Intensive Computation

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- Residue number systems are based on the congruence relation:
 - Two integers a and b are said to be congruent modulo m if m divides exactly the difference of a and b
 - We write $a \equiv b \pmod{m}$
- For example
 - $10 \equiv 7 \pmod{3}$
 - $10 \equiv 4 \pmod{3}$
 - $10 \equiv 1 \pmod{3}$
 - $10 \equiv -2 \pmod{3}$
- The number m is a modulus or base, and we assume that its values exclude 1, which produces only trivial congruences

- In fact:
- If q and r are the quotient and remainder, respectively, of the integer division of a by m that is: a = q:m + r
 - \rightarrow then, by definition, we have $a \equiv r \pmod{m}$
- The number r is said to be the *residue* of a with respect to m, and we shall usually denote this by $r = |a|_m$
- The set of m smallest values, $\{0; 1; 2; ...; m-1\}$, that the residue may assume is called the set of *least positive residues modulo m*

- Suppose we have a set, $\{m_1; m_2; ...; m_N\}$, of N positive and pairwise **relatively prime** moduli
- Let M be the product of the moduli $M=m_1xm_2x...xm_N$
- We write the representation in the form $\langle x1; x2; ...; xN \rangle$, where $xi = |X|_{mi}$, and we indicate the relationship between X and its residues by writing $X \approx \langle x1; x2; ...; xN \rangle$
- Example: in the residue system {2, 3, 5}, M=30 and

- Every number X < M has a unique representation in the residue number system, which is the sequence of residues
 |X|_{mi}: 1 ≤ i ≤ N>
- A partial proof of uniqueness is as follows:
 - Suppose X_1 and X_2 are two different numbers with the **same** *residue representation*
 - Then $|X_1|_{mi} = |X_2|_{mi}$, and so $|X_1 X_2|_{mi} = 0$
 - Therefore $X_1 X_2$ is the least common multiple (**lcm**) of *mi*
 - But if the mi are relatively prime, then their lcm is M, and it must be that $X_1 X_2$ is a multiple of M
 - So it cannot be that $X_1 < M$ and $X_2 < M$
 - Therefore, the representation $< |X|_{mi} : 1 \le i \le N >$ is unique and may be taken as the representation of X

- The number *M* is called the *dynamic range* of the RNS, because the number of numbers that can be represented is *M*
- For unsigned numbers, that range is [0;M 1]
- Representations in a system in which the moduli are not pairwise relatively prime will be not be unique: two or more numbers will have the same representation

Example

	Relatively prime			Relatively non-prime		
N	m1=2	m2=3	m3=5	m1=2	m2=4	m3=6
0	0	0	0	0	0	0
1	1	1	1	1	1	1
2	0	2	2	0	2	2
3	1	0	3	1	3	3
4	0	1	4	0	0	4
5	1	2	0	1	1	5
6	0	0	1	0	2	0
7	1	1	2	1	3	1
8	0	2	3	0	0	2
9	1	0	4	1	1	3
10	0	1	0	0	2	4
11	1	2	1	1	3	5
12	0	0	2	0	0	0
13	1	1	3	1	1	1
14	0	2	4	0	2	2
15	1	0	0	1	3	3

- We defined *standard residue number systems*
- There are also examples of non-standard RNS, the most common of which are the redundant residue number systems
- Such a system is obtained by, essentially, adding extra (redundant) moduli to a standard system
- The dynamic range then consists of a legitimate range, defined by the non-redundant moduli and an illegitimate range
- Redundant number systems of this type are especially useful in fault-tolerant computing

- Ignoring other, more practical, issues, the best moduli are probably prime numbers
- For computer applications, it is important to have moduli-sets that facilitate both efficient representation and balance, meaning that the differences between the moduli should be as small as possible

- Take, for example, the choice of 13 and 17 for the moduli that are adjacent prime numbers
- The dynamic range is 221
- With a straightforward binary encoding:
 - 4 bits will be required to represent 13
 - 5 bits will be required to represent 17

- The representational efficiency is:
 - In the first case 13/16
 - In the second case is 17/32
- If instead we chose 13 and 16, then the representational efficiency:
 - is improved to 16/16 in the second case
 - but at the cost of reduction in the range (down to 208)
- With the better balanced pair, 15 and 16, we would have:
 - a better efficiency 15/16 and 16/16
 - A greater range: 240

- It is also useful to have *moduli that simplify* the implementation of the *arithmetic operations*
- This means that arithmetic on residue digits should not deviate too far from conventional arithmetic, which is just arithmetic modulo a power of two
- A common choice of prime modulus that does not complicate arithmetic and which has good representational efficiency is $mi = 2^i 1$

