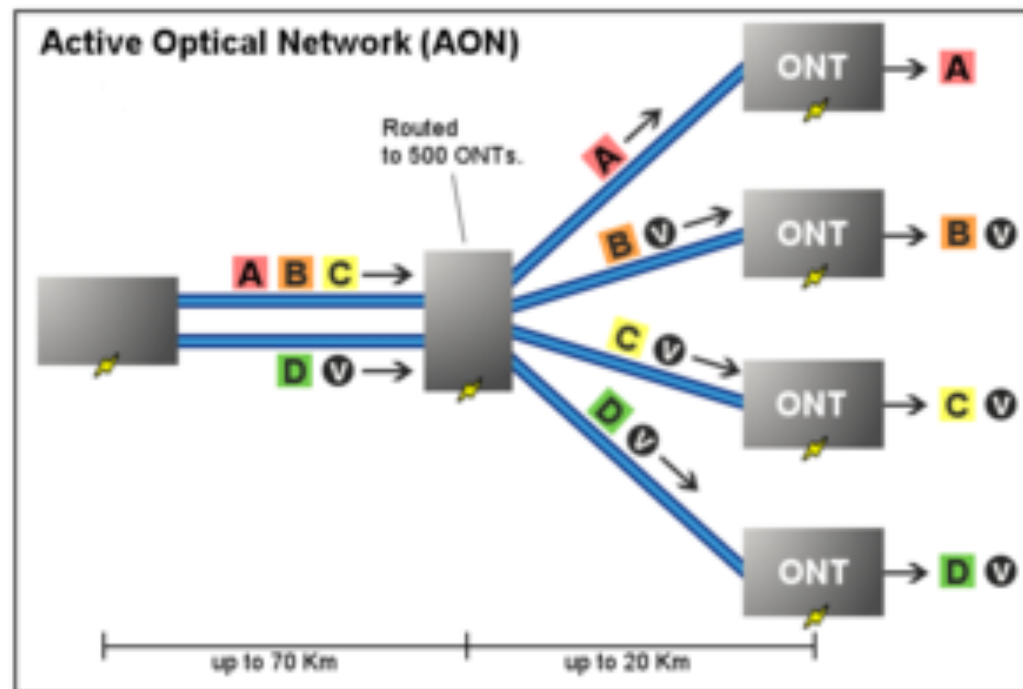
Key: **A** - Data or voice for a single customer. **V** - Video for multiple customers.

- ❑ Optical links from central office to the home
- ❑ Two competing technologies
 - Passive Optical network (PON)
 - Active Optical Network (PAN)
- ❑ Much higher Internet rates; fiber also carries television and phone services

central
network
unit

Active Optical Networks

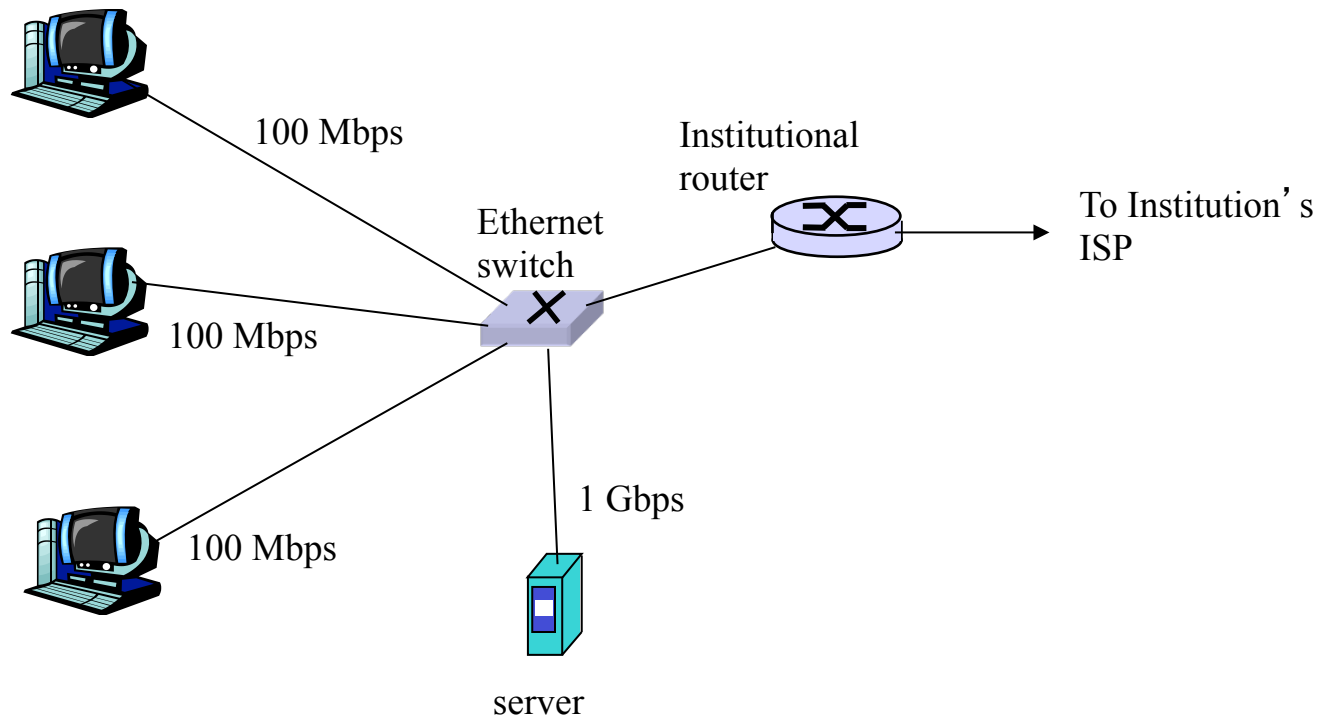
- ❑ An active optical system uses electrically powered switching equipment, such as a router or a switch aggregator, to manage signal distribution and direct signals to specific customers.
- ❑ In such a system, a customer may have a dedicated fiber running to his or her house.



Active vs Passive Optical Networks

- ❑ Passive optical networks, or PONs, have some distinct advantages.
 - They're efficient, in that each fiber optic strand can serve up to 32 users
 - PONs have a low building cost relative to active optical networks along with lower maintenance costs. In active optical networks one aggregator is required every 48 subscribers.
- ❑ Passive optical networks also have some disadvantages.
 - They have less range than an active optical network.
 - PONs also make it difficult to isolate a failure when they occur.
 - Because the bandwidth in a PON is not dedicated to individual subscribers, data transmission speed may slow down during peak usage times.

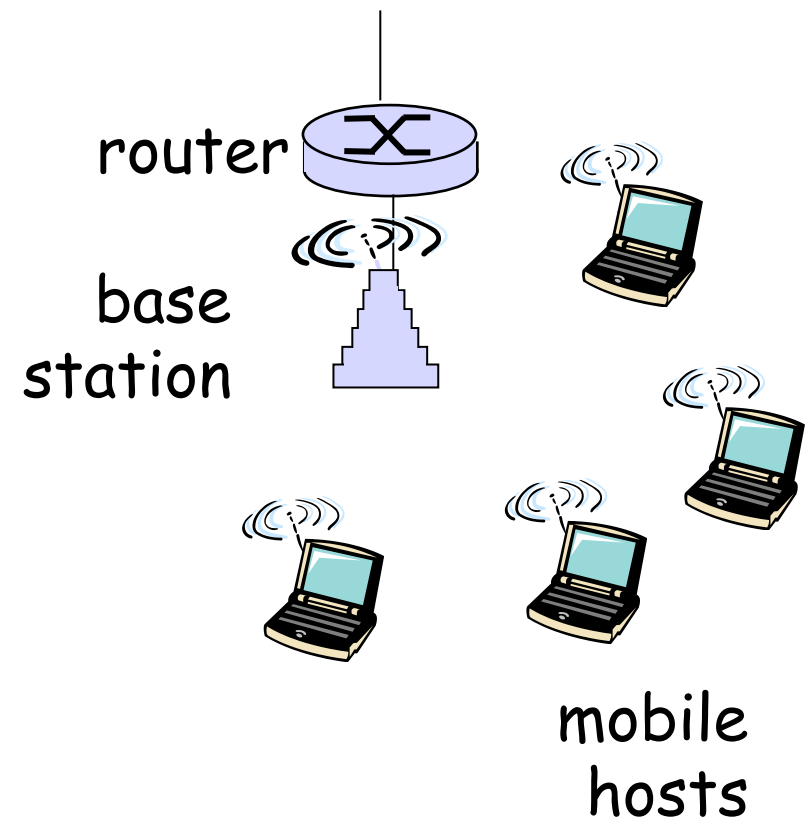
Ethernet Internet access



- ❑ Typically used in companies, universities, etc
- ❑ 10 Mbs, 100Mbps, 1Gbps, 10Gbps Ethernet
- ❑ Today, end systems typically connect into Ethernet switch

Wireless access networks

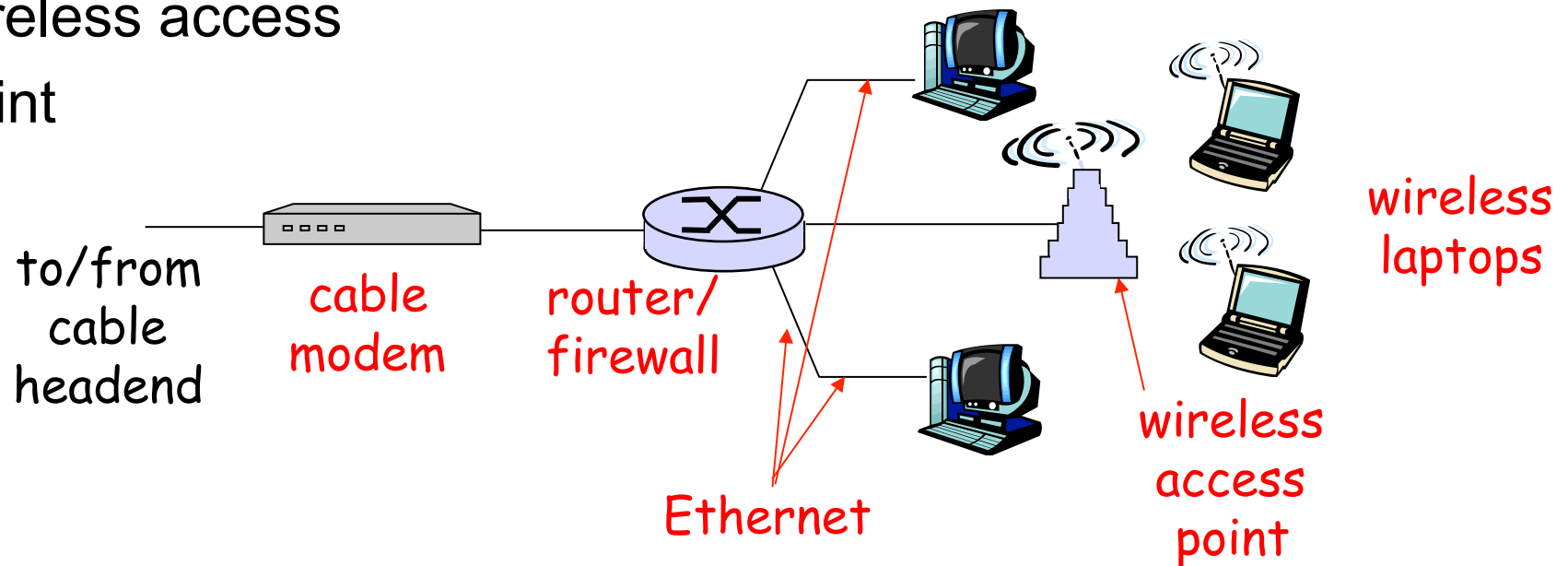
- ❑ shared *wireless* access network connects end system to router
 - via base station aka “access point”
- ❑ **wireless LANs:**
 - 802.11b/g (WiFi): 11 or 54 Mbps
- ❑ **wider-area wireless access**
 - provided by telco operator
 - ~1Mbps over cellular system (EVDO, HSDPA), several tens Mbps LTE
 - WiMAX (10' s Mbps) over wide area
 - Next to come: 5G systems



Home networks

Typical home network components:

- ❑ DSL or cable modem
 - ❑ router/firewall/NAT
 - ❑ Ethernet
 - ❑ wireless access point
- point

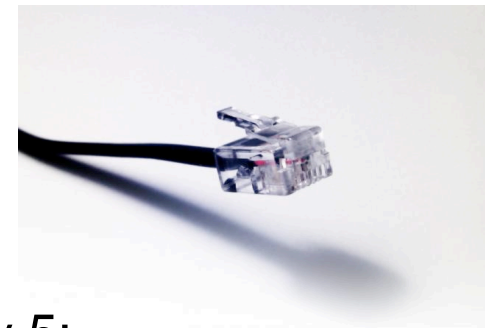


Physical Media

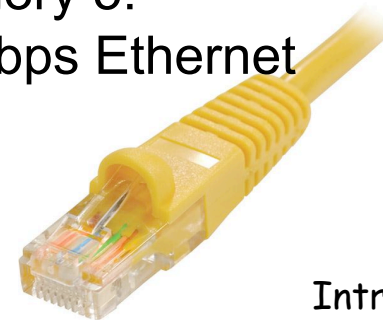
- ❑ **Bit:** propagates between transmitter/rcvr pairs
- ❑ **physical link:** what lies between transmitter & receiver
- ❑ **guided media:**
 - signals propagate in solid media: copper, fiber, coax
- ❑ **unguided media:**
 - signals propagate freely, e.g., radio

Twisted Pair (TP)

- ❑ two insulated copper wires
 - Category 3: traditional phone wires, 10 Mbps Ethernet



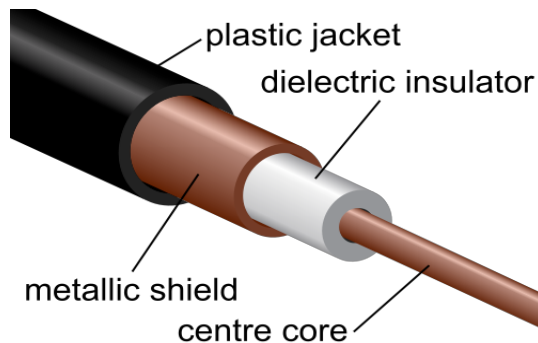
- Category 5: 100Mbps Ethernet



Physical Media: coax, fiber

Coaxial cable:

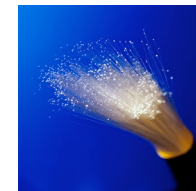
- ❑ two concentric copper conductors
- ❑ bidirectional
- ❑ baseband:
 - single channel on cable
 - legacy Ethernet
- ❑ broadband:
 - multiple channels on cable
 - HFC



http://commons.wikimedia.org/wiki/File:Coaxial_cable_cutaway.svg

Fiber optic cable:

- ❑ glass fiber carrying light pulses, each pulse a bit
- ❑ high-speed operation:
 - ❖ high-speed point-to-point transmission (e.g., 10' s-100' s Gps, but experimented up to tens of terabps)
- ❑ low error rate: repeaters spaced far apart ; immune to electromagnetic noise



http://www.macmynd.com/storage/misc-pics/fiber_optic_cable.jpg

Physical media: radio

- ❑ signal carried in electromagnetic spectrum
- ❑ no physical “wire”
- ❑ bidirectional
- ❑ propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

Radio link types:

- ❑ **terrestrial microwave**
 - ❖ e.g. up to 45 Mbps channels
- ❑ **LAN** (e.g., Wifi)
 - ❖ 11Mbps, 54 Mbps
- ❑ **wide-area** (e.g., cellular)
 - ❖ 3G cellular: ~ 1 Mbps
- ❑ **satellite**
 - ❖ Kbps to 45Mbps channel (or multiple smaller channels)
 - ❖ 270 msec end-end delay
 - ❖ geosynchronous versus low altitude
 - (500 Km dalla superficie terrestre, servono costellazioni di satelliti)

