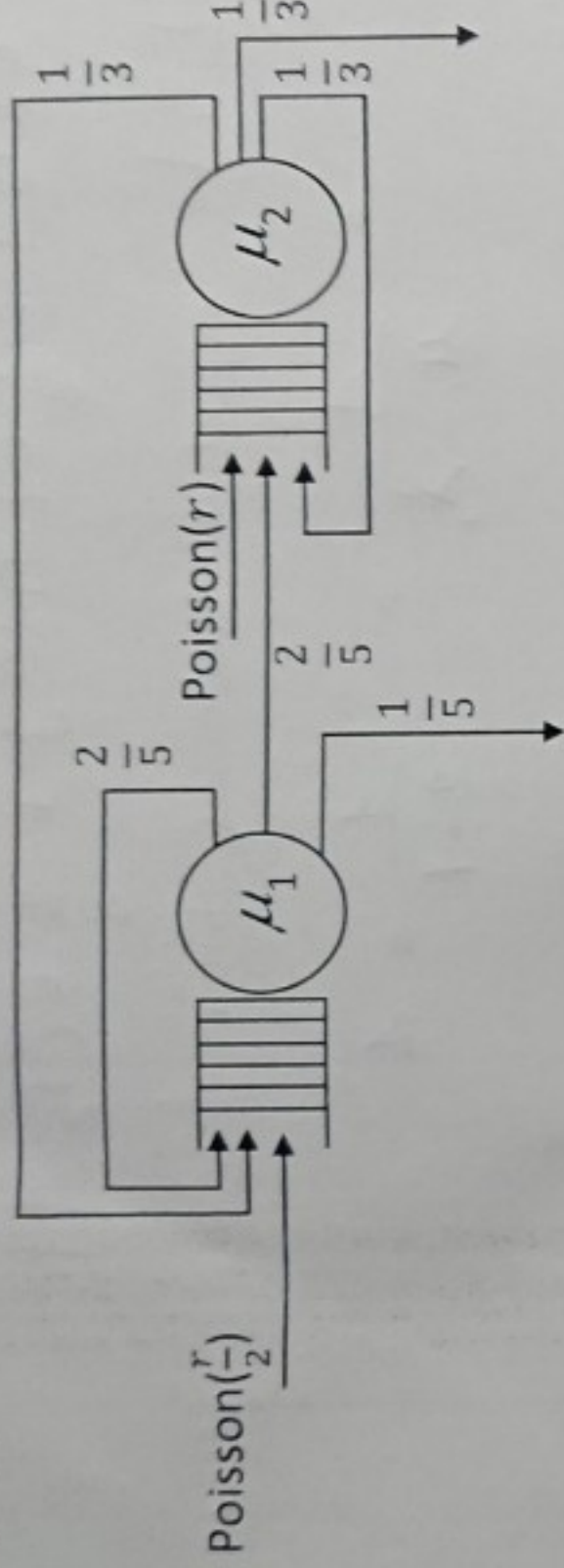


Ex. 1 (10 pts)



Consider the network given in the figure, where two routers have an exponential service rate of $\mu_1 = 2$ packets/second and $\mu_2 = 6$ packets/second, respectively. Each router has an infinite, FCFS buffer, and routing is probabilistic according to the probability values shown in the figure. There are two entry points of average arrival rate $r/2$ and r , respectively, as shown in the figure, which are exponentially distributed.

Answer the following questions:

- What is the maximum allowable value r_{MAX} of the parameter r defining the two arrival rates, that the network can tolerate?
- Set $r = 0.5 r_{MAX}$. What is the mean response time T_{TOT}^1 of a packet entering at router 1?

Solve the problem in the space below, then write the answers in the box.

Answers:

$$r_{MAX} = 0,8$$

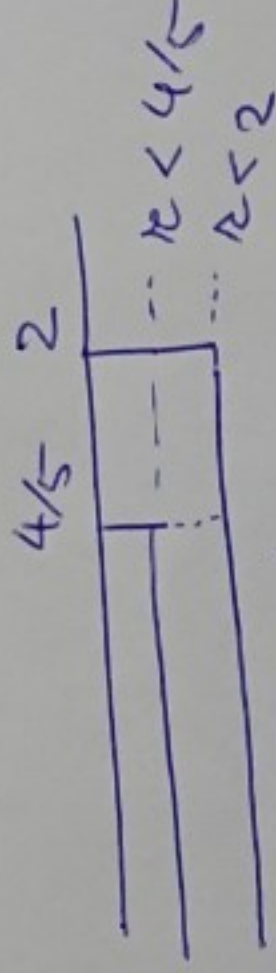
$$T_{TOT}^1 = 45/16$$

$$\left. \begin{aligned} \hat{\lambda}_1 &= \frac{r}{2} + \frac{2}{5} \hat{\lambda}_1 + \frac{1}{3} \hat{\lambda}_2 \\ \hat{\lambda}_2 &= r + \frac{2}{5} \hat{\lambda}_1 + \frac{1}{3} \hat{\lambda}_2 \end{aligned} \right\} \Rightarrow \hat{\lambda}_1 = \frac{5}{2} r ; \hat{\lambda}_2 = 3r$$

To ensure stability we impose $\rho_i < 1 \quad i=1,2$

$$\rho_1 = \frac{\hat{\lambda}_1}{\mu_1} = \frac{5}{4} r < 1$$

$$\rho_2 = \frac{\hat{\lambda}_2}{\mu_2} = \frac{3}{2} r < 1$$



$$r < \frac{4}{5} \rightarrow r_{MAX} = 0,8 \quad (a)$$

$$r = 0.5 r_{MAX} = 0,4$$

$$\Rightarrow \hat{\lambda}_1 = \frac{5}{2} \cdot \frac{4}{5} = 2 ; \hat{\lambda}_2 = 3 \cdot 0,4 = 1,2$$

$$E[T_1^{\text{tot}}] = E[T_1] + \frac{2}{5} E[T_1^{\text{tot}}] + \frac{2}{5} E[T_2^{\text{tot}}]$$

$$E[T_2^{\text{tot}}] = E[T_2] + \frac{1}{3} E[T_1^{\text{tot}}] + \frac{1}{3} E[T_2^{\text{tot}}]$$

where $E[T_1]$ = expected traversal time for a single visit =

$$= \frac{1}{\mu_1 - \hat{\lambda}_1} = \frac{1}{2-1} = 1$$

$$\text{and } E[T_2] = \frac{1}{\mu_2 - \hat{\lambda}_2} = \frac{1}{6-1.2} = \frac{5}{24}$$

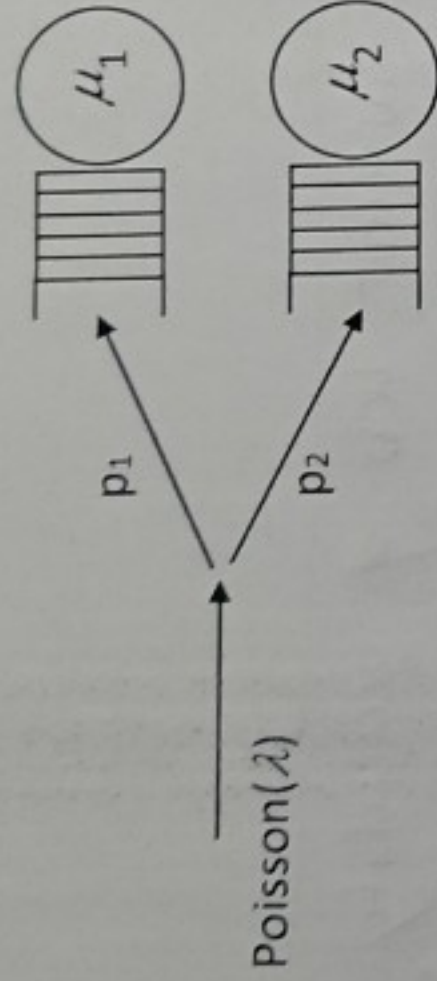
$$\text{We set } \begin{aligned} x &= E[T_1^{\text{tot}}] \\ y &= E[T_2^{\text{tot}}] \end{aligned}$$

$$\left. \begin{aligned} x &= 1 + \frac{2}{5}x + \frac{2}{5}y \\ y &= \frac{5}{24} + \frac{1}{3}x + \frac{1}{3}y \end{aligned} \right\} \Rightarrow \begin{aligned} 3x - 2y &= 1 \\ 16y - 8x &= 5 \end{aligned}$$

$$\begin{aligned} x &= \frac{45}{16} \\ y &= \frac{55}{32} \end{aligned}$$

$$\Rightarrow E[T_1^{\text{tot}}] = x = \frac{45}{16}$$

Ex. 2 (13 pts)



Consider the network given in the figure, where an arrival flow with Poissonian distribution and rate λ is routed probabilistically between two servers with exponential service rates of μ_1 and μ_2 , respectively. Prove or disprove that, in the general case with $\mu_1 \neq \mu_2$, load balancing is the best choice in minimizing WAITING TIME and RESPONSE TIME of the above system.