- Not all pairs of numbers of the form $2^i 1$ are relatively prime
- It can be shown that that 2^j 1 and 2^k 1 are relatively prime if and only if j and k are relatively prime
- For example:

•
$$2^4-1=15$$

$$15 = 3x5$$

•
$$2^5-1=31$$

•
$$2^6 - 1 = 63$$

$$63 = 3x7$$

•
$$2^{7}$$
-1= 127

•
$$2^{8}-1=255$$

- Many moduli sets are based on these choices, but there are other possibilities; for example, moduli-sets of the form $\{2^n-1; 2^n; 2^n+1\}$ are among the most popular in use
- At least four considerations for the selection of moduli
 - The selected moduli must provide an adequate range whilst also ensuring that RNS representations are unique
 - The efficiency of binary representations; a balance between the different moduli in a given moduli-set is also important
 - The **implementations of arithmetic units** for RNS should to some extent be compatible with those for conventional arithmetic, especially given the legacy that exists for the latter
 - The size of individual moduli

- One of the primary advantages of RNS is that certain RNSarithmetic operations do not require carries between digits
- But, this is so only between digits
- Since a digit is ultimately represented in binary, there will be carries between bits, and therefore it is important to ensure that digits (>> the moduli) are not too large

- Small digits make it possible to realize cost-effective tablelookup implementations of arithmetic operations
- But, on the other hand, if the moduli are small, then a large number of them may be required to ensure a sufficient dynamic range
- The choices depend on applications and technologies

Negative numbers

- As with the conventional number systems, any one of the radix complement, diminished-radix complement, or sign-andmagnitude notations may be used in RNS
- The merits and drawbacks of choosing one over the other are similar to those for the conventional notations
- However, the determination of sign is much more difficult with the residue notations, as is magnitude-comparison
- This problem imposes many limitations on the application of RNS and we deal with just the positive numbers

- Addition/subtraction and multiplication are easily implemented with residue notation, depending on the choice of the moduli
- Division is much more difficult due to the difficulties of signdetermination and magnitude-comparison

- Residue addition is carried out by individually adding corresponding digits
- A carry-out from one digit position is not propagated into the next digit position
- As an example, with the moduli-set {2; 3; 5; 7}:
 - the representation of 17 is <1; 2; 2; 3>
 - the representation of 19 is <1; 1; 4; 5>
 - adding the two residue numbers yields <0; 0; 1; 1>, which is the representation for 36 in that system

- Subtraction may be carried out by negating (in whatever is the chosen notation) the subtrahend and adding to the minuend
- This is straightforward for numbers in diminished-radix complement or radix complement notation
- For sign-and-magnitude representation, a slight modification of the algorithm for conventional sign-and-magnitude is necessary:
 - the sign digit is fanned out to all positions
 - addition proceeds as in the case for unsigned numbers but with a conventional sign-and-magnitude algorithm.

- Multiplication too can be performed simply by multiplying corresponding residue digit-pairs, relative to the modulus for their position → multiply digits and ignore or adjust an appropriate part of the result
- As an example, with the moduli-set {2; 3; 5; 7}:
 - 17 > <1; 2; 2; 3>
 - 19 > <1; 1; 4; 5>
 - their product, 323 is <1; 2; 3; 1>

- Basic fixed-point division consists, essentially, of a sequence of subtractions, magnitude-comparisons, and selections of the quotient-digits
- But comparison in RNS is a difficult operation, because RNS is not positional or weighted
- Example:
 - moduli-set {2; 3; 5; 7}
 - the number represented by <0; 0; 1; 1> is almost twice that represented by <1; 1; 4; 5>
 - but this is far from apparent

Conversion

- The most direct way to convert from a conventional representation to a residue one is to divide by each of the given moduli and then collect the remainders, forward conversion
- This is a costly operation if the number is represented in an arbitrary radix and the moduli are arbitrary
- If number is represented in radix-2 (or a radix that is a power of two) and the moduli are of a suitable form (e.g. 2^n-1), then these procedures that can be implemented with more efficiency

Conversion

- The conversion from residue notation to a conventional notation - reverse conversion - is more difficult (conceptually, if not necessarily in the implementation) and so far has been one of the major impediments to the adoption use of RNS
 - One way in which it can be done is to assign weights to the digits of a residue representation and then produce a positional (weighted) mixedradix representation that can then be converted into any conventional form
 - Another approach involves the use of the Chinese Remainder Theorem, which is the basis for many algorithms for conversion from residue to conventional notation