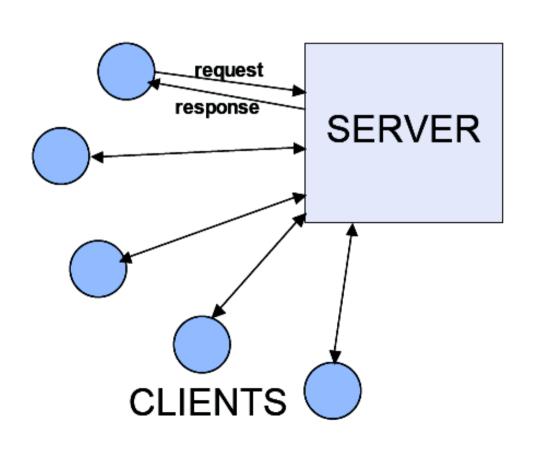
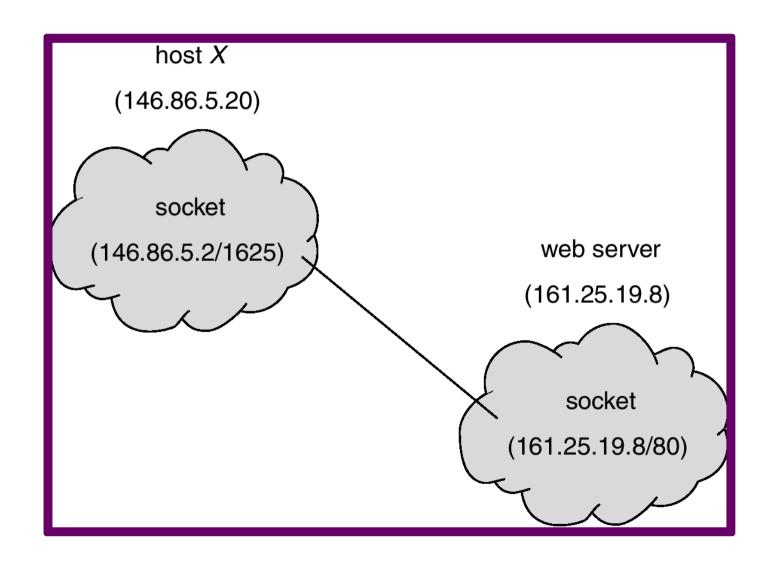
- File I/O:
  - open, close, read, write, seek, fcntl, ...
- Network communication:
  - developers extended set of file descriptors to include network connections.
  - extended read/write to work on these new file descriptors.
  - but other required functionality did not fit into the 'openread-write-close' paradigm.
- -> Socket API

- Server listens for requests from clients
- Server: passive open
- Client: active open
- Example:
  - file server
  - web server
  - print server
  - mail server
  - name server
  - X window server





### **Basics**

- The basic building block for communication is the socket.
- A socket is an <u>endpoint of communication</u> to which a <u>name may be bound</u>.
- Each socket in use has a <u>type</u> and <u>one or more</u> associated processes.

### **Domains**

- Sockets exist within <u>communication domains</u>.
- A communication domain is an <u>abstraction</u> introduced to <u>bundle common properties</u> of processes communicating through sockets, e.g. socket name.
- For example, in the UNIX communication domain sockets are named with UNIX path names; e.g. a socket may be named "/dev/foo".
- Sockets normally exchange data only with sockets in the same domain

### **Domains**

- The 4.4BSD IPC facilities supported <u>four</u> separate communication domains
  - the *UNIX domain*, for on-system communication
  - the <u>Internet domain</u>, which is used by processes which communicate using the Internet standard communication protocols
  - the NS domain, which is used by processes which communicate using the Xerox standard communication protocols
  - the ISO OSI protocols

## Socket Types

- Sockets are typed according to the communication properties visible to a user.
- Processes are presumed to communicate only between sockets of the same type
- Four types of sockets currently are available

#### Stream Sockets

- A <u>stream socket</u> provides for the <u>bidirectional</u>, <u>reliable</u>, <u>sequenced</u>, and <u>unduplicated</u> flow of data <u>without record boundaries</u>.
- Aside from the bidirectionality of data flow, a pair of connected stream sockets provides an interface nearly identical to that of pipes

## Datagram Sockets

- A <u>datagram socket</u> supports <u>bidirectional</u> flow of data which is <u>not promised to be sequenced</u>, <u>reliable</u>, <u>or unduplicated</u>.
- Messages may be dropped, duplicated, and, possibly, delivered in an order different from the order in which they were was sent.
- An important characteristic of a datagram socket is that <u>record boundaries in data are</u> <u>preserved</u>.

