## Exercises on the topic of class 19

## Exercises with solutions

Ex. 1. Minimize the Mealy automaton specified by the following table:

|  | 0 | 1 |
| :---: | :---: | :---: |
| S0 | $\mathrm{S} 1 / 0$ | $\mathrm{~S} 0 / 0$ |
| S1 | $\mathrm{S} 2 / 0$ | $\mathrm{~S} 4 / 0$ |
| S2 | $\mathrm{S} 2 / 0$ | $\mathrm{~S} 1 / 0$ |
| S3 | $\mathrm{S} 1 / 0$ | $\mathrm{~S} 3 / 1$ |
| S4 | $\mathrm{S} 1 / 0$ | $\mathrm{~S} 0 / 0$ |
| S5 | $\mathrm{S} 2 / 0$ | $\mathrm{~S} 3 / 1$ |

SOLUTION:
Since the starting state is not specified, we cannot speak about unreachable states. Let us them build the triangular table:

| S1 | $\begin{aligned} & (\mathrm{S} 1, \mathrm{~S} 2) \\ & (\mathrm{S} 0, \mathrm{~S} 4) \\ & \hline \end{aligned}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| S2 | $\begin{aligned} & (\mathrm{S} 1, \mathrm{~S} 2) \\ & (\mathrm{S} 0, \mathrm{~S} 1) \\ & \hline \end{aligned}$ | (S1,S4) |  |  |  |
| S3 | X | X | X |  |  |
| S4 | 0 | $\begin{aligned} & \hline(\mathrm{S} 1, \mathrm{~S} 2) \\ & (\mathrm{S} 0, \mathrm{~S} 4) \\ & \hline \end{aligned}$ | $\begin{aligned} & \hline(\mathrm{S} 1, \mathrm{~S} 2) \\ & (\mathrm{S} 0, \mathrm{~S} 1) \\ & \hline \end{aligned}$ | X |  |
| S5 | X | X | X | (S1,S2) | X |
|  | S0 | S1 | S2 | S3 | S4 |

Since S0 and S4 are equivalent, the pointers to such pairs have to be deleted. In the pointers left, we note that ( $\mathrm{S} 1, \mathrm{~S} 2$ ) depends from ( $\mathrm{S} 1, \mathrm{~S} 4$ ), and vice versa; hence, also these pairs are equivalent. This entails that also ( $\mathrm{S} 0, \mathrm{~S} 1$ ) and ( $\mathrm{S} 3, \mathrm{~S} 5$ ) are equivalent, so as allthe remaining pairs. Hence, the minimal automaton is made up from states $\mathrm{Q} 0=\{\mathrm{S} 0, \mathrm{~S} 1, \mathrm{~S} 2, \mathrm{~S} 4\}$ and $\mathrm{Q} 1=\{\mathrm{S} 3$, S5\} and can be drawn as follows:


Ex. 2. Find the minimal Moore automaton equivalent to the following Mealy automaton, with starting state $S_{0}$


## SOLUTION:

We first consider all pairs state/output (for transforming the Mealy automaton into the Moore one). Then, the resulting Moore automaton has table

