## Chapter 8 roadmap

8.1 What is network security? 8.2 Principles of cryptography 8.3 Message integrity, authentication 8.4 Securing e-mail **8.5** Securing TCP connections: SSL **8.6** Network layer security: IPsec 8.7 Securing wireless LANs 8.8 Operational security: firewalls and IDS

# **Digital signatures**

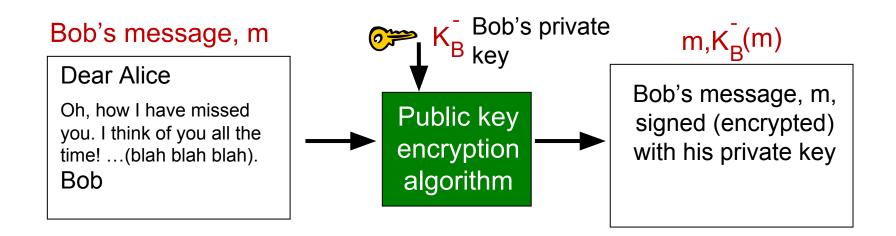
cryptographic technique analogous to handwritten signatures:

- sender (Bob) digitally signs document, establishing he is document owner/creator.
- verifiable, nonforgeable: recipient (Alice) can prove to someone that Bob, and no one else (including Alice), must have signed document

## **Digital signatures**

#### simple digital signature for message m:

Bob signs m by encrypting with his private key K
<sub>B</sub>, creating "signed" message, K
<sub>B</sub>(m)



# **Digital signatures**

- suppose Alice receives msg m, with signature: m,  $K_{B}(m)$
- Alice verifies m signed by Bob by applying Bob's public key  $K_{B}^{+}$  to  $K_{B}^{-}(m)$  then checks  $K_{B}^{+}(K_{B}(m)) = m$ .
- If K<sup>+</sup><sub>B</sub>(K<sup>-</sup><sub>B</sub>(m)) = m, whoever signed m must have used Bob's private key.
  - Alice thus verifies that:
    - Bob signed m
    - no one else signed m
    - Bob signed m and not m'

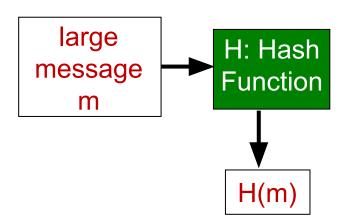
non-repudiation:

 Alice can take m, and signature K<sub>B</sub>(m) to court and prove that Bob signed m



computationally expensive to public-key-encrypt long messages

- *goal:* fixed-length, easy- tocompute digital "fingerprint"
- apply hash function H to m, get fixed size message digest, H(m).



Hash function properties:

- many-to-1
- produces fixed-size msg digest (fingerprint)
- given message digest x, computationally infeasible to find m such that x = H (m)

#### Internet checksum: poor crypto hash function

Internet checksum has some properties of hash function:

- produces fixed length digest (16-bit sum) of message
- is many-to-one

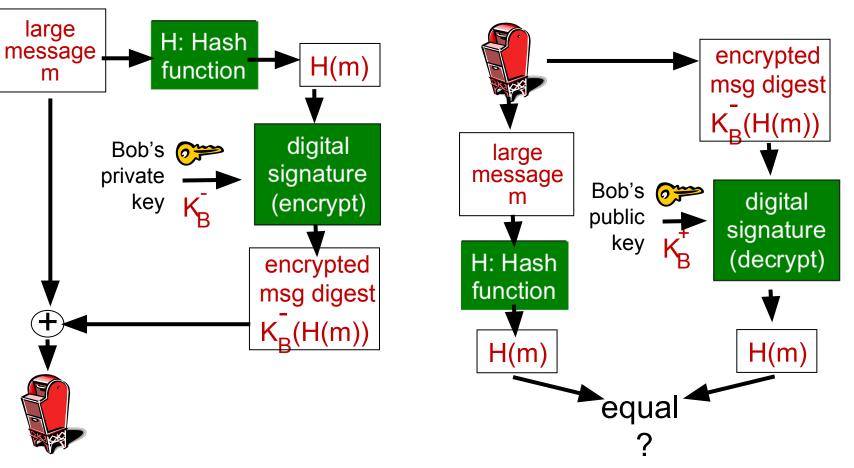
But given message with given hash value, it is easy to find another message with same hash value:

<u>message</u>	ASCII format	<u>message</u>	ASCII format
I O U 1	49 4F 55 31	I O U <u>9</u>	49 4F 55 <u>39</u>
00.9	30 30 2E 39	0 0 . <u>1</u>	30 30 2E <u>31</u>
9 B O B	39 42 D2 42	9 B O B	39 42 D2 42
	B2 C1 D2 AC	different messages —	B2 C1 D2 AC
	k	out identical checksums!	

#### Digital signature = signed message digest

Bob sends digitally signed message:

Alice verifies signature, integrity of digitally signed message:

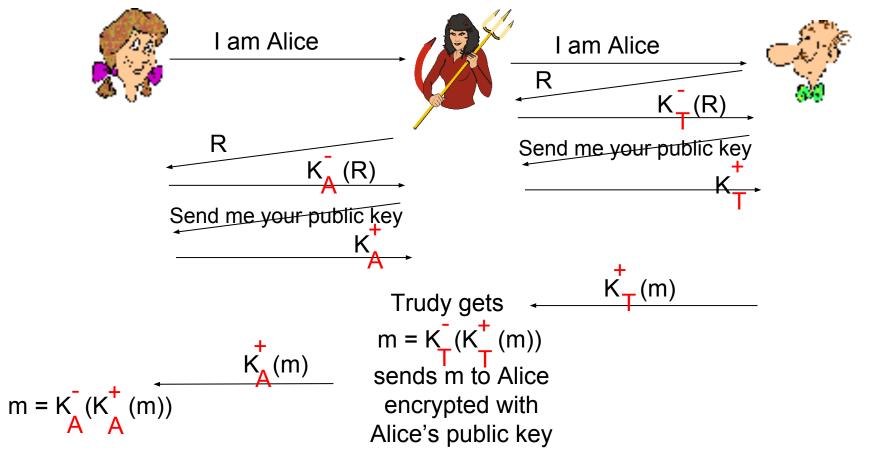


### Hash function algorithms

- MD5 hash function widely used (RFC 1321)
  - computes 128-bit message digest in 4-step process.
  - arbitrary 128-bit string x, appears difficult to construct msg m whose MD5 hash is equal to x
- SHA-1 is also used
  - US standard [NIST, FIPS PUB 180-1]
  - 160-bit message digest

## Recall: ap5.0 security hole

*man (or woman) in the middle attack:* Eve poses as Alice (to Bob) and as Bob (to Alice)

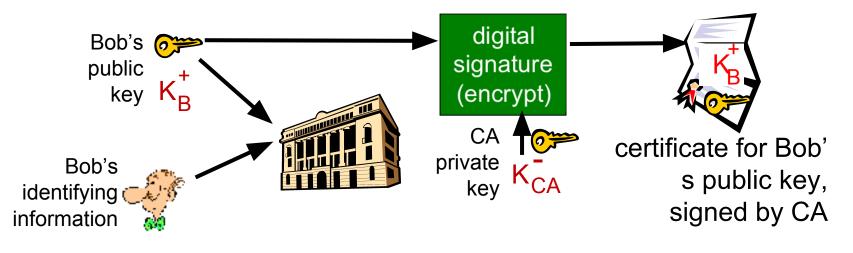


#### **Public-key certification**

- motivation: Eve plays pizza prank on Bob
  - Eve creates e-mail order: Dear Pizza Store, Please deliver to me four pepperoni pizzas. Thank you, Bob
  - Eve signs order with her private key
  - Eve sends order to Pizza Store
  - Eve sends to Pizza Store her public key, but says it's Bob's public key
  - Pizza Store verifies signature; then delivers four pepperoni pizzas to Bob
  - Bob doesn't even like pepperoni

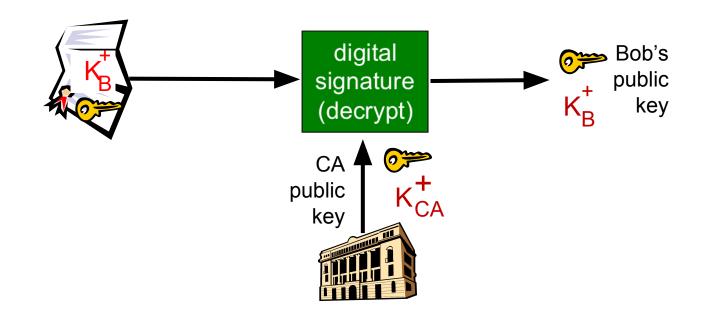
### **Certification authorities**

- certification authority (CA): binds public key to particular entity, E.
- ✤ E (person, router) registers its public key with CA.
  - E provides "proof of identity" to CA.
  - CA creates certificate binding E to its public key.
  - certificate containing E's public key digitally signed by CA CA says "this is E's public key"



#### **Certification authorities**

- when Alice wants Bob's public key:
  - gets Bob's certificate (Bob or elsewhere).
  - apply CA's public key to Bob's certificate, get Bob's public key

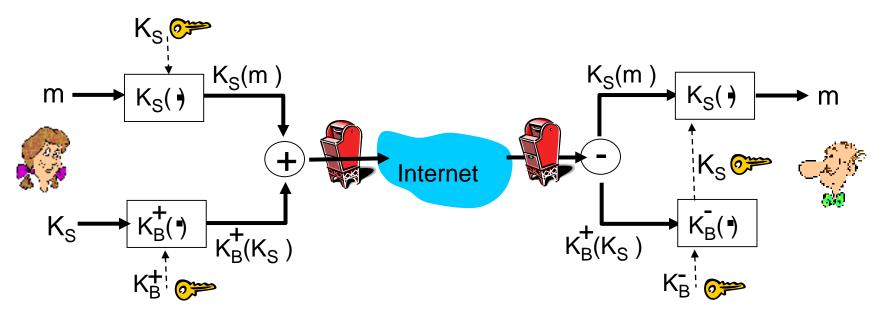


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#### Secure e-mail

Alice wants to send confidential e-mail, m, to Bob.

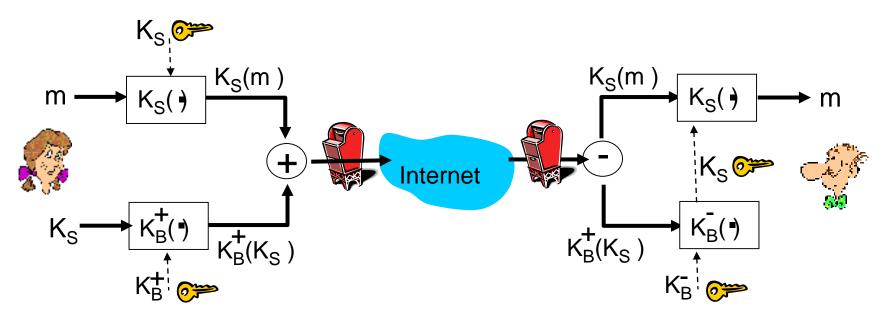


#### Alice:

- \* generates random symmetric private key, K<sub>S</sub>
- encrypts message with K<sub>s</sub> (for efficiency)
- also encrypts K<sub>s</sub> with Bob's public key
- \* sends both  $K_{S}(m)$  and  $K_{B}(K_{S})$  to Bob

#### Secure e-mail

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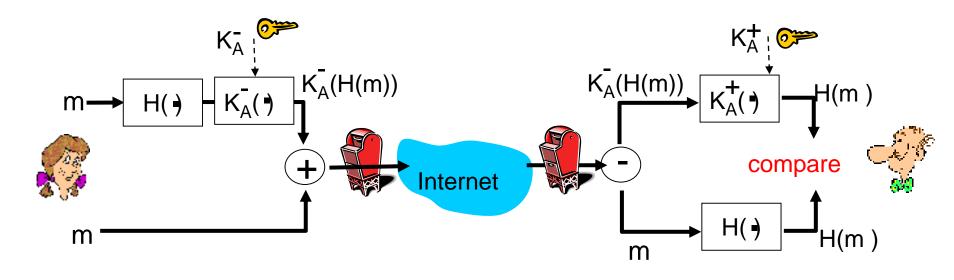


#### Bob:

- $\diamond$  uses his private key to decrypt and recover K<sub>s</sub>
- uses  $K_s$  to decrypt  $K_s(m)$  to recover m

#### Secure e-mail (continued)

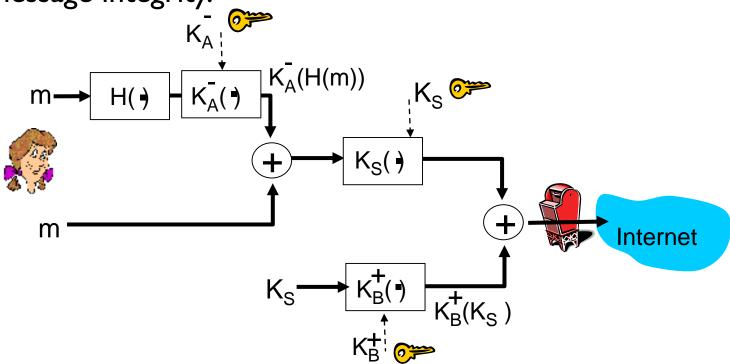
Alice wants to provide sender authentication message integrity



- Alice digitally signs message
- sends both message (in the clear) and digital signature

#### Secure e-mail (continued)

Alice wants to provide secrecy, sender authentication, message integrity.



Alice uses three keys: her private key, Bob's public key, newly created symmetric key