Physical media performance evolution (update: 2014) –On the move

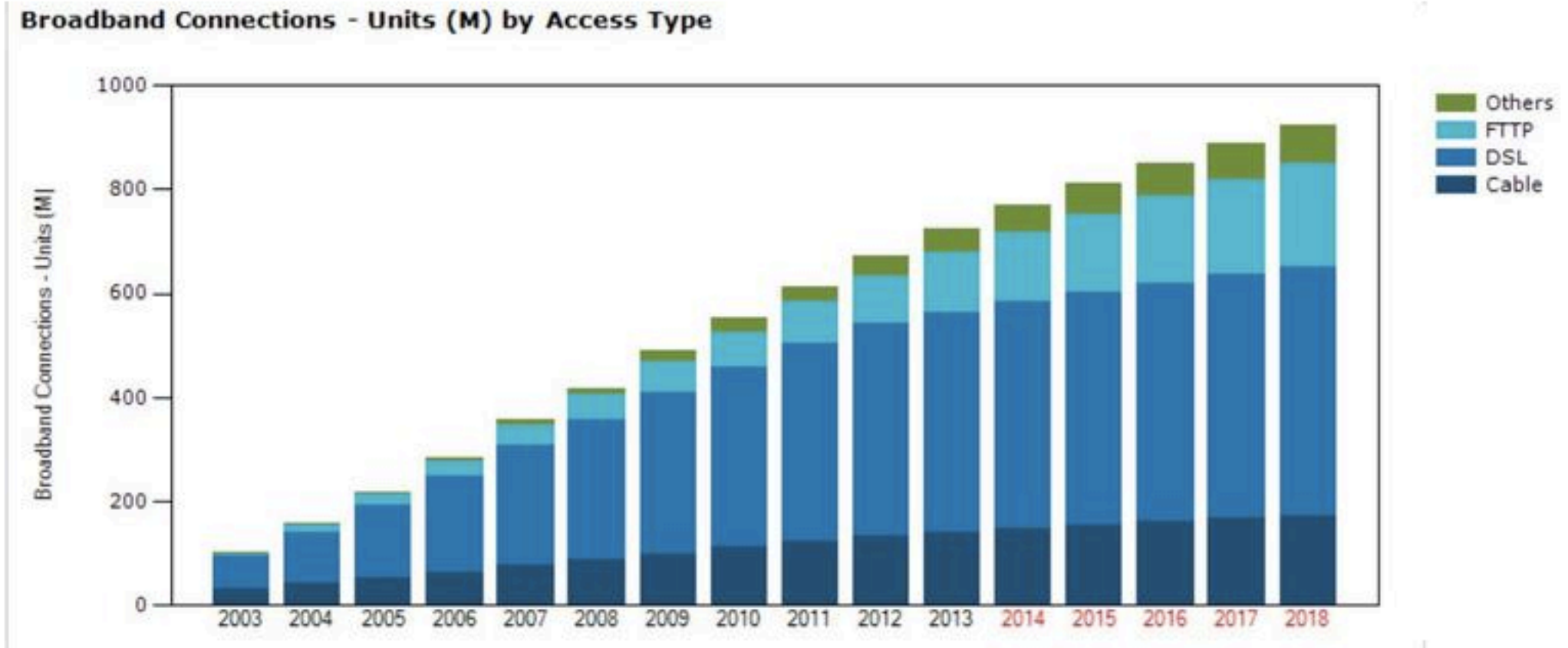
Generation		Technology	Maximum Download Speed	Typical Download Speed
2G	G	GPRS	0.1Mbit/s	<0.1Mbit/s
	E	EDGE	0.3Mbit/s	0.1Mbit/s
3G	3G	3G (Basic)	0.3Mbit/s	0.1Mbit/s
	H	HSPA	7.2Mbit/s	1.5Mbit/s
	H+	HSPA+	21Mbit/s	4Mbit/s
	H+	DC-HSPA+	42Mbit/s	8Mbit/s
4G	4G_U	LTE	100Mbit/s	15Mbit/s

Physical media performance evolution (update: 2014) –Access technologies

WiFi, Ethernet, Fiber to the “home”, DSL...Maximum current speeds or technologies tested to enter the market within a couple of years

- ❑ DSL (G.Fast technology) 1Gbps
 - By 2016
 - Combined with fiber; access to broadband network within 50m to reach such speeds
- ❑ Ethernet: 25Gbps (40Gbps under standardization). With more lines: currently 100Gbps, standards towards 400Gbps
- ❑ WiFi IEEE 802.11ac Up to 1Gbps to come
- ❑ Fiber
 - Technologies tested up to few tens of terabps
 - 1Gbps per home more than enough (current threshold per user satisfaction >10Mbps)
- ❑ Cellular systems evolution
 - Tens-hundred of Mbps

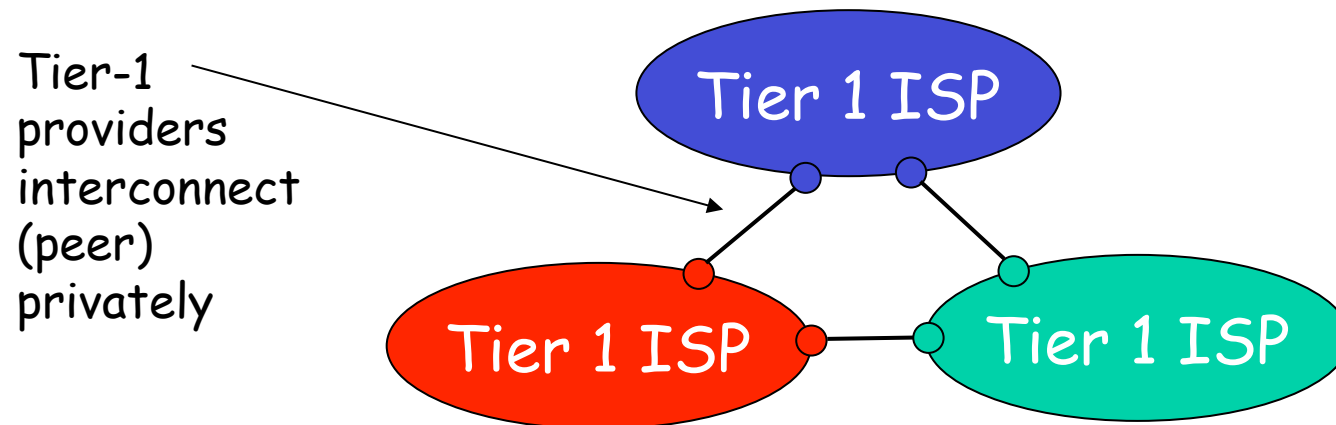
Physical media performance evolution (update: 2014)—different types of media



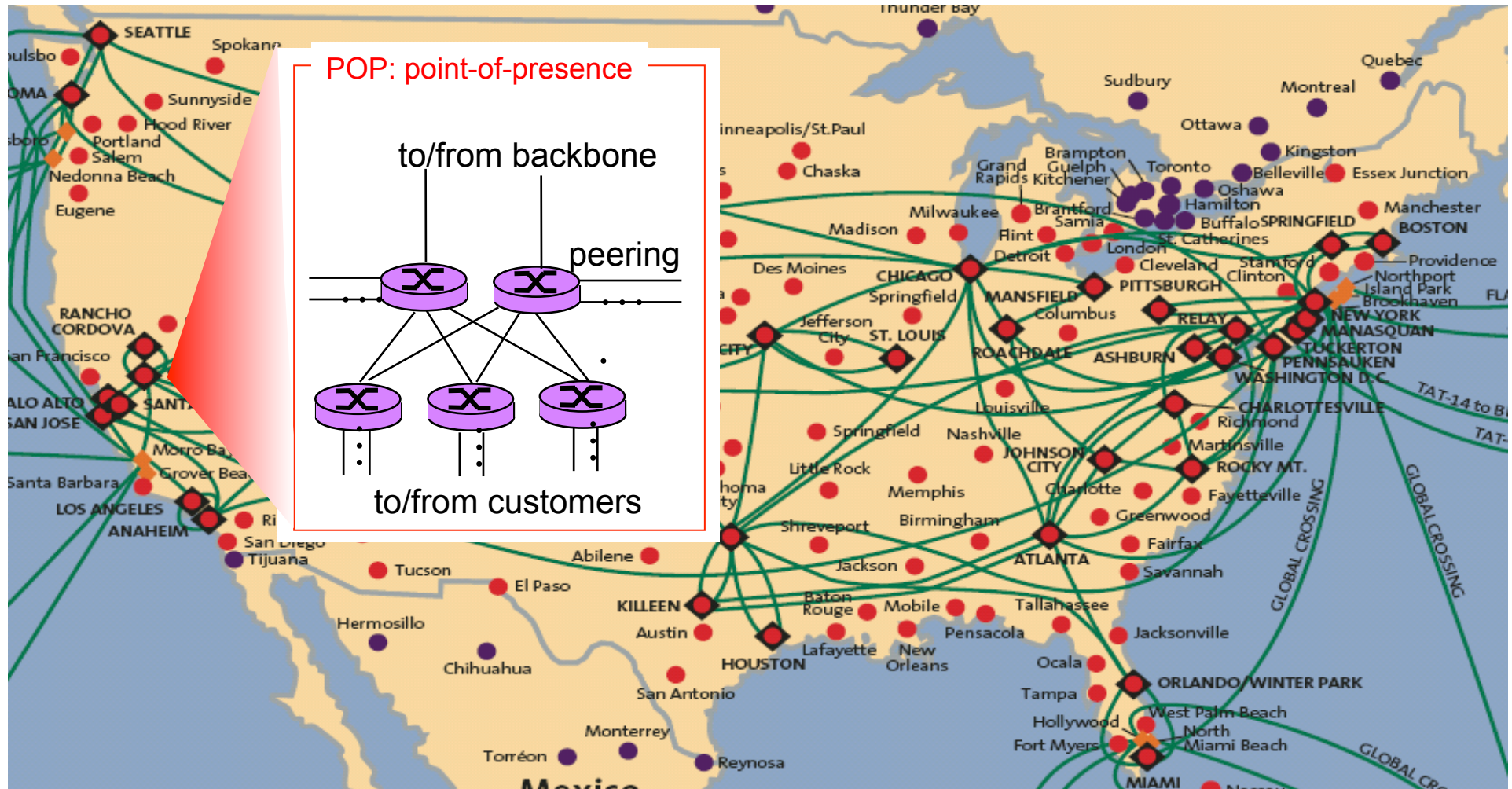
DSL is the most widely used broadband connection technology, and it's growing, but fiber-optic links are growing faster.

Internet structure: network of networks

- roughly hierarchical
- at center: “tier-1” ISPs (e.g., Verizon, Sprint, AT&T, Cable and Wireless), national/international coverage
 - treat each other as equals



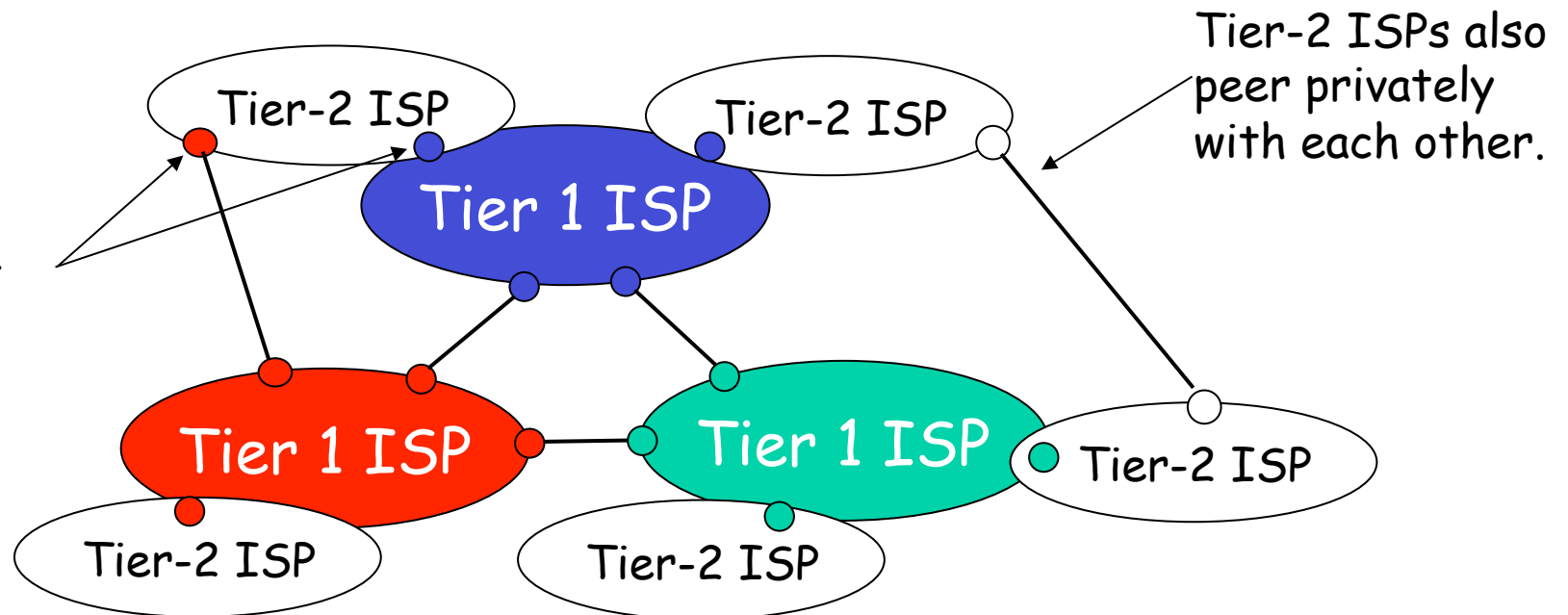
Tier-1 ISP: e.g., Sprint



Internet structure: network of networks

- “Tier-2” ISPs: smaller (often regional) ISPs
 - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

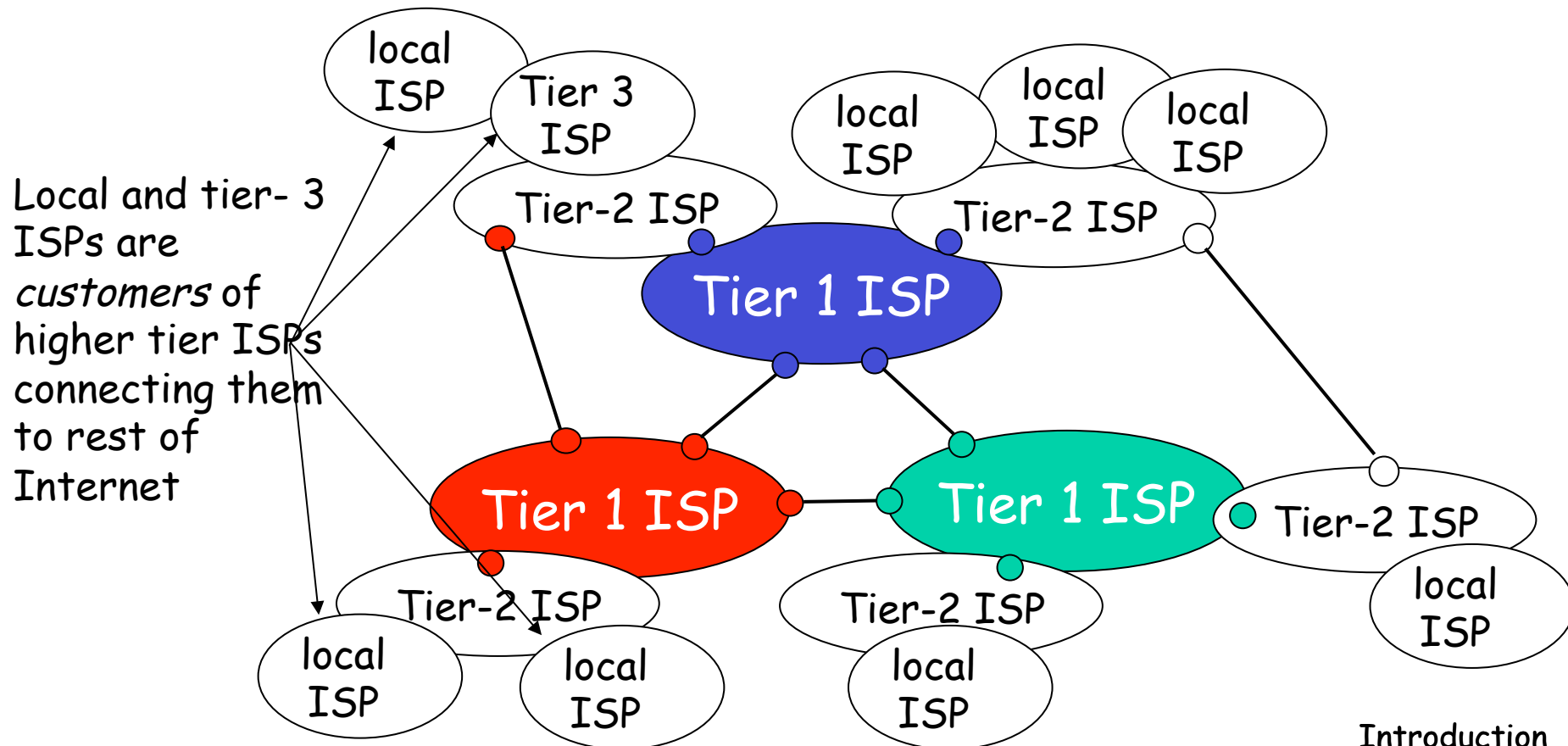
Tier-2 ISP pays tier-1 ISP for connectivity to rest of Internet
□ tier-2 ISP is customer of tier-1 provider



