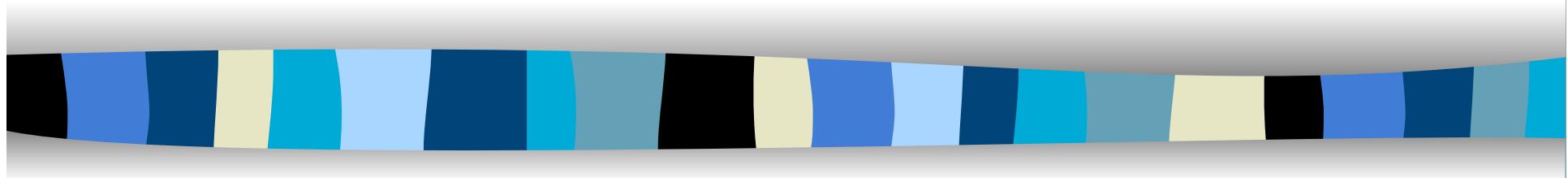


Formal Methods in software development



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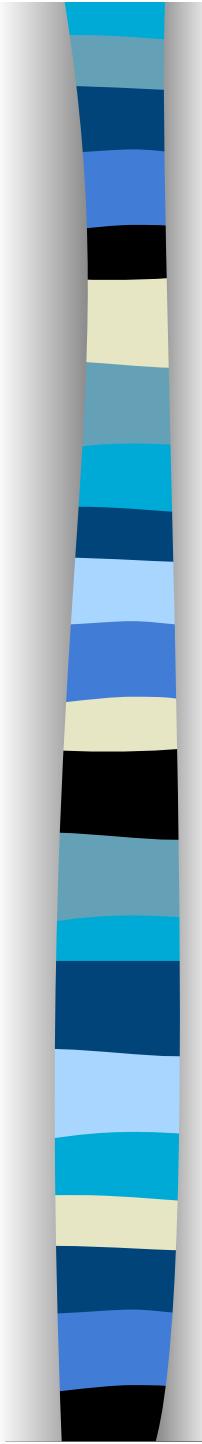
Looking for a program: example

Given ($|x > 0|$) $S (|y. y < x|)$

many possible solutions for S :

$y = 0;$

```
y = 0;  
while (y * y < x) {  
    y = y + 1;  
}  
y = y - 1;
```



Example (assignment and concatenation)

Write down a program P such that

- (a) $(\top) P (y = x + 2)$
- (b) $(\top) P (z > x + y + 4)$

Example (assignment and concatenation)

```
(| T |)
(| x+1 + 1 = x + 2 |)
t = x + 1;
(| t + 1 = x + 2 |)
z = t + 1;
(| z = x + 2 |)
y = z;
(| y = x + 2 |)
```

Example (if then else)

For each of the specifications below, write code for P and prove the partial correctness of the specified input/output behaviour:

- (a) $(\top) P \{z = \max(w, x, y)\}$, where $\max(w, x, y)$ denotes the largest of w , x and y .
- (b) $(\top) P \{((x = 5) \rightarrow (y = 3)) \wedge ((x = 3) \rightarrow (y = -1))\}$.

Using the method bottom-up

Syntesizing the program computing n^3

We look for assignments and build loop

First postcondition

$$c=N^3$$

We transform it in order to obtain the postcondition of a loop $I \wedge \neg C$

$$c=x^3 \wedge x=N$$

Metaprogram

1. establish I
2. while $x \neq N$ do
3. Preserve I while making x closer to N done;
invariant I $c=x^3$ variant $N-x$

Syntesizing the program computing n^3

3. Preserve I while making x closer to N done;
invariant I $c=x^3$ variant $N-x$

becomes

3. $x, c := x+1, E$ done;
invariant I $c=x^3$ variant $N-x$

Where , to preserve invariant

$I \wedge x \neq N \Rightarrow [x, c := x+1, E] \ I$
invariant I $c=x^3$ variant $N-x$

Syntesizing the program computing n^3

```
[x,c := x+1, E] I =
[x,c := x+1, E] c=x3 =
E =(x+1)3 =
E =x3+3x2+3x+1 =
E =c+3x2+3x+1

E =c+d
d=3x2+3x+1
invariant I ∧ I2 c=x3 ∧ d=3x2+3x+1 variant N-x
```

Where , to preserve invariant

$$c=x^3 \wedge d=3x^2+3x+1 \Rightarrow [x,c := x+1, c+d] (c=x^3)$$

But is I_2 invariant?

What is d ?

Syntesizing the program computing n^3

```
[x,c,d := x+1, c+d, E'] I2 =  
E' = 3(x2+2x+1) + 3(x+1) + 1 =  
E' = d+6x+6
```

```
E' = d+e  
e = 6x+6
```

invariant $I \wedge I_2 \wedge I_3 \quad c=x^3 \wedge d=3x^2+3x+1 \wedge e=6x+6$ variant $N-x$

```
[x,c,d,e := x+1, c+d, d+e, E''] I3 =  
E'' = 6(x+1)+6 =  
E'' = e+6
```

inizialization

```
[x,c,d,e := 0, c, d, E] I
```

Using the method bottom-up

Syntesizing the program computing n^3

Program

```
1. x,c,d,e := 0,0,1,6;  
2. while x≠N do  
3.   x,c,d,e := x+1,c+d,d+e,e+6 done;
```



Bibliography

Lessons 1-4 Propositional Natural deduction (Huth Ryan chapter 1)

Lesson 5 OBDD (Huth Ryan chapter 6)

Lessons 6-7 Predicate Natural deduction (Huth Ryan chapter 2)

Lessons 8-10 Temporal logic (Huth Ryan chapter 3)

Lessons 11-13 Denotational semantics (dens.pdf)

Lesson 14 Operational semantics and languages

Lesson 15 Reactive systems

Lesson 16 Bisimulations (Milner: Communicating and mobile systems, intro2ccs.pdf)

Lesson 17 Hennessy Milner Logic CCS

Lesson 18 Hoare logic (Huth Ryan chapter 4)

Lesson 19 Synthesysing programs (Monin)