|  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| State $_{\mathrm{Me}}$ Output $_{\mathrm{Me}}$ |  | State ${ }_{\mathrm{Mo}}$ Output $_{\mathrm{Mo}}$ |  | Input = 0 | Input = 1 |
| $\mathrm{S}_{0}$ | 0 | 0 | $\mathrm{~T}_{0}$ | 00 | $\mathrm{~T}_{0}$ |
| $\mathrm{~S}_{0}$ | 0 | 1 | $\mathrm{~T}_{1}$ | 01 | $\mathrm{~T}_{0}$ |
| $\mathrm{~S}_{0}$ | 1 | 0 | $\mathrm{~T}_{2}$ | 10 | $\mathrm{~T}_{5}$ |
| $\mathrm{~S}_{0}$ | 1 | 1 | $\mathrm{~T}_{3}$ | 11 | $\mathrm{~T}_{0}$ |
| $\mathrm{~S}_{1}$ | 0 | 0 | $\mathrm{~T}_{4}$ | 00 | $\mathrm{~T}_{1}$ |
| $\mathrm{~S}_{1}$ | 0 | 1 | $\mathrm{~T}_{5}$ | 01 | $\mathrm{~T}_{5}$ |
| $\mathrm{~S}_{1}$ | 1 | 0 | $\mathrm{~T}_{6}$ | 10 | $\mathrm{~T}_{1}$ |
| $\mathrm{~S}_{1}$ | 1 | 1 | $\mathrm{~T}_{7}$ | 11 | $\mathrm{~T}_{11}$ |
| $\mathrm{~S}_{2}$ | 0 | 0 | $\mathrm{~T}_{8}$ | 00 | $\mathrm{~T}_{11}$ |
| $\mathrm{~S}_{2}$ | 0 | 1 | $\mathrm{~T}_{9}$ | 01 | $\mathrm{~T}_{11}$ |
| $\mathrm{~S}_{2}$ | 1 | 0 | $\mathrm{~T}_{10}$ | 10 | $\mathrm{~T}_{10}$ |
| $\mathrm{~S}_{2}$ | 1 | 1 | $\mathrm{~T}_{11}$ | 11 | $\mathrm{~T}_{10}$ |

We can assume any clone of S 0 to be the starting state; for example, let's pick T0. Hence, $\mathrm{T}_{2}$, $\mathrm{T}_{3}, \mathrm{~T}_{4}, \mathrm{~T}_{6}, \mathrm{~T}_{7}, \mathrm{~T}_{9}$ can be canceled because unreachable. Then, the automaton becomes

| Stato $_{\mathbf{t}}$ | Input $_{\mathbf{t}}=\mathbf{0}$ | Input $_{\mathbf{t}}=\mathbf{1}$ |
| :---: | :---: | :---: |
| $\mathrm{T}_{0} / 00$ | $\mathrm{~T}_{0}$ | $\mathrm{~T}_{5}$ |
| $\mathrm{~T}_{1} / 01$ | $\mathrm{~T}_{0}$ | $\mathrm{~T}_{5}$ |
| $\mathrm{~T}_{5} / 01$ | $\mathrm{~T}_{1}$ | $\mathrm{~T}_{11}$ |
| $\mathrm{~T}_{8} / 00$ | $\mathrm{~T}_{10}$ | $\mathrm{~T}_{8}$ |
| $\mathrm{~T}_{10} / 10$ | $\mathrm{~T}_{10}$ | $\mathrm{~T}_{8}$ |
| $\mathrm{~T}_{11} / 11$ | $\mathrm{~T}_{10}$ | $\mathrm{~T}_{8}$ |

At a first sight, a minimization would seems possible (e.g., by grouping together $<\mathrm{T}_{8}, \mathrm{~T}_{10}, \mathrm{~T}_{11}>$ and $\left\langle\mathrm{T}_{0}, \mathrm{~T}_{1}\right\rangle$, since they have the same row in the above table). This is however not possible since the groupable states have different outputs, whereas those with same output lead to
distinguishable states (and so are not groupable). Indeed, this is confirmed by the triangular table:


Then, ( $\mathrm{T}_{0}, \mathrm{~T}_{8}$ ) must be marked with X since, by reading 0 , we reach $\left(\mathrm{T}_{0}, \mathrm{~T}_{10}\right)$ that is marked with $X$; similarly, ( $T_{1}, T_{5}$ ) must be marked with $X$ becasue always with 0 it reaches ( $\mathrm{T}_{0}, \mathrm{~T}_{1}$ ).
So, all states are distinguishable and the minimal Moore automaton is:


## Exercises without solutions

Ex. 1. Given the following automaton with initial state T0:


Minimize it and give, through temporal diagrams, output and state transitions obtained with input 110001001001.

Ex. 2. Minimize the following automaton:


