

## Operations on Naturals

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## Arithmetics in base 2



All operations are done like in base 10, but *modulo 2*

Es.:  $(1 + 1)_2 = 10_2$

Hence, also carries and borrows work *modulo 2!!*

### SUM:

In base 2, we have:

$$0+0 = 0, \text{ carry} = 0$$

$$0+1 = 1+0 = 1, \text{ carry} = 0$$

$$1+1 = 0, \text{ carry} = 1$$

## Sum (example)



Sum in base 2 the numbers 110001 and 10111.

Solution:

$$\begin{array}{r} 110001 + \\ 10111 = \\ \hline 1001000 \end{array}$$

Indeed:  $110001_2 = (2^5+2^4+1)_{10} = (32+16+1)_{10} = 49_{10}$

$$10111_2 = (2^4+2^2+2+1)_{10} = (16+4+2+1)_{10} = 23_{10}$$

$$1001000_2 = (2^6+2^3)_{10} = (64+8)_{10} = 72_{10}$$


## Overflow



In the previous example, if the numbers were represented with 6 bits, the result would not have been representable (it asks for 7 bits) → *overflow*

No overflow if working with 7 or more bits

*When summing naturals, once we fix the size of the format, we have an overflow if and only if the carry resulting from the sum of the MSBs is 1*

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### Subtraction

Within naturals, subtraction is defined only if the minuend is at least as big as the subtrahend, that is  
 $m - s$  is defined only if  $m \geq s$


Subtraction	Difference	Borrow
0-0	0	0
1-1	0	0
1-0	1	0
0-1	1	1

If there is a borrow and the previous bit is 1, the latter is turned into 0;

If there is a borrow and the previous bit is 0, the latter is turned into 1 and so on for all the following 0's, until we find a 1. This latter bit is turned into 0 and subtraction goes on over the modified minuend;

If no 1 is met, then the subtrahend is bigger than the minuend and so the subtraction is not possible in the naturals.

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### Subtraction


Examples:

$$\begin{array}{r} 11010 - \\ 00100 = \\ \hline 10110 \\ (26-4=22)_{10} \end{array}$$

$$\begin{array}{r} 11000 - \\ 10001 = \\ \hline 00111 \\ (24-17=7)_{10} \end{array}$$

$$\begin{array}{r} 01011 - \\ 10100 = \\ \hline 011001 = \\ 001111 \\ (40-25=15)_{10} \end{array}$$

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### Multiplication

As we are used to, but in base 2:


- Partial products
- Shifting of the partial products
- Sum of the shifted partial products

$$\begin{array}{r} 1011 \times \\ 1101 = \\ \hline 1011 \\ 0000 - \\ 1011 - - \\ 1011 - - - \\ \hline 10001111 \end{array}$$

$$\begin{array}{r} 11_{10} \times \\ 13_{10} = \\ \hline 33 \\ 11 - \\ \hline 143_{10} \end{array}$$

OBS.: the result has a double length!

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### Division

More complex than multiplication  
 Same way as we are used to, but always in base 2!

147:11 = 13 remainder 4

$$\begin{array}{r} 147 \\ 11 \overline{) 147} \\ \underline{11} \phantom{0} \\ 37 \\ \underline{33} \phantom{0} \\ 4 \end{array}$$

$$\begin{array}{r} 10010011 \\ \underline{1011} \phantom{0000} \\ 001110 \phantom{00} \\ \underline{1011} \phantom{00} \\ 001111 \phantom{0} \\ \underline{1011} \phantom{0} \\ 100 \end{array}$$

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