#### Raw Sockets

- A <u>raw socket</u> provides users <u>access to the</u> <u>underlying communication protocols</u> which support socket abstractions.
- These sockets are <u>normally datagram oriented</u>,
- Not intended for the general user; they have been provided mainly for those interested in developing new communication protocols, or for gaining access to some of the more esoteric facilities of an existing protocol.

### Sequenced Packet Sockets

- A sequenced packet socket is similar to a stream socket, with the exception that record boundaries are preserved.
- This interface is provided only as part of the NS socket abstraction.

### **Socket Creation**

```
s = socket(domain, type, protocol);
```

- Create a socket in the specified domain and of the specified type.
- A particular protocol may also be requested.
- If the protocol is left unspecified (a value of 0), the system will select an appropriate protocol. from those protocols which comprise the communication domain and which

### **Socket Creation**

```
s = socket(domain, type, protocol);
```

 The user is returned a descriptor (a small integer number) which may be used in later system calls which operate on sockets.

### **Socket Creation**

```
s = socket(domain, type, protocol);
```

- The domain is specified as one of:
  - AF\_UNIX (Unix Domain)
  - AF\_INET (Internet Domain)
  - AF\_NS (NS Domain)
- The socket types are:
  - SOCK STREAM
  - SOCK\_DGRAM
  - SOCK RAW
  - SOCK\_SEQPACKET

# Examples

```
s = socket(AF_INET, SOCK_STREAM, 0);
```

Creates a stream socket in the Internet domain

```
s = socket(AF UNIX, SOCK DGRAM, 0);
```

- Creates a datagram socket for on-machine use (Unix Domain)
- The default protocol (last argument to the socket call is 0) should be correct for almost every situation

### **Socket Names**

- A socket is created without a name.
- Until a name is bound to a socket, processes have no way to reference it and, consequently, no messages may be received on it.
- Communicating processes are bound by an association. In the Internet (and NS) domains, an association is composed of local and foreign addresses, and local and foreign ports.
- In the UNIX domain, an association is composed of local and foreign path names.

#### **Socket Names**

- In the Internet domain there may never be duplicate protocol, local address, local port, foreign address, foreign port> tuples.
- UNIX domain sockets need not always be bound to a name, but when bound there may never be duplicate protocol, local pathname, foreign pathname> tuples.

# Binding Names

```
bind(s, name, namelen);
```

- The <u>bind()</u> system call allows a process to specify half of an association, <local address, local port> (or <local pathname>).
- The <u>connect()</u> and <u>accept()</u> primitives are used to complete a socket's association.
- The bound name is a <u>variable length byte string</u> which is interpreted by the supporting protocol(s).

# Binding Names

- In the Internet domain names contain an Internet address and port number.
- In the UNIX domain, names contain a path name and a family, which is always AF\_UNIX.

### Example

```
#include <sys/un.h>
struct sockaddr_un addr;
strcpy(addr.sun_path, "/tmp/foo");
addr.sun_family = AF_UNIX;
len = strlen(addr.sun_path) +
       sizeof (addr.sun_family)
bind(s, (struct sockaddr *) &addr, len);
```

# Binding Names

- File name referred to in addr.sun\_path is created as a socket in the system file space.
- The caller must, therefore, have write permission in the directory where addr.sun\_path is to reside, and this file should be deleted by the caller when it is no longer needed.

### Example

```
#include <sys/types.h>
#include <netinet/in.h>
...
struct sockaddr_in sin;
...
bind(s, (struct sockaddr *) &sin, sizeof (sin));
```

- Connection establishment is usually asymmetric, with one process a "client" and the other a"server".
- The server binds a socket to a well-known address and then passively "listens"
- The client requests services from the server by initiating a "connection" to the server's socket.
- On the client side the connect call is used to initiate a connection.

```
// Unix Domain
struct sockaddr un server;
connect(s, (struct sockaddr *)&server,
 strlen(server.sun_path) + sizeof
  (server.sun_family));
// Internet Domain
struct sockaddr_in server;
connect(s, (struct sockaddr *)&server, sizeof
  (server));
```

- server would contain either the UNIX
  pathname, Internet address and port number of
  the server to which the client process wishes to
  speak.
- If the client process's socket is unbound at the time of the connect call, the system will automatically select and bind a name to the socket if necessary

- An error is returned if the connection was unsuccessful (any name automatically bound by the system, however, remains).
- Otherwise, the socket is associated with the server and data transfer may begin.

### Server Side

 For the server to receive a client's connection it must perform two steps after binding its socket.

```
listen(s, 5);
```

- Means that the server is willing to listen for incoming connection requests
- The second parameter specifies the maximum number of outstanding connections which may be queued awaiting acceptance

### Server Side

A server may accept a connection:

```
struct sockaddr_in from;
...
fromlen = sizeof (from);
newsock = accept(s, (struct sockaddr *)&from, &fromlen);
```

### Server Side

- A new descriptor is returned on receipt of a connection (along with a new socket).
- fromlen:
  - input; how much space is associated with from
  - output: size of the name
- The second parameter may be a null pointer.

- Accept will not return until a connection is available or the system call is interrupted by a signal to the process.
- Further, there is no way for a process to indicate it will accept connections from only a specific source

### Data Transfer

Normal read and write system calls are usable:

```
write(s, buf, sizeof (buf));
read(s, buf, sizeof (buf));
```

#### But also:

```
send(s, buf, sizeof (buf), flags);
recv(s, buf, sizeof (buf), flags);
```

### Data Transfer

- send/recv flags:
  - MSG\_OOB send/receive out of band data
  - MSG\_PEEK look at data without reading
  - MSG\_DONTROUTE send data without routing packets
- When MSG\_PEEK is specified with a recv call, any data present is returned to the user, but treated as still "unread".
- Next read or recv call applied to the socket will return the data previously previewed.

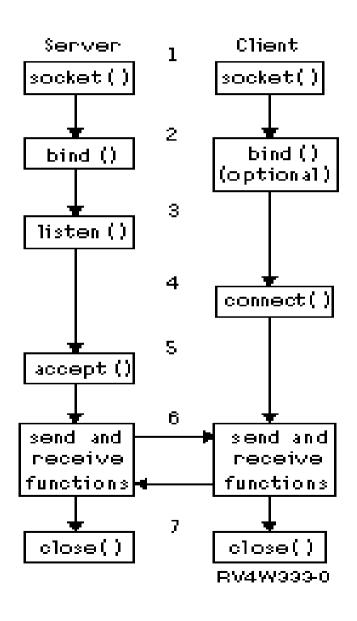
# Closing

Once a socket is no longer of interest, it may be discarded

```
shutdown(s, how);
close(s);
```

- how:
  - SHUT RD
  - SHUT WR
  - SHUT\_RDWR

### Socket Calls Flow



## Datagram Sockets

- connect() on datagram sockets returns immediately
  - The system simply records the peer's address
  - On a stream socket a connect request initiates the connection.
- Only one connected address is permitted for each socket at one time
  - a second connect will change the destination
- accept() and listen() are not used with datagram sockets.

#### Connectionless Sockets

- Only with datagram sockets (!)
- sendto() specifies a destination address

- to and tolen indicate the address of recipient
- recvfrom() receives messages on an unconnected datagram socket

#### netdb

- Routines for
  - mapping host names to network addresses, network names to network numbers
  - protocol names to protocol numbers
  - service names to port numbers and the appropriate protocol
- The file <netdb.h> must be included

#### **Host Names**

Internet host name to address mapping

```
struct hostent {
              /* official name of host */
 char *h_name;
 char **h aliases; /* alias list */
 int h_addrtype;
                    /* host address type
                       (e.g., AF_INET) */
               /* length of address */
 int h_length;
 char **h_addr_list; /* list of addresses,
                       null terminated */
#define h_addr h_addr_list[0] /* first address,
                       network byte order */
```

#### **Host Names**

```
gethostbyname(const char *name)
```

takes an host name and returns a hostent structure

```
gethostbyaddr(const char *addr,
  int len, int type)
```

 maps Internet host addresses (AF\_INET, AF\_INET6) into a hostent structure

#### **Network Names**

```
struct netent {
                  /* official name of net */
 char *n name;
 char **n aliases; /* alias list */
               /* net address type */
 int n_addrtype;
                    /* network number,
 int n net;
                     host byte order */
};
getnetbyname(const char *name);
getnetbynumber(long net,
                    int type);
```

