

Eigenvalues, Eigenvectors and Applications

Intensive Computation

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INTRODUCTION

Eigenvalue Problems

- Solving linear systems $A\mathbf{x} = \mathbf{b}$ is one part of numerical linear algebra, and involves manipulating the rows of a matrix
- The second main part of numerical linear algebra is about find **eigenvalues** and **eigenvectors**
- This is done by *transforming a matrix to leave its eigenvalues unchanged*
- The standard algebraic **eigenvalue problem** is:
Given an $n \times n$ matrix A , find a scalar and a nonzero vector x such that

$$A\mathbf{x} = \lambda\mathbf{x}$$

where λ is an **eigenvalue** of A and non-zero x is the corresponding **eigenvector**

Eigenvalue Problems

Eigenvalue problems occur in many areas of science and engineering:

- The *natural modes and frequencies of vibration* of a structure are determined by the eigenvectors and eigenvalues of an appropriate matrix
- The *stability of the structure* is determined by the locations of the eigenvalues, thus their computation is of critical interest
- Eigenvalues can be very useful in *analyzing numerical methods* (the convergence analysis of iterative methods for solving systems of algebraic equations, and the stability analysis of methods for solving systems of differential equations)
- *Graph theory*

Eigenvalue Problems

- An **eigenvector** of a matrix **determines a direction** in which the *effect of the matrix is particularly simple*:
 - The **matrix expands or shrinks** any vector lying in that direction by a scalar multiple, and
 - the **expansion or contraction factor** is given by the **corresponding eigenvalue**
- Thus, **eigenvalues and eigenvectors** provide a means of **understanding the complicated behavior** of a general linear transformation by decomposing it into simpler actions

Eigenvalue Problems

- Although many examples involve only **real matrices**, both the theory and computational procedures are generally applicable to **complex matrices**
- The notation difference for complex matrices is that the **conjugate transpose**, A^H , is used instead of the **transpose**, A^T
- The set of **all the eigenvalues** of a matrix A , denoted by $\lambda(A)$ is called the **spectrum** of A
- The **maximum modulus** of the eigenvalues, $\max\{|\lambda| : \lambda \text{ in } \lambda(A)\}$, is called the **spectral radius** of A , denoted by $\rho(A)$

Eigenvalue Problems

- The equation $Ax = \lambda x$ is equivalent to

$$(A - \lambda I)x = 0$$

- This homogeneous equation has a **nonzero solution** x if and only if its matrix is **singular**, that is the eigenvalues of A are the values λ such that

$$\det(A - \lambda I) = 0$$

- $\det(A - \lambda I)$ is a polynomial of degree n in λ , the **characteristic polynomial** of A , and *its roots are the eigenvalues* of A

Eigenvalue Problems

- An $n \times n$ matrix A always has n eigenvalues (*Fundamental Theorem of Algebra*), but they need be neither *distinct* nor *real*
- The product of the eigenvalues is $\det A = \prod_{i=1}^n \lambda_i$
- The sum of the eigenvalues is $\sum_{i=1}^n a_{ii} = \sum_{i=1}^n \lambda_i$ called *trace*

POWER METHOD

Computing Eigenvalues and Eigenvectors

- Many numerical methods for **computing eigenvalues and eigenvectors** are based on **reducing the original matrix** to a simpler form, whose eigenvalues and eigenvectors are then **easily determined**
- **Finding the eigenvalues and eigenvectors** is equivalent to **transforming** the underlying system of equations into a special set of coordinate axes in which the **matrix is diagonal**
- The **eigenvalues** are the entries of the **diagonal** matrix
- The eigenvectors are the new set of coordinate axes

Similarity Transformations

- We need to identify:
 - what types of transformations preserve eigenvalues
 - for what types of matrices the eigenvalues are easily determined
- A matrix B is *similar* to a matrix A if there is a nonsingular matrix T such that

$$B = T^{-1}AT$$

- Then

$$By = \lambda y \rightarrow T^{-1}ATy = \lambda y \rightarrow A(Ty) = \lambda(Ty)$$

so that **A and B have the same eigenvalues**, and if y is an eigenvector of B , then **$x = Ty$ is an eigenvector of A**

Similarity Transformations

- Similarity transformations:
 - Preserve eigenvalues
 - Do not preserve eigenvectors
 - But the eigenvectors are still easily recovered
- Note that the converse is not true
 - two matrices that are **similar** must have the **same eigenvalues**
 - but two matrices that have the **same eigenvalues** are *not necessarily similar*

Similarity Transformations

- The **eigenvalues of a diagonal matrix are its diagonal entries**, and the eigenvectors are the corresponding columns of the identity matrix I
- Note that:
 - Diagonal form simplifies eigenvalue problems for general matrices by similarity transformations
 - But **some matrices cannot be transformed into diagonal form** by a similarity transformation
- Fortunately:
 - **every matrix** can be transformed into **triangular** by a similarity transformation
 - The **eigenvalues of a triangular matrix are also the diagonal entries**

Computing Eigenvalues and Eigenvectors

- There are several methods that are designed to compute all of the eigenvalues of a matrix and require a great deal of work
- In practice, one may need **only one or a few eigenvalues** and corresponding eigenvectors
- The simplest method for computing a single eigenvalue and eigenvector of a matrix is the **power method**, which takes successively higher powers of the matrix times an initial starting vector

Power Method

- Assume that the matrix has a **unique eigenvalue λ_1 of maximum modulus**, with corresponding eigenvector u_1
- Let us consider the iteration scheme starting from a **given nonzero vector x_0**

$$x_k = Ax_{k-1}$$

- The sequence **converges** to a multiple of u_1 , the eigenvector corresponding to the dominant eigenvalue λ_1

Power Method

- In fact, if we express the starting vector x_0 as a linear combination, $x_0 = \sum_{i=1}^n \alpha_i u_i$ where u_i are eigenvectors of A , then

$$\begin{aligned} x_k &= Ax_{k-1} = A^2 x_{k-2} = \cdots = A^k x_0 = A^k \sum_{i=1}^n \alpha_i u_i = \\ &= \sum_{i=1}^n \alpha_i A^k u_i = \sum_{i=1}^n \lambda_i^k \alpha_i u_i = \lambda_1^k (\alpha_1 u_1 + \sum_{i=2}^n (\lambda_i / \lambda_1)^k \alpha_i u_i) \end{aligned}$$

- Since $|\lambda_i / \lambda_1| < 1$ for $i > 1$, successively higher powers go to zero, leaving only the component corresponding to u_1

Power Method

- To avoid eventual **overflow** (or underflow if the dominant eigenvalue is less than 1 in magnitude), it is better to **normalize** the approximate eigenvector at each iteration
- We can require its largest component to have modulus 1
- This step gives the **iteration scheme**

$$y_k = Ax_{k-1}$$

$$x_k = y_k / \|y_k\|_\infty$$

- With this normalization $\|y_k\|_\infty \rightarrow |\lambda_1|$ and $x_k \rightarrow u_1 / \|u_1\|_\infty$

Deflation Methods

- Suppose that an eigenvalue λ_1 and corresponding eigenvector x_1 for a matrix A have been computed
- We can **compute additional eigenvalues** $\lambda_2, \dots, \lambda_n$ of A , by a process called **deflation**, which removes the known eigenvalue
- Construct a matrix B with eigenvalues $0, \lambda_2, \dots, \lambda_n$ (*deflate* the matrix A , remove λ_1).
- Then λ_2 can be obtained by the power method

Deflation Methods

- Let H be any nonsingular matrix such that $Hx_1 = \alpha e_1$, a scalar multiple of the first column of the identity matrix I
- Then the **similarity transformation** determined by H transforms A into the form

$$HAH^{-1} = \begin{bmatrix} \lambda_1 & b^T \\ 0 & B \end{bmatrix}$$

- where B is a matrix of order $n - 1$ having eigenvalues $\lambda_2, \dots, \lambda_n$
- For example, good choice for H can be an appropriate Householder transformation

Deflation Methods

- We can work with B to compute the **next eigenvalue** λ_2 and the **eigenvector** y_2
- To compute the **second eigenvector** x_2 of matrix A , we need to add an element to vector y_2 (that consist of $n-1$ elements), that is $x_2 = (s_2 \ y_2)$
- s_2 can be the element

$$x_2 = H^{-1} \begin{bmatrix} \alpha \\ y_2 \end{bmatrix} \quad \text{where} \quad \alpha = \frac{b^T y_2}{\lambda_2 - \lambda_1}$$

- is an eigenvector corresponding to λ_2 for the original matrix A , provided $\lambda_2 \neq \lambda_1$
- **Process can be repeated to find additional eigenvalues and eigenvectors**

Deflation Methods

An **alternative approach** is the following:

- A given, λ_1 and x_1 known (e.g., by power method)
- Denote by a^T the first row of A (or the p -th row of A), i.e.,
 $a^T = (a_{11}, a_{12}, \dots, a_{1n})$
- Consider $B = A - x_1^* a^T$
- Where x_1^* is the vector x_1 normalized by dividing by its first element (and $x_1^* a^T$ is an $n \times n$ matrix)
- We can verify that $B x_1 = 0$ that is **0 is eigenvalue of B** (instead of λ_1) and **$\lambda_2, \dots, \lambda_n$ are still eigenvalues (of A)**
- Then we compute the value of λ_2 by the power method

Deflation Methods

A variant of the previous **alternative approach** is:

- A given, λ_1 and x_1 known (e.g., by power method)
- Denote by v_1 any vector such that $v_1^T x_1 = \lambda_1$
- Then the matrix $A - x_1 v_1^T$ has eigenvalues of $0, \lambda_2, \dots, \lambda_n$
- Then we compute the value of λ_2 by the power method
- There are several possible choices for v_1

Inverse Iteration

- For some applications, the **smallest eigenvalue** of a matrix is required rather than the largest
- We use the fact that the *eigenvalues of A^{-1} are the reciprocals of those of A*
- Hence the smallest eigenvalue of A is the reciprocal of the largest eigenvalue of A^{-1}

- In fact
$$\min_{i=1,\dots,n} |\lambda_i(A)| = \min_{i=1,\dots,n} \left| \frac{1}{\lambda_i(A^{-1})} \right| = \frac{1}{\max_{i=1,\dots,n} |\lambda_i(A^{-1})|}$$

- We therefore use the inverse iteration scheme

$$Ay_k = x_{k-1} \Rightarrow y_k = A^{-1}x_{k-1}$$

Convergence

- The **convergence rate** of the power method depends on the ratio $|\lambda_2|/|\lambda_1|$, where λ_2 is the eigenvalue having second-largest modulus
- **The smaller $|\lambda_2|/|\lambda_1|$, the faster the convergence**
- Hence the power method will converge
 - quickly if $|\lambda_2|/|\lambda_1|$ is small
 - slowly if $|\lambda_2|/|\lambda_1|$ is close to 1

EXAMPLES

Gould Index of Accessibility

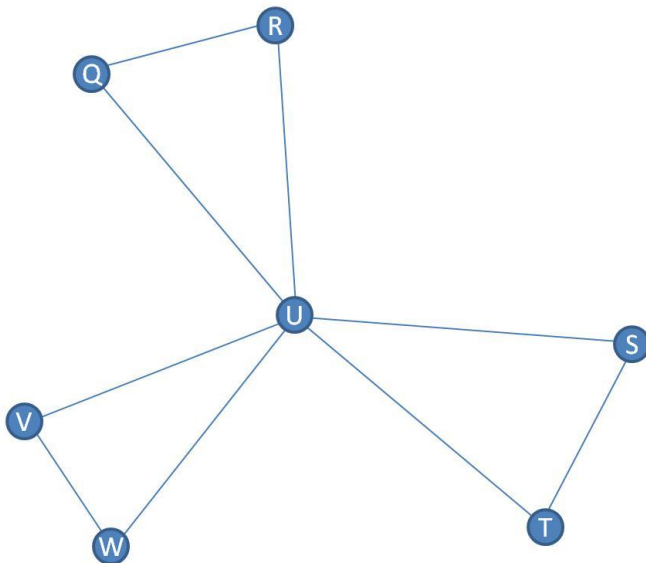
- https://www.math.washington.edu/~morrow/336_11/papers/leo.pdf
- <http://matrixapps.blogspot.it/2010/07/gould-index-matrix-application-to.html>
- The method proposed by Peter Gould (1967), also known as eigenvector centrality, is one method of computing the "centrality", or approximate importance, of each node in a graph
- The assumption is that each node's centrality is the sum of the centrality values of the nodes that it is connected to

Gould Index of Accessibility

- First we look at the method, then we will attempt to justify it
- We begin with the **adjacency matrix** A of the graph
- It is usual to define the entries a_{ii} , the diagonal, as 0
- We **replace the diagonal zeros with ones**
- The index that Gould proposes uses the (normalized) **eigenvector from the principle eigenvalue** of the **modified adjacency matrix** $B=A+I$
- The i th entry corresponds to the i th vertex and this is its **accessibility rank**

Gould Index of Accessibility

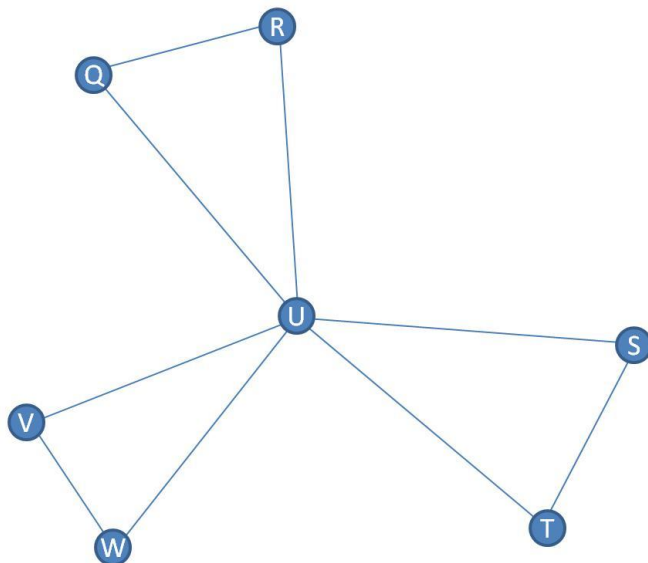
- Consider a graph that represents a set of towns (the vertices) and the travel routes between those towns (the edges)
- Historical geographers were interested in which town would become the trade center for this region
- Make the **adjacency matrix** for the graph and **place a 1 in each diagonal position**



<i>Q</i>	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>	
1	1	0	0	1	0	0	<i>Q</i>
1	1	0	0	1	0	0	<i>R</i>
0	0	1	1	1	0	0	<i>S</i>
0	0	1	1	1	0	0	<i>T</i>
1	1	1	1	1	1	1	<i>U</i>
0	0	0	0	1	1	1	<i>V</i>
0	0	0	0	1	1	1	<i>W</i>

Gould Index of Accessibility

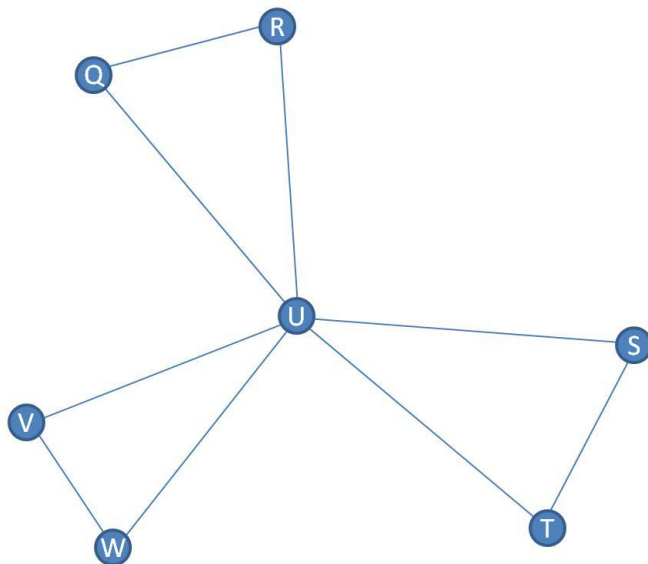
- Find the **largest eigenvalue** of the matrix
- We get eigenvalues (2, 0, 4, -1, 0, 2, 0), where the third one, 4, has the largest absolute value
- Find the **eigenvector associated** with the eigenvalue of 4: (0.3162, 0.3162, 0.3162, 0.3162, 0.6325, 0.3162, 0.3162)



<i>Q</i>	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>	
1	1	0	0	1	0	0	<i>Q</i>
1	1	0	0	1	0	0	<i>R</i>
0	0	1	1	1	0	0	<i>S</i>
0	0	1	1	1	0	0	<i>T</i>
1	1	1	1	1	1	1	<i>U</i>
0	0	0	0	1	1	1	<i>V</i>
0	0	0	0	1	1	1	<i>W</i>

Gould Index of Accessibility

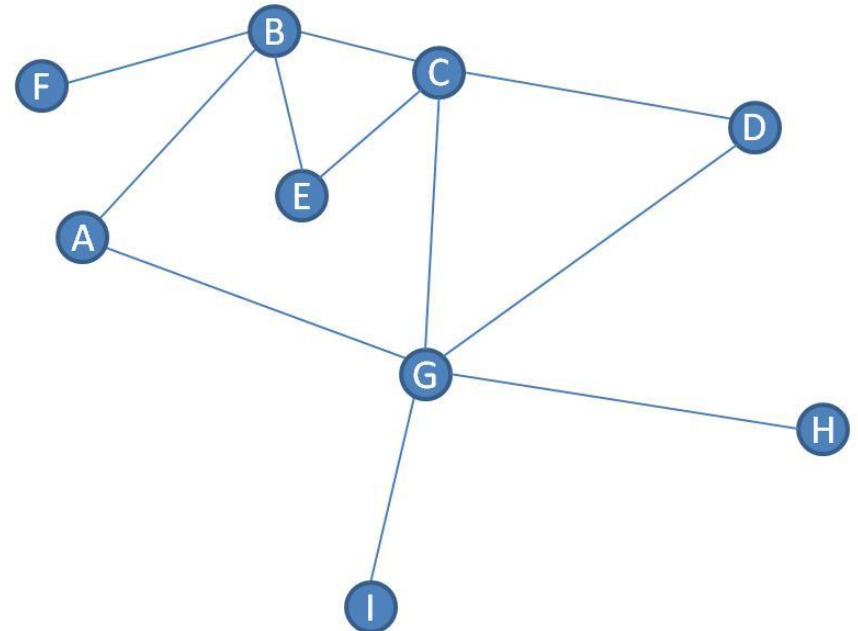
- You can normalize this vector by dividing by the sum of the entries, 2.5297, to get
 $(Q, R, S, T, U, V, W) = (0.125, 0.125, 0.125, 0.125, 0.25, 1.25, 1.25)$
- These are the Gould indices of each of the vertices
- They describe how strongly each vertex is connected to the other vertices



<i>Q</i>	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>	
1	1	0	0	1	0	0	<i>Q</i>
1	1	0	0	1	0	0	<i>R</i>
0	0	1	1	1	0	0	<i>S</i>
0	0	1	1	1	0	0	<i>T</i>
1	1	1	1	1	1	1	<i>U</i>
0	0	0	0	1	1	1	<i>V</i>
0	0	0	0	1	1	1	<i>W</i>

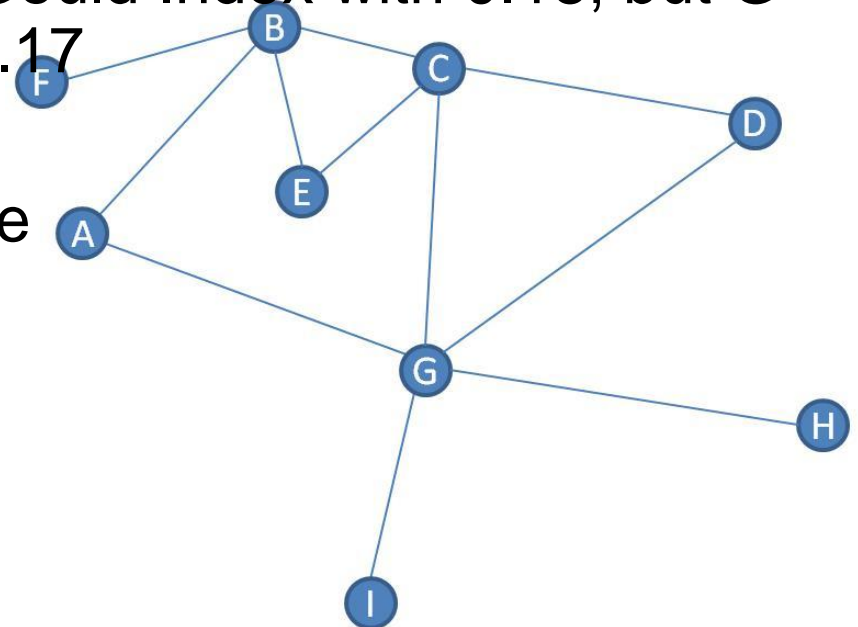
Gould Index of Accessibility

- Now consider a more realistic graph
- In this case, would G be the trade center, or would it be C?
- The eigenvalues which are (4.01, -1.37, 1.71, 1, -0.37, -0.56, 2.58, 1, 1)



Gould Index of Accessibility

- The eigenvector associated with the largest eigenvalue is (0.2941, 0.4097, 0.5023, 0.3249, 0.3026, 0.1359, 0.4770, 0.1583, 0.1583)
- Normalizing we have (A, B, C, D, E, F, G, H, I) = (0.1064, 0.1482, 0.1818, 0.1176, 0.1095, 0.0492, 0.1726, 0.0573, 0.0573)
- We see that C has the largest Gould Index with 0.18, but G comes in a close second with 0.17
- Note that I and H have the same smallest index, as expected



Algebraic Connectivity of graphs

- The importance of the **algebraic connectivity** of a graph is due to the fact that it is a good **parameter to measure**, to a certain extent, how well **a graph is connected**
- The algebraic connectivity is for
 - application on trees,
 - application on **hard problems in graph theory** (the expanding properties of graphs, weighted graphs, absolute algebraic connectivity, isoperimetric number, genus and other invariants of a graph)
 - the study of the **asymptotic behavior for random graphs**;
 - applications on **combinatorial optimization problems** (the maximum cut problem and the traveling salesman problem)
 - the **theory of elasticity**
 - the correspondence between continuous and discrete mathematics

Algebraic Connectivity of graphs

- Based on theory of Fiedler (1970s)
- We define the Laplacian matrix $L(G)$ of the graph $G(N,E)$ as:
 - $L(G)(i,i) = \text{degree of node } i \text{ (number of incident edges)}$
 - $L(G)(i,j) = -1$ if $i \neq j$ and there is an edge (i,j)
 - $L(G)(i,j) = 0$ otherwise
- We have $L(G)=D-A$ where D is the diagonal matrix of node degrees and A is the adjacency matrix of graph $G(N,E)$

Algebraic Connectivity of Graphs

- $L(G)$ is **symmetric** \rightarrow This means the **eigenvalues** of $L(G)$ are real and its **eigenvectors** are real and orthogonal
- The **eigenvalues** of $L(G)$ are **nonnegative**:
 - $0 = \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$
- The number of connected components of G is equal to the number of λ_i equal to 0
- G is connected if and only if $\lambda_2 \neq 0$

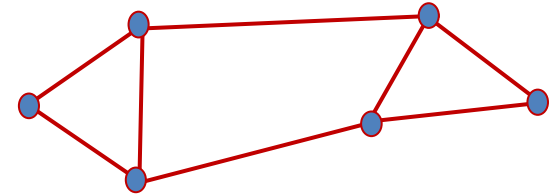
Algebraic Connectivity of Graphs

- The **second smallest eigenvalue** λ_2 of $L(G)$ is called **algebraic connectivity** of G
- The **eigenvector** associated with the algebraic connectivity has been named the ***Fiedler vector***
- The **Fiedler vector** can be used to **partition a graph**

Graph Partitioning

Spectral Bi-partitioning Algorithm

- Build Laplacian matrix L of the graph
- Find the **second smallest eigenvalue** and the corresponding **eigenvector**
- Map vertices to corresponding components of the Fiedler vector
- **Grouping**
 - Sort components of reduced 1-dimensional vector.
 - Identify clusters by splitting the sorted vector in two



Graph Partitioning

How to choose a splitting point?

- Split at 0 (positive and negative values)
- Mean or median value
- Partitioning a graph into k clusters can be done by **Recursive bi-partitioning** (Hagen et al., '91)
- Recursively apply bi-partitioning algorithm in a hierarchical divisive manner
- Example: Image Segmentation
 - Uses 2nd (smallest) eigenvector to define optimal cut
 - Recursively generates two clusters with each cut

Markov chain

- http://webSPACE.ship.edu/deensley/m318/ppt/Section_49.pdf
- A vector with nonnegative entries that add up to 1 is called a **probability vector**
- A **stochastic matrix** is a square matrix whose columns are probability vectors
- A **Markov chain** is a sequence of probability vectors x_0, x_1, \dots together with a stochastic matrix P , such that

$$x_1 = Px_0, x_2 = Px_1, x_3 = Px_2, \dots$$

Markov chain

Example

- Consider the following model of population movement between a city and the suburbs:
 - each year 5% of city dwellers move the suburbs and
 - 3% of suburbanites move to the city
- If in 2001 58.2% of the population lived in the city and 41.8% lived in the suburbs, what is the ***population distribution 20 years later?***

Markov chain

Example

- Consider matrix $M = \begin{pmatrix} 0.95 & 0.03 \\ 0.05 & 0.97 \end{pmatrix}$ and vector $x_0 = \begin{pmatrix} 0.582 \\ 0.418 \end{pmatrix}$

to represent the population distribution in 2001

- In this case $x_1 = Mx_0 = \begin{pmatrix} 0.565 \\ 0.435 \end{pmatrix}$ gives the population distribution in 2002
- In general $x_n = M^n x_0$ **gives the population distribution in n years after 2001**

Markov chain

- The difference between x_i and x_{i+1} gets smaller every step
- Let P be stochastic matrix. A **steady-state vector** is a probability vector x such that $Px=x$
- A **steady-state vector** is also called **equilibrium vector**
- If state x is achieved, the system stays there
- **A nonzero steady-state vector** is in fact an **eigenvector of eigenvalue 1 of P** (since it satisfies $Px=x$)

Markov chain

- The **eigenvalues** are $\lambda_1 = 1$ and $\lambda_2 = 0.92$, and the corresponding **eigenvectors** are

$$v_1 = \begin{pmatrix} -0.514496 \\ -0.857493 \end{pmatrix} \quad v_2 = \begin{pmatrix} -0.707107 \\ 0.707107 \end{pmatrix}$$

- So we form

$$P = \begin{pmatrix} 0.514496 & -0.707107 \\ -0.85749 & 0.707107 \end{pmatrix} \quad D = \begin{pmatrix} 1 & 0 \\ 0 & 0.92 \end{pmatrix}$$

Markov chain

- Now we have $Mx = \lambda x \rightarrow MP = PD \rightarrow M = PDP^{-1}$
- And $M \cdot M = PDP^{-1}PDP^{-1} = PD^2P^{-1}$
- In general $M^k = PD^kP^{-1}$
- **Having** $M^k = P \begin{pmatrix} 1 & 0 \\ 0 & (0.92)^k \end{pmatrix} P^{-1}$
- **It follows** $\lim_{k \rightarrow \infty} M^k = P \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} P^{-1} = \begin{pmatrix} 0.375 & 0.375 \\ 0.625 & 0.625 \end{pmatrix}$

Markov chain

- For any stochastic vector x we will have

$$\lim_{k \rightarrow \infty} M^k x = \begin{pmatrix} 0.375 & 0.375 \\ 0.625 & 0.625 \end{pmatrix} x =$$
$$\begin{pmatrix} 0.375 v_1 & 0.375 v_2 \\ 0.625 v_1 & 0.625 v_2 \end{pmatrix} = \begin{pmatrix} 0.375 & 0.375 \\ 0.625 & 0.625 \end{pmatrix}$$