Internet structure: network of networks

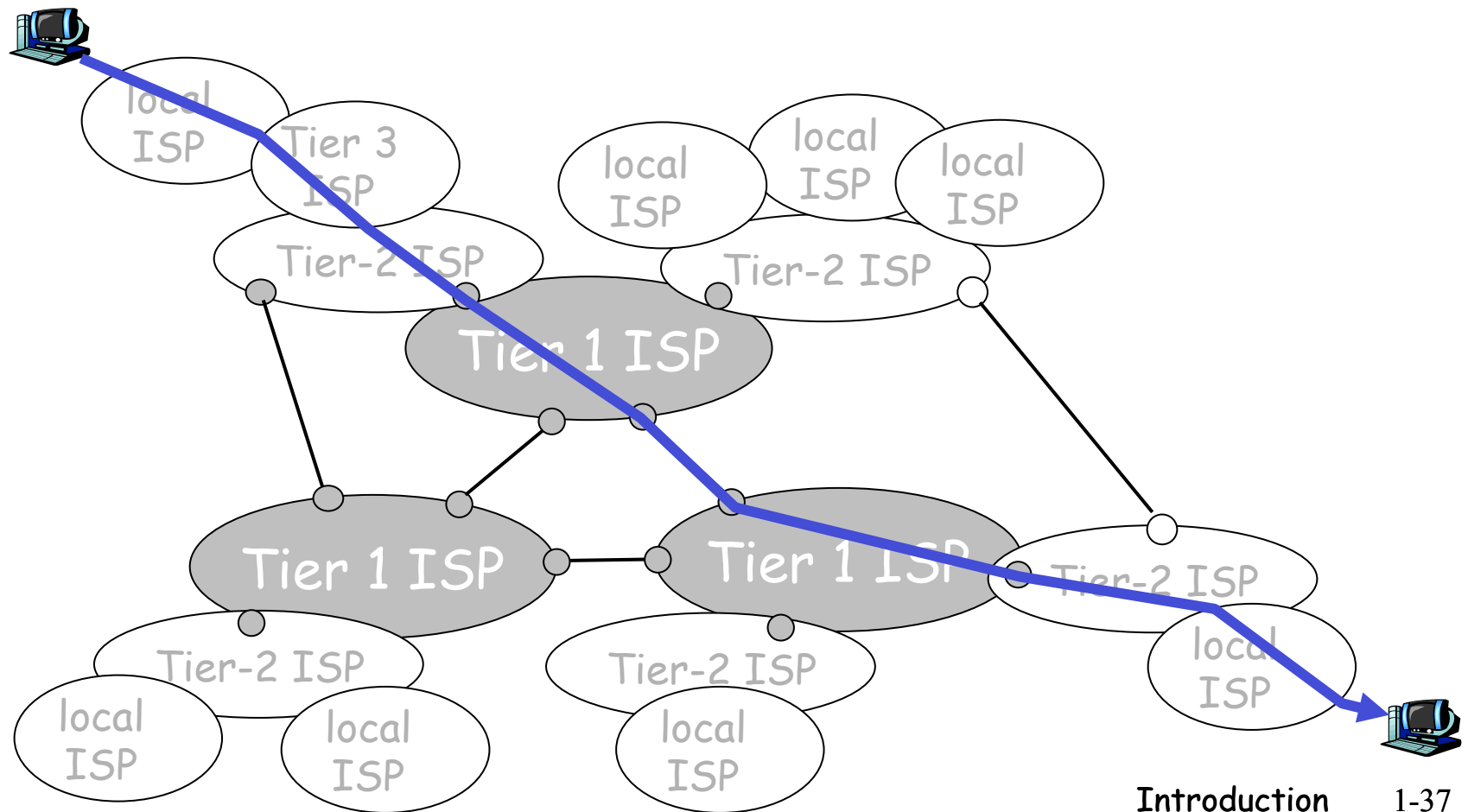
□ “Tier-3” ISPs and local ISPs

- last hop (“access”) network (closest to end systems)



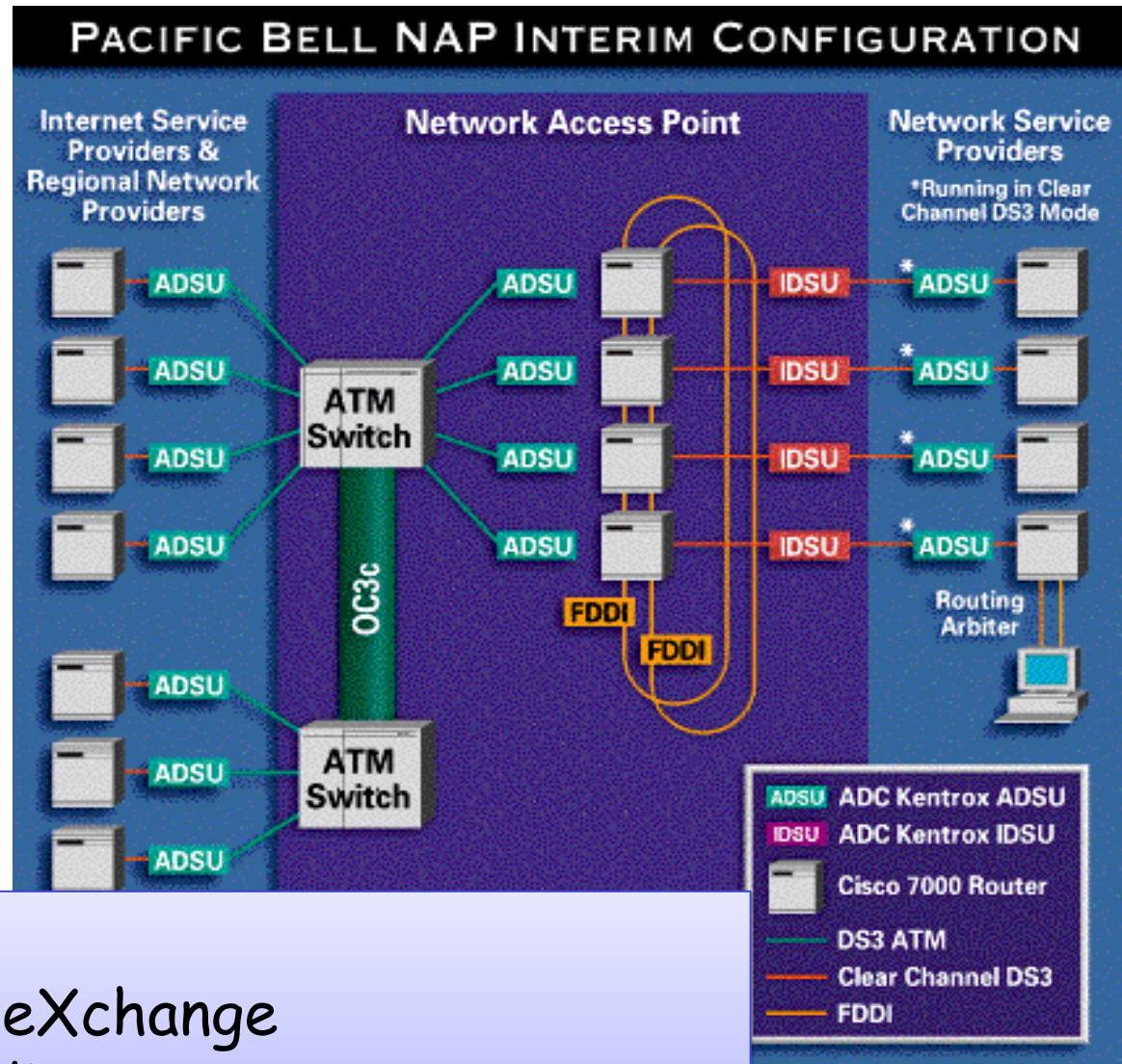
Internet structure: network of networks

- a packet passes through many networks!



A NAP: just another router...?

Pacific Bell S. Francisco NAP



In Italia:

- MIX Milan Internet eXchange
- NaMeX Nautilus Mediterranean Exchange Point

Chapter 1: roadmap

1.1 What *is* the Internet?

1.2 Network edge

1.3 Network core

1.4 Network access and physical media

1.5 Internet structure and ISPs

1.6 Delay & loss in packet-switched networks

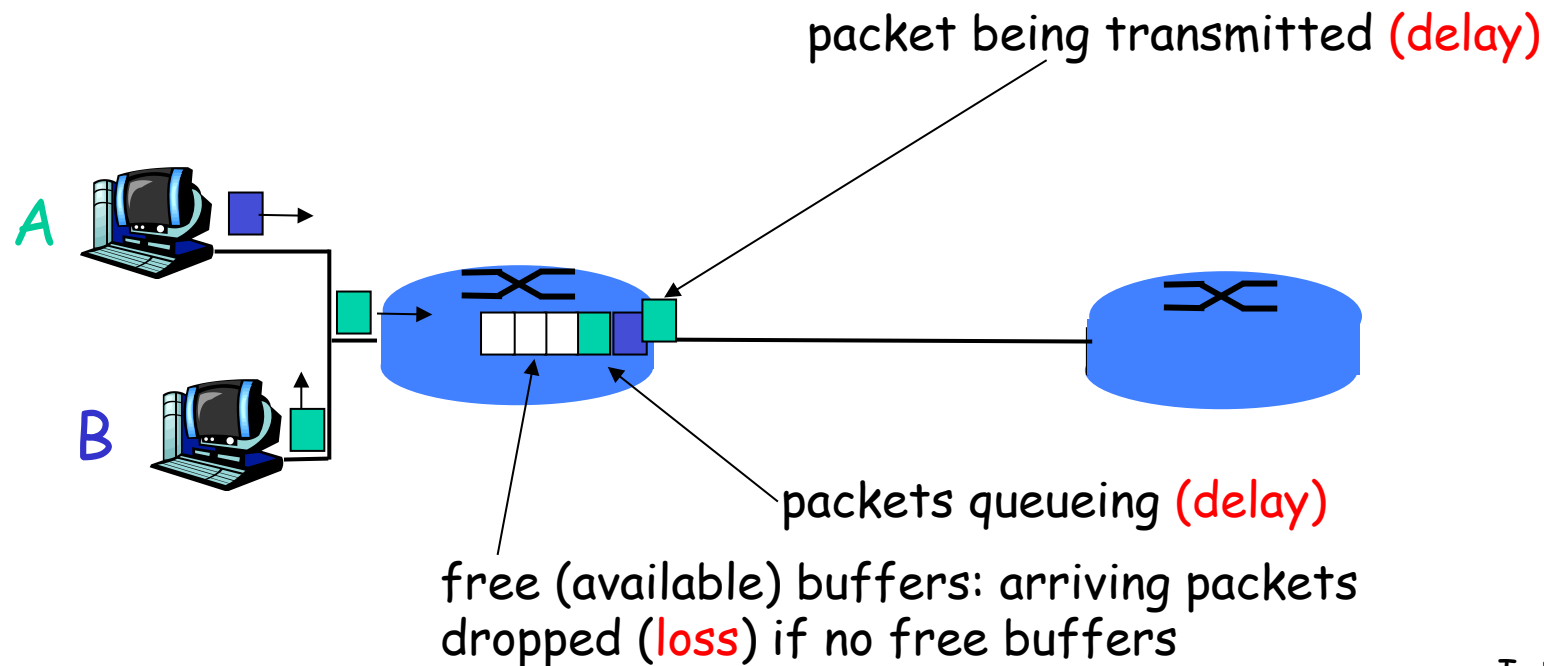
1.7 Protocol layers, service models

1.8 History

How do loss and delay occur?

packets *queue* in router buffers

- ❑ packet arrival rate to link exceeds output link capacity
- ❑ packets queue, wait for turn



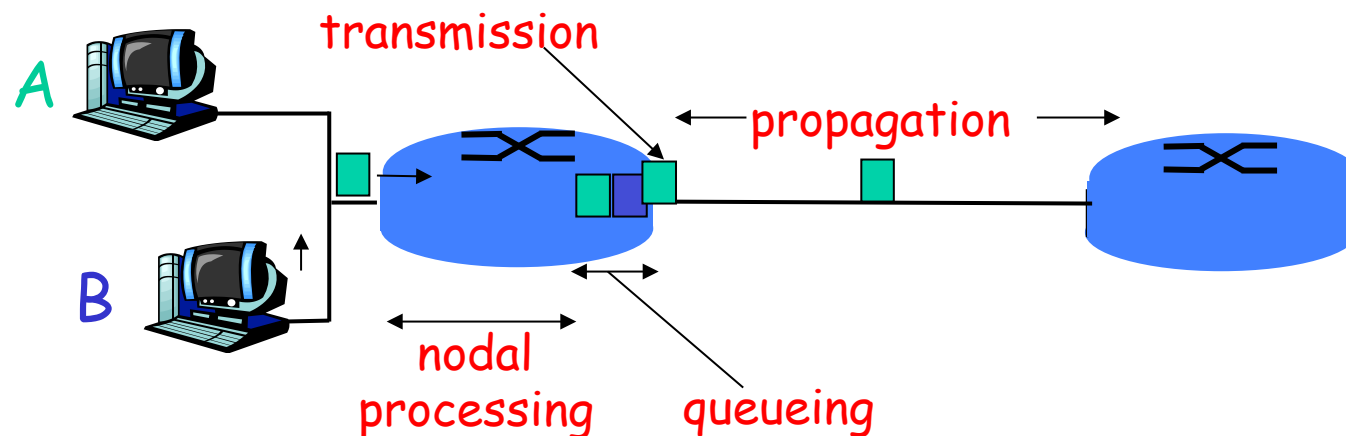
Four sources of packet delay

❑ 1. nodal processing:

- check bit errors
- determine output link

❑ 2. queueing

- time waiting at output link for transmission
- depends on congestion level of router



Delay in packet-switched networks

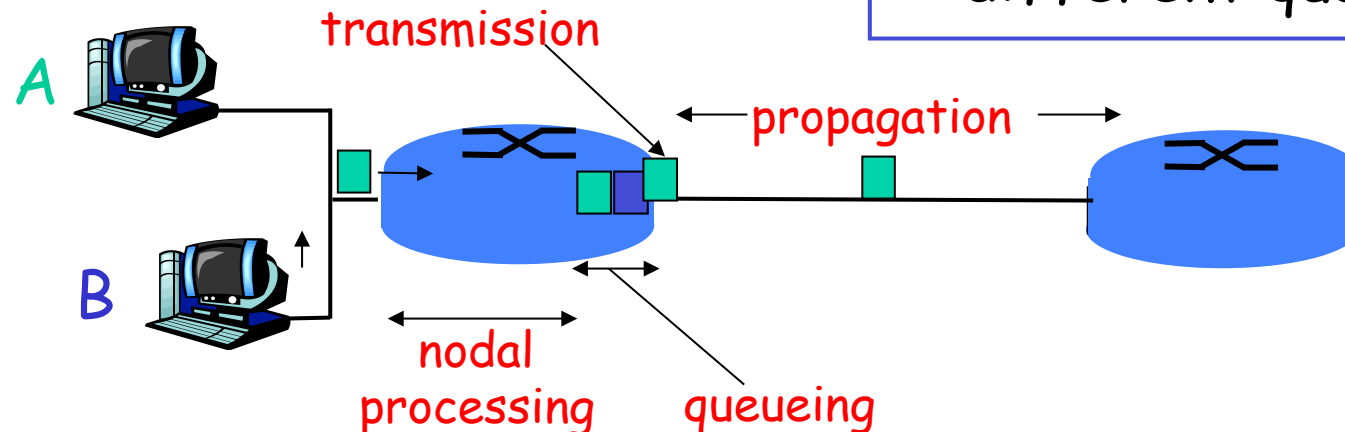
3. Transmission delay:

- R = link bandwidth (bps)
- L = packet length (bits)
- time to send bits into link = L/R

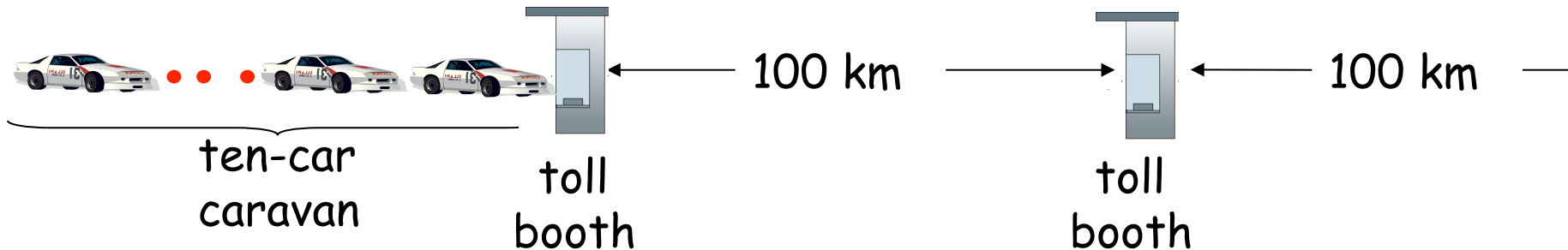
4. Propagation delay:

- d = length of physical link
- s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- propagation delay = d/s

Note: s and R are very different quantities!

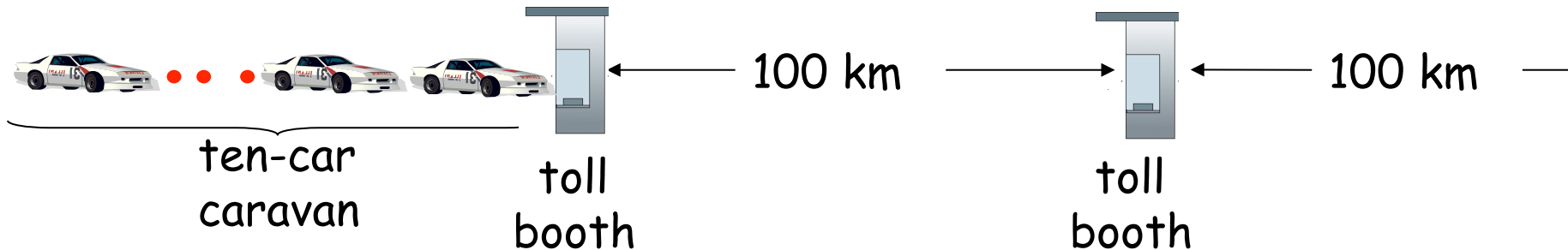


Caravan analogy



- ❑ cars “propagate” at 100 km/hr
 - ❑ toll booth takes 12 sec to service car (transmission time)
 - ❑ car~bit; caravan ~ packet
 - ❑ Q: How long until caravan is lined up before 2nd toll booth?
- ❑ Time to “push” entire caravan through toll booth onto highway = $12 \times 10 = 120$ sec
 - ❑ Time for last car to propagate from 1st to 2nd toll both: $100\text{km}/(100\text{km/hr}) = 1$ hr
 - ❑ A: 62 minutes

Caravan analogy (more)



- ❑ Cars now “propagate” at 1000 km/hr
- ❑ Toll booth now takes 1 min to service a car
- ❑ **Q: Will cars arrive to 2nd booth before all cars serviced at 1st booth?**
- ❑ **Yes!** After 7 min, 1st car at 2nd booth and 3 cars still at 1st booth.
- ❑ 1st bit of packet can arrive at 2nd router before packet is fully transmitted at 1st router!
 - See Ethernet applet at [AWL Web site](#)

Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

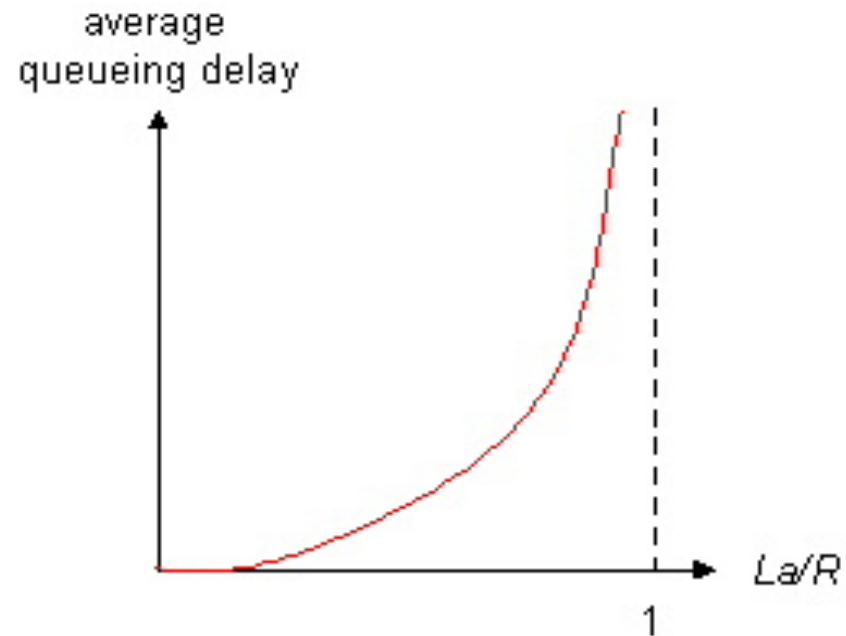
- d_{proc} = processing delay
 - typically a few microsecs or less
- d_{queue} = queuing delay
 - depends on congestion
- d_{trans} = transmission delay
 - $= L/R$, significant for low-speed links
- d_{prop} = propagation delay
 - a few microsecs to hundreds of msecs

Delay for each hop!!!

Queueing delay (revisited)

- R =link bandwidth (bps)
 - L =packet length (bits)
 - a =average packet arrival rate
- rate

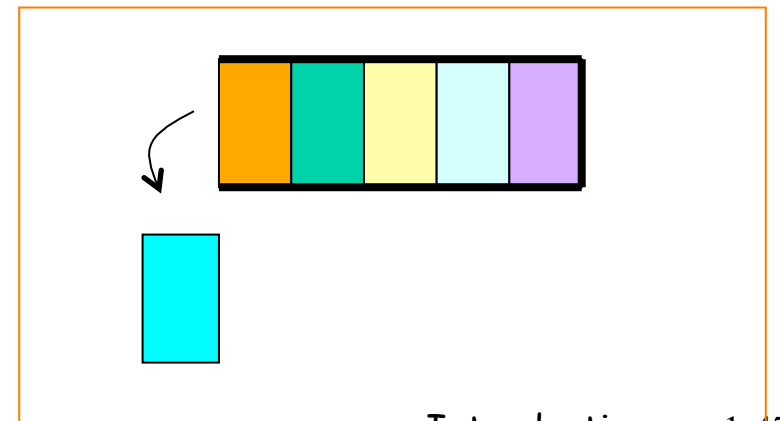
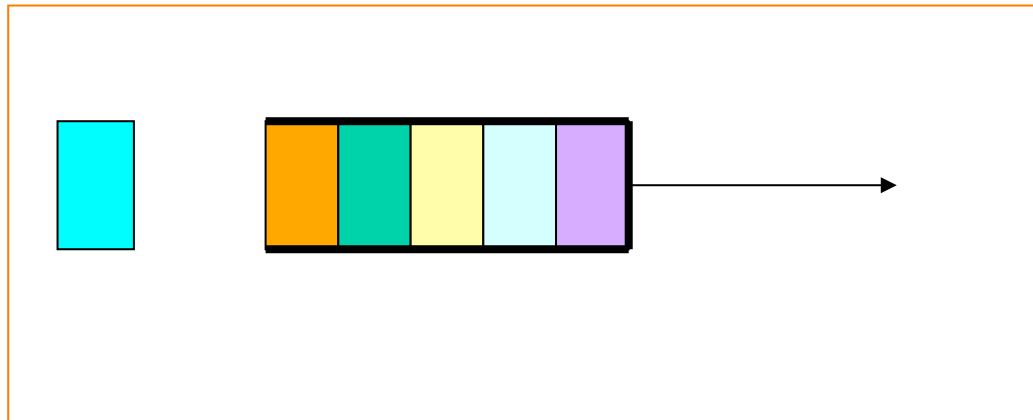
traffic intensity = La/R



- $La/R \sim 0$: average queueing delay small
- $La/R \rightarrow 1$: delays become large
- $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!