#### **Protocol Names**

#### Service Names

```
struct servent {
 char *s name; /* official service name */
 char **s aliases; /* alias list */
 int s_port; /* port number,
                      network byte order */
 char *s_proto; /* protocol to use */

    getservbyname(const char *name,

                const char *proto);
getservbyport(int port, const char *proto);
sp = getservbyname("telnet", (char *) 0));
sp = getservbyname("telnet", "tcp");
```

#### **Endiannes**

#### Big Endian

 the most significant byte of any multibyte data field is stored at the lowest memory address

#### Little Endian

 the least significant byte of any multibyte data field is stored at the lowest memory address

#### Host Independent Formats

- Intel CPUs are Little Endian, while the network byte order is Big Endian
- Macros to convert "host" order to "network byte order

```
char c1 = 1;
char c2 = 2;
short s = 255; // 0x00FF
long l = 0x44332211;
```

```
Offset : Memory dump
0x00000 : 01 02 FF 00
0x00004 : 11 22 33 44
```

# Network Byte Order

- htonl(val)
  - convert 32-bit quantity from host to network byte order
- htons(val)
  - convert 16-bit quantity from host to network byte order

# Network Byte Order

- ntohl(val)
  - convert 32-bit quantity from network to host byte order
- ntohs(val)
  - convert 16-bit quantity from network to host byte order

```
#include <stdio.h>
#include <netdb.h>
#include <stdlib.h>
unsigned long ResolveName(char name[])
  struct hostent *host;
  if ((host = gethostbyname(name)) == NULL){
    fprintf(stderr, "gethostbyname() failed");
    exit(1);
  return *((unsigned long *)host->h_addr_list[0]);
```

```
unsigned short ResolveService(char service[],
   char protocol[])
  struct servent *serv;
  unsigned short port;
  if ((port = atoi(service)) == 0) {
     if ((serv = getservbyname(service, protocol)) ==
  NULL) {
        fprintf(stderr, "getservbyname() failed");
        exit(1);
     else port = serv->s_port;
  else port = htons(port);
  return port;
```

# Multiplexing

 select() allows multiplexing i/o requests among multiple sockets and/or files

```
#include <sys/time.h>
#include <sys/types.h>
fd_set readmask, writemask, exceptmask;
struct timeval timeout;
select(nfds, &readmask, &writemask, &exceptmask,
 &timeout);
```

- select() takes pointers to three sets
  - one for the <u>set of file descriptors</u> for which the caller wishesto be able to read data on
  - one for those descriptors to which data is to be written
  - one for which exceptional conditions are pending (out-of-band data is the only exceptional condition currently implemented bythe socket
- If the user is not interested in certain conditions the corresponding argument should be NULL.

- Each set is actually a structure containing an array of long integer bit masks
- The size of the array is set by the definition FD\_SETSIZE
- The macros

```
- FD_SET(fd, &mask)
```

- FD\_CLR(fd, &mask)

allow adding and removing file descriptor fd in the set mask.

The set should be zeroed before use

```
- FD_ZERO(&mask)
```

 nfds specifies the range of file descriptors (i.e. one plus the value of the largest descriptor) to be examined

- A timeout value may be specified
- If timeout (struct timeval) is set to 0, select returns immediately
- If the last parameter is a NULL pointer, the selection will block indefinitely
  - returns only when a descriptor is selectable or when a signal

- select()returns:
  - the number of file descriptors selected
  - 0 if the select call returns due to the timeout expiring
  - -1 if terminated because of an error or interruption

- The status of a file descriptor may be tested
  - FD\_ISSET(fd, &mask)
  - returns a non-zero value if fd is a member of the set mask, and 0 if it is not

```
#include <sys/time.h>
#include <sys/types.h>
...

fd_set read_template;
struct timeval wait;
...
```

```
for (;;) {
   wait.tv_sec = 1; /* one second */
   wait.tv_usec = 0;
   FD_ZERO(&read_template);
   FD_SET(s1, &read_template);
   FD_SET(s2, &read_template);
   nb = select(FD_SETSIZE, &read_template,
        (fd_set *) 0, (fd_set *) 0, &wait);
```

```
if (nb <= 0) {
   if (nb<0) perror("select")</pre>
  else printf("Timeout.\n);
  continue;
  (FD_ISSET(s1, &read_template)) {
  sts=ReadDataFromSocket(s1)
  (FD_ISSET(s2, &read_template)) {
  sts=ReadDataFromSocket(s2)
```

- select() provides a synchronous multiplexing scheme.
- Asynchronous notification of output completion, input availability, and exceptional conditions is possible through use of signals (SIGIO and SIGURG)