Answers:

Load balancing gives optimal response time (Yes/No):

Load balancing gives optimal waiting time (Yes/No):

Case of load balancing: $E[V_1] \cdot E[S_1] = E[V_2] \cdot E[S_2] \Rightarrow \frac{P_1}{\mu_1} = \frac{P_2}{\mu_2}$

we set $p \triangleq p_1 \Rightarrow p_2 = 1-p \Rightarrow \frac{p}{\mu_1} = \frac{1-p}{\mu_2}$ from which

$$p \triangleq P_1 = \frac{p\lambda}{\mu_1} = \frac{(1-p)\lambda}{\mu_2} = P_2$$

$$E[\tau] = p \cdot E[\tau_1] + (1-p) \cdot E[\tau_2]$$

$$E[\tau_Q] = p E[\tau_{Q1}] + (1-p) E[\tau_{Q2}] =$$

$$E[\tau_1] = \left(\frac{p}{1-p} \right) \cdot \frac{1}{p\lambda}$$

$$E[\tau] = \left(\frac{p}{1-p} \right) \cdot \left(\frac{1}{\lambda} + \frac{1}{\lambda} \right) = \frac{2p}{\lambda(1-p)}$$

$$E[\tau_2] = \frac{p}{1-p} \cdot \frac{1}{(1-p)\lambda} = \frac{1}{\lambda(1-p)^2} \cdot p = \frac{p}{\lambda(1-p)^2}$$

$$= \frac{2p}{\lambda(1-p)} - \frac{1}{\lambda} (p+p) =$$

$$= \frac{2p}{\lambda(1-p)} - \frac{2p}{\lambda} = \frac{2p(1-1+p)}{\lambda(1-p)} =$$

$$= \frac{2p^2}{\lambda(1-p)}$$

In summary in case of load balancing we have $p = p_1 = p_2$, and

$$E[T] = \frac{2p}{\lambda(1-p)}$$

$$E[T_Q] = \frac{2p^2}{\lambda(1-p)}$$

Let us consider $\mu_1 > \mu_2$ without loss of generality and let $p_1 = 1$ $p_2 = 0$ (we send all the jobs to the faster server $\Rightarrow$ UNBALANCED LOAD)

We have

$$E[T] = \frac{1}{\mu_1 - \lambda}$$

$$\text{and } E[T_Q] = \frac{1}{\mu_1 - \lambda} - \frac{1}{\mu_1}$$

Consider now a numerical example with $\lambda = 1$, $\mu_1 = 10$ and $\mu_2 = 1$

Balanced case

$$E[T]_B = \frac{2p}{\lambda(1-p)}$$

$$\frac{2 \cdot \frac{1}{11}}{1(10/11)} = 1/5$$

$$p = \frac{p_1 \lambda}{\mu_1}$$

$$\begin{cases} p_1 + p_2 = 1 \\ p_1 = \frac{\mu_1}{\mu_2} p_2 \end{cases} \Rightarrow \begin{cases} p_1 = \mu_1 / (\mu_1 + \mu_2) = \frac{10}{11} \\ p_2 = \mu_2 / (\mu_1 + \mu_2) = 1/11 \end{cases}$$

$$\Rightarrow p = \frac{10/11}{10} = \frac{1}{11}$$

$$E[T_Q]_B = \frac{2 \cdot \frac{1}{11}}{(1 - 1/11)} = \frac{2 \cdot \frac{1}{11}}{10} = \frac{1}{55}$$

Same numerical example but unbalanced case

$$E[T]_{UB} = \frac{1}{10-1} = \frac{1}{9} < E[T_B] = \frac{1}{5} \rightarrow \text{better}$$

$$E[T_Q]_{UB} = \frac{1}{9} - \frac{1}{10} = \frac{1}{90} < E[T_Q]_B = \frac{1}{55}$$

if unbalanced case if load is concentrated on the faster server.