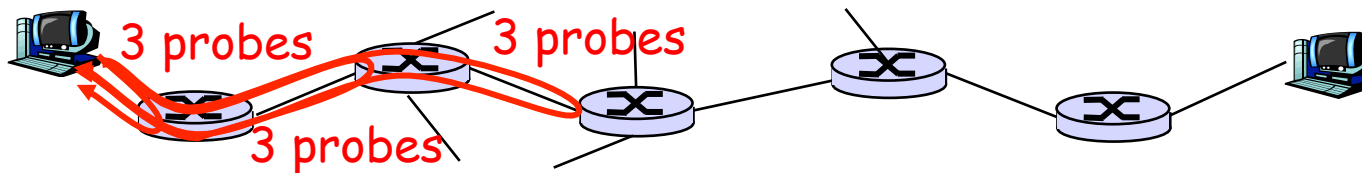
Packet loss

- ❑ queue (→buffer) preceding link in buffer has finite capacity
- ❑ when packet arrives to full queue, packet is dropped (→lost)
- ❑ lost packet may be retransmitted by previous node, by source end system, or not retransmitted at all



“Real” Internet delays and routes


- ❑ What do “real” Internet delay & loss look like?
- ❑ Trace route program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



“Real” Internet delays and routes


traceroute: gaia.cs.umass.edu to www.eurecom.fr

Three delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu




1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 * * *
18 * * *
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

trans-oceanic link



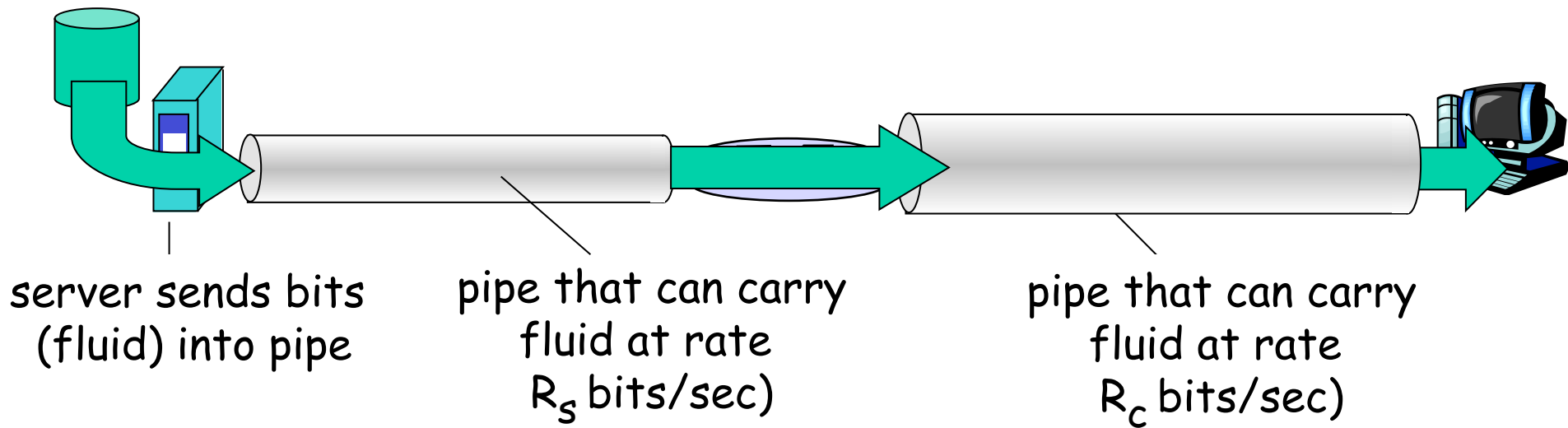
* means no reponse (probe lost, router not replying)



Name and address of router, round trip delays (3 samples)

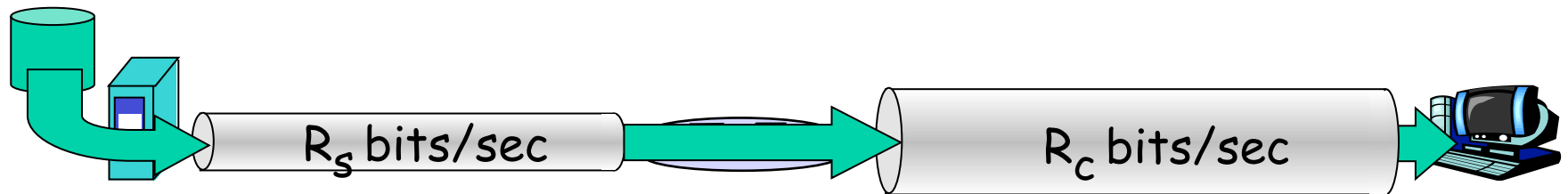
Throughput

- ❑ *throughput*: rate (bits/time unit) at which bits transferred between sender/receiver
 - *instantaneous*: rate at given point in time
 - *average*: rate over longer period of time

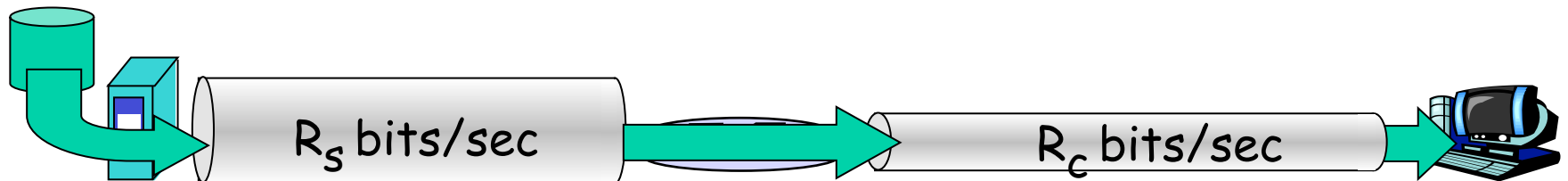


Throughput (more)

- $R_s < R_c$ What is average end-end throughput?



- $R_s > R_c$ What is average end-end throughput?

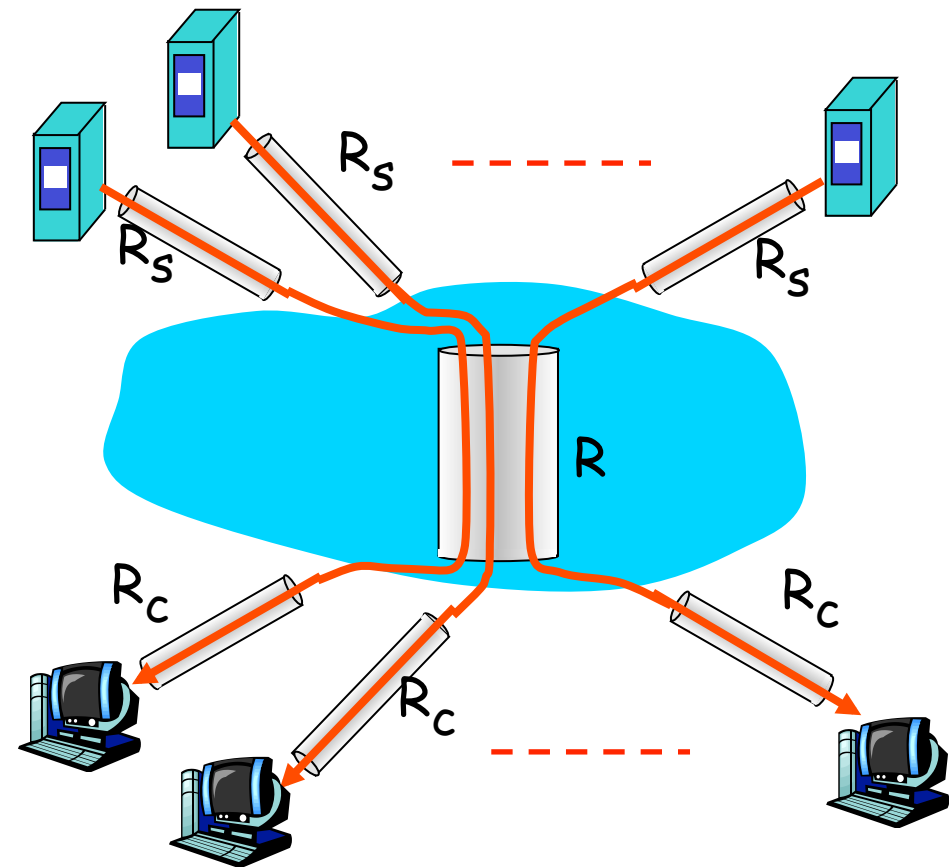


bottleneck link

link on end-end path that constrains end-end throughput

Throughput: Internet scenario

- per-connection end-end throughput:
 $\min(R_c, R_s, R/10)$
- in practice: R_c or R_s is often bottleneck



10 connections (fairly) share backbone
bottleneck link R bits/sec