# Closing Sockets

```
s=connect(...);
if( fork() ) { /* Child */
    while(gets(buffer) >0) write(s,buf,strlen(buffer));
    close(s);
    exit(0);
else { /* Parent */
    while((l=read(s,buffer,sizeof(buffer)) do_something(l,buffer);
    wait(0);
    exit(0);
```

#### Socket Shutdown

```
s=connect(...);
if( fork() ) { /* Child */
    while( gets(buffer) >0) write(s,buf,strlen(buffer));
    close(s);
    shutdown(s,SHUT WR);
    exit(0);
else { /* Parent */
    while((l=read(s,buffer,sizeof(buffer)) do_something(l,buffer);
    wait(0);
    exit(0);
```

# setsockopt/getsockopt

```
int setsockopt(int s, int level, int optname,
  const void *optval, int optlen);
int getsockopt(int s, int level, int optname,
  void *optval, socklen_t *optlen)M
```

- Manipulate the options associated with a socket.
- Options may exist at multiple protocol levels; they are always present at the uppermost socket level (SOL\_SOCKET)

# setsockopt/getsockopt

- A server waits 2 MSL (maximum segment lifetime) for old connection.
- If not properly terminated, a further bind() will return EADDRINUSE.

# setsockopt/getsockopt

• Before bind():

#### Other options:

- SO\_ERROR get error status
- SO\_KEEPALIVE send periodic keep-alives
- SO\_LINGER close() on non-empty buffer
- SO\_SNDBUF send buffer size
- SO\_RCVBUF receive buffer size

# Non Blocking I/O

 Once a socket has been created it may be marked as non-blocking

```
#include <fcntl.h>
int s;
s = socket(AF_INET, SOCK_STREAM, 0);
if (fcntl(s, F SETFL, FNDELAY) < 0)
 perror("fcntl F_SETFL, FNDELAY");
exit(1)
```

# NonBlocking I/O

- NB: must check for errno==EWOULDBLOCK
- If an operation, such as a send, cannot be done in its entirety the data that can be sent immediately will be processed, and the return value will indicate the amount actually sent

#### SIGIO

- Allows a process to be notified via a signal when a socket (or more generally, a file descriptor) has data waiting to be read
- Three steps:
  - set up a SIGIO signal handler
  - set the process id (or process group id) which is to receive notification of pending input to itself
  - enable asynchronous notification of pending I/O (another fcntl() call)

```
#include <fcntl.h>
int io_handler();
signal(SIGIO, io_handler);
if (fcntl(s, F_SETOWN, getpid()) < 0) {
 perror("fcntl F_SETOWN");
 exit(1);
if (fcntl(s, F_SETFL, FASYNC) < 0) {</pre>
 perror("fcntl F_SETFL, FASYNC");
 exit(1);
```

- When a signal is sent to a process while performing a sockets function, several things may occur depending on whether the socket function is defined as a slow function.
- A slow function is a function that can block indefinitely:
  - write(), recv(), send(), recvfrom(), recvmsg(), sendmsg(), accept().
  - All other sockets functions are fast

- Fast functions are not interrupted by a signal
- The signal is raised when these socket functions exit.

- Slow functions are interrupted by a signal if they are blocked waiting for IO (if they are processing IO, they are not interrupted).
  - They are interrupted in the middle of processing by the raising of a signal.
  - They stop what processing they are doing and return the error EINTR.
  - They do not complete the IO that was initiated.
  - The user program must re-initiate any desired IO explicitly.

- There are three signals that can be generated by actions on a socket:
  - SIGPIPE
  - SIGURG
  - SIGIO

- A SIGPIPE is generated when a send()/write()
   operation is attempted on a broken socket.
  - E.g. a socket which has been shutdown().
- The default action is to terminate the process.
- The target of the signal is the process attempting the send()/write().

- SIGIO is somewhat more complex to set up :
  - fcntl(..,F\_SETFL,FASYNC) to enable Async. I/O
  - fnctl(..,F\_SETOWN, pid) to set target process (group) id.
- A SIGIO signal is generated whenever new I/O can complete on a socket

- SIGIO signal is generated when
  - new data arrives at the socket
  - data can again be sent on the socket
  - the socket is either partially or completely shutdown or when
  - a listen socket has a connection request posted on it

— ...

- A SIGURG indicates that an urgent condition is present on a socket.
- Either the arrival of out of band data or the presence of control status information
  - fnctl(..,F\_SETOWN, pid) to set target process (group) id.